RAID4S: Adding SSDs to RAID Arrays

RAID (Redundant Array of Independent Disks)

RAID algorithms provide better performance and/or data reliability by combining several disk drives.

RAID4 and RAID5 store the parity computed across data blocks to protect against a single disk failure.

Computing parity for small I/Os significantly reduces performance. A common technique for mitigating this problem is to transform small writes into larger writes. However some small writes are inevitable, particularly as disk drives fill up.

Solid State Drives (SSDs)

SSDs compared to disk drives: - Constructed of flash RAM thus not mechanical devices

- Faster random I/O
- Lower power consumption
- Wear out with each write
- More expensive

SSDs are not cheap enough to replace all disk drives. This work integrates SSDs into disk-based storage systems.

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Time

One sn - Read - Write - Data RAID4: RAID5 RAID4 access

- For M s - RAID
- RAID
- RAID
- RAID4S
- RAID
- RAID

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Small Write Analysis- M of - N = - N = <th></th>	
nall write = 4 I/Os old data and parity new data and parity and parity accessed in parallel : parity drive is bottleneck : data and parity accessed in parallel S: parity offloaded to SSD and all data es are parallelized small writes to N disks: 4: 2M $4S: 2 \lceil M/N \rceil$ $5: 2*2 \lceil M/(N+1) \rceil$ S speedup over:	Write Analysis
small writes to N disks:Throw4: $2M$ - 64k4S: $2 \lceil M/N \rceil$ - Each5: $2*2 \lceil M/(N+1) \rceil$ - OurS speedup over:- our	+ I/Os d parity nd parity ccessed in parallel is bottleneck arity accessed in parallel baded to SSD and all data lelized $- N = 4 d$ Assumption Each so Small wr parallelized
4: M/ [M/N] 5: 2 [M/(N+1)] / [M/N] - ₩e	to N disks: Through - 64KB s - Each I/ - Outstar complete - Hardwa - Western

Calculated Throughpu

tanding I/Os at the controller disk drives for RAID4S + 1 SSD 5 disk drives for RAID 4 and RAID5

tions:

small write incurs a seek and transfer time. rites are to different stripes. Thus, the I/O is zed as much as possible.

put calculation: small writes O incurs a seek and some transfer time Inding I/Os are queued at the controller and are ted in parallel, if possible.

are: Digital WD20EADS (low power disk drive) È 2.2

Figure 2. Throughput of small writes. RAID4S completes N small writes in the time that RAID4 completes one small write, as long as the SSD is N times faster than hard disks in the array.

RAID4S is up to 2X faster than RAID5 and 4X faster than RAID4

Conclusions

Flash SSDs provide better performance and new opportunities for data storage in distributed systems. This work replaces the parity disk drive in a RAID4 system with SSD. Initial results show improvements of up to (N/2)X speedup over RAID5 and NX speedup over RAID4.

Future Work

More general workloads - Mix small and large writes - Workload traces Build RAID4S with hard drives and flash and measure: - Throughput - Latency - Power Investigate reliability impact of RAID4S

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U	t												
	16	8.92709804	5.10119888	2.23177451									
	Throughput with Multiple Random Write Streams												
9		+	+		+								
75													
.5	•	0	0		0		+ 0	RAID 4S RAID 5 RAID 4					
25	- 0 -	-0											
0	0	5	1	0	15								
		:	# Outstar	nding I/Os									