

RECOMMENDED ELLIPTIC CURVES FOR FEDERAL GOVERNMENT USE

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This collection of elliptic curves is recommended for Federal government use and contains choices of private key length and underlying fields.

§1. PARAMETER CHOICES

1.1 Choice of Key Lengths

The principal parameters for elliptic curve cryptography are the elliptic curve E and a designated point G on E called the *base point*. The base point has order r , a large prime. The number of points on the curve is $n = fr$ for some integer f (the *cofactor*) not divisible by r . For efficiency reasons, it is desirable to take the cofactor to be as small as possible.

All of the curves given below have cofactors 1, 2, or 4. As a result, the private and public keys are approximately the same length. Each length is chosen to correspond to the cryptovvariable length of a common symmetric cryptologic. In each case, the private key length is, at least, approximately twice the symmetric cryptovvariable length.

1.2 Choice of Underlying Fields

For each cryptovvariable length, there are given two kinds of fields.

- A *prime field* is the field $GF(p)$ which contains a prime number p of elements. The elements of this field are the integers modulo p , and the field arithmetic is implemented in terms of the arithmetic of integers modulo p .

- A *binary field* is the field $GF(2^m)$ which contains 2^m elements for some m (called the *degree* of the field). The elements of this field are the bit strings of length m , and the field arithmetic is implemented in terms of operations on the bits.

The following table gives the sizes of the various underlying fields. By $\|p\|$ is meant the length of the binary expansion of the integer p .

<u>Symmetric CV Length</u>	<u>Example Algorithm</u>	<u>Prime Field</u>	<u>Binary Field</u>
80	SKIPJACK	$\ p\ = 192$	$m = 163$
112	Triple-DES	$\ p\ = 224$	$m = 233$
128	AES Small	$\ p\ = 256$	$m = 283$
192	AES Medium	$\ p\ = 384$	$m = 409$
256	AES Large	$\ p\ = 521$	$m = 571$

1.3 Choice of Basis

To describe the arithmetic of a binary field, it is first necessary to specify how a bit string is to be interpreted. This is referred to as choosing a *basis* for the field. There are two common types of bases: a *polynomial basis* and a *normal basis*.

- A polynomial basis is specified by an irreducible polynomial modulo 2, called the *field polynomial*. The bit string $(a_{m-1} \dots a_2 a_1 a_0)$ is taken to represent the polynomial

$$a_{m-1} t^{m-1} + \dots + a_2 t^2 + a_1 t + a_0$$

over $GF(2)$. The field arithmetic is implemented as polynomial arithmetic modulo $p(t)$, where $p(t)$ is the field polynomial.

- A normal basis is specified by an element θ of a particular kind. The bit string $(a_0 \ a_1 \ a_2 \ \dots \ a_{m-1})$ is taken to represent the element

$$a_0 \theta + a_1 \theta^2 + a_2 \theta^{2^2} + \dots + a_{m-1} \theta^{2^{m-1}}.$$

Normal basis field arithmetic is not easy to describe or efficient to implement in general, but is for a special class called *Type T low-complexity* normal bases. For a given field degree m , the choice of T specifies the basis and the field arithmetic (see Appendix 2).

There are many polynomial bases and normal bases from which to choose. The following procedures are commonly used to select a basis representation.

- *Polynomial Basis:* If an irreducible *trinomial* $t^m + t^k + 1$ exists over $GF(2)$, then the field polynomial $p(t)$ is chosen to be the irreducible trinomial with the lowest-degree middle term t^k . If no irreducible trinomial exists, then one selects instead a *pentanomial* $t^m + t^a + t^b + t^c + 1$. The particular pentanomial chosen has the following properties: the second term t^a has the lowest degree among all irreducible pentanomials of degree m ; the third term t^b has the lowest degree among all irreducible pentanomials of degree m and second term t^a ; and the fourth term t^c has the lowest degree among all irreducible pentanomials of degree m , second term t^a , and third term t^b .
- *Normal Basis:* Choose the Type T low-complexity normal basis with the smallest T .

For each binary field, the parameters are given for the above basis representations.

1.4 Choice of Curves

Two kinds of curves are given:

- *Pseudo-random* curves are those whose coefficients are generated from the output of a seeded cryptographic hash. If the seed value is given along with the coefficients, it can be verified easily that the coefficients were indeed generated by that method.
- *Special curves* whose coefficients and underlying field have been selected to optimize the efficiency of the elliptic curve operations.

For each size, the following curves are given:

- A pseudo-random curve over $GF(p)$.
- A pseudo-random curve over $GF(2^m)$.
- A special curve over $GF(2^m)$ called a *Koblitz curve* or *anomalous binary curve*.

The pseudo-random curves are generated via the SHA-1 based method given in the ANSI X9.62 and IEEE P1363 standards. (The generation and verification processes are given in Appendices 4 through 7.)

1.5 Choice of Base Points

Any point of order r can serve as the base point. Each curve is supplied with a sample base point $G = (G_x, G_y)$. Users may want to generate their own base points to ensure cryptographic separation of networks.

§2. CURVES OVER PRIME FIELDS

For each prime p , a pseudo-random curve

$$E : y^2 \equiv x^3 - 3x + b \pmod{p}$$

of prime order r is listed.¹ (Thus, for these curves, the cofactor is always $f = 1$.) The following parameters are given:

- The prime modulus p
- The order r
- The 160-bit input seed s to the SHA-1 based algorithm
- The output c of the SHA-1 based algorithm
- The coefficient b (satisfying $b^2 c \equiv -27 \pmod{p}$)
- The base point x coordinate G_x
- The base point y coordinate G_y

The integers p and r are given in decimal form; bit strings and field elements are given in hex.

¹The selection $a = -3$ for the coefficient of x was made for reasons of efficiency; see IEEE P1363.

Curve P-192

$p = 62771017353866807638357894232076664160839087\backslash$
00390324961279

$r = 62771017353866807638357894231760590137671947\backslash$
73182842284081

$s = 3045ae6f\ c8422f64\ ed579528\ d38120ea\ e12196d5$

$c =$ 3099d2bb
bfc b2538 542dcd5f b078b6ef 5f3d6fe2 c745de65

$b =$ 64210519
e59c80e7 0fa7e9ab 72243049 feb8deec c146b9b1

$G_x =$ 188da80e
b03090f6 7cbf20eb 43a18800 f4ff0afd 82ff1012

$G_y =$ 07192b95
ffc8da78 631011ed 6b24cdd5 73f977a1 1e794811

Curve P-224

$p = 26959946667150639794667015087019630673557916\backslash$
260026308143510066298881

$r = 26959946667150639794667015087019625940457807\backslash$
714424391721682722368061

$s = \text{bd713447 99d5c7fc dc45b59f a3b9ab8f 6a948bc5}$

$c =$ 5b056c7e 11dd68f4
0469ee7f 3c7a7d74 f7d12111 6506d031 218291fb

$b =$ b4050a85 0c04b3ab
f5413256 5044b0b7 d7bfd8ba 270b3943 2355ffb4

$G_x =$ b70e0cbd 6bb4bf7f
321390b9 4a03c1d3 56c21122 343280d6 115c1d21

$G_y =$ bd376388 b5f723fb
4c22dfe6 cd4375a0 5a074764 44d58199 85007e34

Curve P-256

$p = 11579208921035624876269744694940757353008614\backslash$
3415290314195533631308867097853951

$r = 11579208921035624876269744694940757352999695\backslash$
5224135760342422259061068512044369

$s = c49d3608 86e70493 6a6678e1 139d26b7 819f7e90$

$c =$ 7efba166 2985be94 03cb055c
75d4f7e0 ce8d84a9 c5114abc af317768 0104fa0d

$b =$ 5ac635d8 aa3a93e7 b3ebbd55
769886bc 651d06b0 cc53b0f6 3bce3c3e 27d2604b

$G_x =$ 6b17d1f2 e12c4247 f8bce6e5
63a440f2 77037d81 2deb33a0 f4a13945 d898c296

$G_y =$ 4fe342e2 fe1a7f9b 8ee7eb4a
7c0f9e16 2bce3357 6b315ece cbb64068 37bf51f5

Curve P-384

$p = 39402006196394479212279040100143613805079739\backslash$
 $27046544666794829340424572177149687032904726\backslash$
 $6088258938001861606973112319$

$r = 39402006196394479212279040100143613805079739\backslash$
 $27046544666794690527962765939911326356939895\backslash$
 $6308152294913554433653942643$

$s = a335926a a319a27a 1d00896a 6773a482 7acdac73$

$c =$ $79d1e655 f868f02f$
 $ff48dcde e14151dd b80643c1 406d0ca1 0dfe6fc5$
 $2009540a 495e8042 ea5f744f 6e184667 cc722483$

$b =$ $b3312fa7 e23ee7e4$
 $988e056b e3f82d19 181d9c6e fe814112 0314088f$
 $5013875a c656398d 8a2ed19d 2a85c8ed d3ec2aef$

$G_x =$ $aa87ca22 be8b0537$
 $8eb1c71e f320ad74 6e1d3b62 8ba79b98 59f741e0$
 $82542a38 5502f25d bf55296c 3a545e38 72760ab7$

$G_y =$ $3617de4a 96262c6f$
 $5d9e98bf 9292dc29 f8f41dbd 289a147c e9da3113$
 $b5f0b8c0 0a60b1ce 1d7e819d 7a431d7c 90ea0e5f$

Curve P-521

$p =$ 68647976601306097149819007990813932172694353\
00143305409394463459185543183397656052122559\
64066145455497729631139148085803712198799971\
6643812574028291115057151

$r =$ 68647976601306097149819007990813932172694353\
00143305409394463459185543183397655394245057\
74633321719753296399637136332111386476861244\
0380340372808892707005449

$s =$ d09e8800 291cb853 96cc6717 393284aa a0da64ba

$c =$ 0b4 8bfa5f42
0a349495 39d2bdfc 264eeeeb 077688e4 4fbf0ad8
f6d0edb3 7bd6b533 28100051 8e19f1b9 ffbe0fe9
ed8a3c22 00b8f875 e523868c 70c1e5bf 55bad637

$b =$ 051 953eb961
8e1c9a1f 929a21a0 b68540ee a2da725b 99b315f3
b8b48991 8ef109e1 56193951 ec7e937b 1652c0bd
3bb1bf07 3573df88 3d2c34f1 ef451fd4 6b503f00

$G_x =$ c6 858e06b7
0404e9cd 9e3ecb66 2395b442 9c648139 053fb521
f828af60 6b4d3dba a14b5e77 efe75928 fe1dc127
a2ffa8de 3348b3c1 856a429b f97e7e31 c2e5bd66

$G_y =$ 118 39296a78
9a3bc004 5c8a5fb4 2c7d1bd9 98f54449 579b4468
17afbd17 273e662c 97ee7299 5ef42640 c550b901
3fad0761 353c7086 a272c240 88be9476 9fd16650

§3. CURVES OVER BINARY FIELDS

For each field degree m , a pseudo-random curve is given, along with a Koblitz curve. The pseudo-random curve has the form

$$E : y^2 + x y = x^3 + x^2 + b,$$

and the Koblitz curve has the form

$$E_a : y^2 + x y = x^3 + a x^2 + 1$$

where $a = 0$ or 1 .

For each pseudo-random curve, the cofactor is $f = 2$. The cofactor of each Koblitz curve is $f = 2$ if $a = 1$ and $f = 4$ if $a = 0$.

The coefficients of the pseudo-random curves, and the coordinates of the base points of both kinds of curves, are given in terms of both the polynomial and normal basis representations discussed in §1.3.

For each m , the following parameters are given:

Field Representation:

- The normal basis type T
- The field polynomial (trinomial or pentanomial)

Koblitz Curve:

- The coefficient a
- The base point order r
- The base point x coordinate G_x
- The base point y coordinate G_y

Pseudo-random curve:

- The base point order r

Pseudo-random curve (Polynomial Basis representation):

- The coefficient b
- The base point x coordinate G_x
- The base point y coordinate G_y

Pseudo-random curve (Normal Basis representation):

- The 160-bit input seed s to the SHA-1 based algorithm
- The coefficient b (*i.e.*, the output of the SHA-1 based algorithm)
- The base point x coordinate G_x
- The base point y coordinate G_y

Integers (such as T , m , and r) are given in decimal form; bit strings and field elements are given in hex.

Degree 163 Binary Field

$$T = 4$$

$$p(t) = t^{163} + t^7 + t^6 + t^3 + 1$$

Curve K-163

$$a = 1$$

$$r = 5846006549323611672814741753598448348329118574063$$

Polynomial Basis:

$$G_x = \quad 2 \text{ fe13c053 7bbc11ac aa07d793 de4e6d5e 5c94eee8}$$

$$G_y = \quad 2 \text{ 89070fb0 5d38ff58 321f2e80 0536d538 ccdaa3d9}$$

Normal Basis:

$$G_x = \quad 0 \text{ 5679b353 caa46825 fea2d371 3ba450da 0c2a4541}$$

$$G_y = \quad 2 \text{ 35b7c671 00506899 06bac3d9 dec76a83 5591edb2}$$

Curve B-163

$$r = 5846006549323611672814742442876390689256843201587$$

Polynomial Basis:

$$b = \quad 2 \text{ 0a601907 b8c953ca 1481eb10 512f7874 4a3205fd}$$

$$G_x = \quad 3 \text{ f0eba162 86a2d57e a0991168 d4994637 e8343e36}$$

$$G_y = \quad 0 \text{ d51fbc6c 71a0094f a2cdd545 b11c5c0c 797324f1}$$

Normal Basis:

$s =$ 85e25bfe 5c86226c db12016f 7553f9d0 e693a268
 $b =$ 6 645f3cac f1638e13 9c6cd13e f61734fb c9e3d9fb
 $G_x =$ 0 311103c1 7167564a ce77ccb0 9c681f88 6ba54ee8
 $G_y =$ 3 33ac13c6 447f2e67 613bf700 9daf98c8 7bb50c7f

Degree 233 Binary Field

$T = 2$
 $p(t) = t^{233} + t^{74} + 1$

Curve K-233

$a = 0$
 $r = 34508731733952818937173779311385127605709409888622521\backslash$
26328087024741343

Polynomial Basis:

$G_x =$ 172 32ba853a 7e731af1
29f22ff4 149563a4 19c26bf5 0a4c9d6e efad6126
 $G_y =$ 1db 537dece8 19b7f70f
555a67c4 27a8cd9b f18aeb9b 56e0c110 56fae6a3

Normal Basis:

$G_x =$ 0fd e76d9dcd 26e643ac
26f1aa90 1aa12978 4b71fc07 22b2d056 14d650b3
 $G_y =$ 064 3e317633 155c9e04
47ba8020 a3c43177 450ee036 d6335014 34cac978

Curve B-233

$r = 69017463467905637874347558622770255558398127373450135 \setminus$
55379383634485463

Polynomial Basis:

$b =$ 066 647ede6c 332c7f8c
0923bb58 213b333b 20e9ce42 81fe115f 7d8f90ad
 $G_x =$ 0fa c9dfcbac 8313bb21
39f1bb75 5fef65bc 391f8b36 f8f8eb73 71fd558b
 $G_y =$ 100 6a08a419 03350678
e58528be bf8a0bef f867a7ca 36716f7e 01f81052

Normal Basis:

$s =$ 74d59ff0 7f6b413d 0ea14b34 4b20a2db 049b50c3
 $b =$ 1a0 03e0962d 4f9a8e40
7c904a95 38163adb 82521260 0c7752ad 52233279
 $G_x =$ 18b 863524b3 cdfefb94
f2784e0b 116faac5 4404bc91 62a363ba b84a14c5
 $G_y =$ 049 25df77bd 8b8ff1a5
ff519417 822bfedf 2bbd7526 44292c98 c7af6e02

Degree 283 Binary Field

$$T = 6$$

$$p(t) = t^{283} + t^{12} + t^7 + t^5 + 1$$

Curve K-283

$$a = 0$$

$$r = 38853377844514581418389238136470378132848117337930613 \backslash \\ 24295874997529815829704422603873$$

Polynomial Basis:

$$G_x = \begin{array}{cccc} 503213f & 78ca4488 & 3f1a3b81 & 62f188e5 \\ 53cd265f & 23c1567a & 16876913 & b0c2ac24 & 58492836 \end{array}$$

$$G_y = \begin{array}{cccc} 1ccda38 & 0f1c9e31 & 8d90f95d & 07e5426f \\ e87e45c0 & e8184698 & e4596236 & 4e341161 & 77dd2259 \end{array}$$

Normal Basis:

$$G_x = \begin{array}{cccc} 3ab9593 & f8db09fc & 188f1d7c & 4ac9fcc3 \\ e57fcd3b & db15024b & 212c7022 & 9de5fcd9 & 2eb0ea60 \end{array}$$

$$G_y = \begin{array}{cccc} 2118c47 & 55e7345c & d8f603ef & 93b98b10 \\ 6fe8854f & feb9a3b3 & 04634cc8 & 3a0e759f & 0c2686b1 \end{array}$$

Curve B-283

$r = 77706755689029162836778476272940756265696259243769048 \backslash$
89109196526770044277787378692871

Polynomial Basis:

$b =$ 27b680a c8b8596d a5a4af8a 19a0303f
ca97fd76 45309fa2 a581485a f6263e31 3b79a2f5
 $G_x =$ 5f93925 8db7dd90 e1934f8c 70b0dfec
2eed25b8 557eac9c 80e2e198 f8cdbecd 86b12053
 $G_y =$ 3676854 fe24141c b98fe6d4 b20d02b4
516ff702 350eddb0 826779c8 13f0df45 be8112f4

Normal Basis:

$s = 77e2b073 70eb0f83 2a6dd5b6 2dfc88cd 06bb84be$
 $b =$ 157261b 894739fb 5a13503f 55f0b3f1
0c560116 66331022 01138cc1 80c0206b dafbc951
 $G_x =$ 749468e 464ee468 634b21f7 f61cb700
701817e6 bc36a236 4cb8906e 940948ea a463c35d
 $G_y =$ 62968bd 3b489ac5 c9b859da 68475c31
5bafcdc4 ccd0dc90 5b70f624 46f49c05 2f49c08c

Degree 409 Binary Field

$$T = 4$$

$$p(t) = t^{409} + t^{87} + 1$$

Curve K-409

$$a = 0$$

$$r = 33052798439512429947595765401638551991420234148214060\backslash$$
$$96423243950228807112892491910506732584577774580140963\backslash$$
$$66590617731358671$$

Polynomial Basis:

$$G_x = \quad \quad \quad 060f05f \ 658f49c1 \ ad3ab189$$
$$0f718421 \ 0efd0987 \ e307c84c \ 27accfb8 \ f9f67cc2$$
$$c460189e \ b5aaaa62 \ ee222eb1 \ b35540cf \ e9023746$$

$$G_y = \quad \quad \quad 1e36905 \ 0b7c4e42 \ acba1dac$$
$$bf04299c \ 3460782f \ 918ea427 \ e6325165 \ e9ea10e3$$
$$da5f6c42 \ e9c55215 \ aa9ca27a \ 5863ec48 \ d8e0286b$$

Normal Basis:

$$G_x = \quad \quad \quad 1b559c7 \ cba2422e \ 3affe133$$
$$43e808b5 \ 5e012d72 \ 6ca0b7e6 \ a63aeafb \ c1e3a98e$$
$$10ca0fcf \ 98350c3b \ 7f89a975 \ 4a8e1dc0 \ 713cec4a$$

$$G_y = \quad \quad \quad 16d8c42 \ 052f07e7 \ 713e7490$$
$$eff318ba \ 1abd6fef \ 8a5433c8 \ 94b24f5c \ 817aeb79$$
$$852496fb \ ee803a47 \ bc8a2038 \ 78ebf1c4 \ 99afd7d6$$

Degree 571 Binary Field

$$T = 10$$

$$p(t) = t^{571} + t^{10} + t^5 + t^2 + 1$$

Curve K-571

$$a = 0$$

$$r = 19322687615086291723476759454659936721494636648532174\backslash$$
$$99328617625725759571144780212268133978522706711834706\backslash$$
$$71280082535146127367497406661731192968242161709250355\backslash$$
$$5733685276673$$

Polynomial Basis:

$$G_x = \begin{array}{cccccc} 26eb7a8 & 59923fbc & 82189631 & & & \\ f8103fe4 & ac9ca297 & 0012d5d4 & 60248048 & 01841ca4 & \\ 43709584 & 93b205e6 & 47da304d & b4ceb08c & bbd1ba39 & \\ 494776fb & 988b4717 & 4dca88c7 & e2945283 & a01c8972 & \end{array}$$

$$G_y = \begin{array}{cccccc} 349dc80 & 7f4fbf37 & 4f4aeade & & & \\ 3bca9531 & 4dd58cec & 9f307a54 & ffc61efc & 006d8a2c & \\ 9d4979c0 & ac44aea7 & 4fbеbb9 & f772aedc & b620b01a & \\ 7ba7af1b & 320430c8 & 591984f6 & 01cd4c14 & 3ef1c7a3 & \end{array}$$

Normal Basis:

$$\begin{aligned} G_x = & \quad 04bb2db \ a418d0db \ 107adae0 \\ & 03427e5d \ 7cc139ac \ b465e593 \ 4f0bea2a \ b2f3622b \\ & c29b3d5b \ 9aa7a1fd \ fd5d8be6 \ 6057c100 \ 8e71e484 \\ & bcd98f22 \ bf847642 \ 37673674 \ 29ef2ec5 \ bc3ebcf7 \\ G_y = & \quad 44cbb57 \ de20788d \ 2c952d7b \\ & 56cf39bd \ 3e89b189 \ 84bd124e \ 751ceff4 \ 369dd8da \\ & c6a59e6e \ 745df44d \ 8220ce22 \ aa2c852c \ fcbbef49 \\ & ebaa98bd \ 2483e331 \ 80e04286 \ feaa2530 \ 50caff60 \end{aligned}$$

Curve B-571

$$\begin{aligned} r = & \ 38645375230172583446953518909319873442989273297064349 \backslash \\ & 98657235251451519142289560424536143999389415773083133 \backslash \\ & 88112192694448624687246281681307023452828830333241139 \backslash \\ & 3191105285703 \end{aligned}$$

Polynomial Basis:

$$\begin{aligned} b = & \quad 2f40e7e \ 2221f295 \ de297117 \\ & b7f3d62f \ 5c6a97ff \ cb8ceff1 \ cd6ba8ce \ 4a9a18ad \\ & 84ffabbd \ 8efa5933 \ 2be7ad67 \ 56a66e29 \ 4afd185a \\ & 78ff12aa \ 520e4de7 \ 39baca0c \ 7ffeff7f \ 2955727a \\ G_x = & \quad 303001d \ 34b85629 \ 6c16c0d4 \\ & 0d3cd775 \ 0a93d1d2 \ 955fa80a \ a5f40fc8 \ db7b2abd \\ & bde53950 \ f4c0d293 \ cdd711a3 \ 5b67fb14 \ 99ae6003 \\ & 8614f139 \ 4abfa3b4 \ c850d927 \ e1e7769c \ 8eec2d19 \end{aligned}$$

$G_y =$ 37bf273 42da639b 6dccfffe
b73d69d7 8c6c27a6 009cbbca 1980f853 3921e8a6
84423e43 bab08a57 6291af8f 461bb2a8 b3531d2f
0485c19b 16e2f151 6e23dd3c 1a4827af 1b8ac15b

Normal Basis:

$s =$ 2aa058f7 3a0e33ab 486b0f61 0410c53a 7f132310

$b =$ 3762d0d 47116006 179da356

88eeaccf 591a5cde a7500011 8d9608c5 9132d434

26101a1d fb377411 5f586623 f75f0000 1ce61198

3c1275fa 31f5bc9f 4be1a0f4 67f01ca8 85c74777

$G_x =$ 0735e03 5def5925 cc33173e

b2a8ce77 67522b46 6d278b65 0a291612 7dfea9d2

d361089f 0a7a0247 a184e1c7 0d417866 e0fe0feb

0ff8f2f3 f9176418 f97d117e 624e2015 df1662a8

$G_y =$ 04a3642 0572616c df7e606f

ccadaecf c3b76dab 0eb1248d d03fbdfc 9cd3242c

4726be57 9855e812 de7ec5c5 00b4576a 24628048

b6a72d88 0062eed0 dd34b109 6d3acbb6 b01a4a97

APPENDIX 1: IMPLEMENTATION OF MODULAR ARITHMETIC

The prime moduli in the above examples are of a special type (called *generalized Mersenne numbers*) for which modular multiplication can be carried out more efficiently than in general. This appendix provides the rules for implementing this faster arithmetic, for each of the prime moduli appearing in the examples.

The usual way to multiply two integers (mod m) is to take the integer product and reduce it (mod m). One therefore has the following problem: given an integer A less than m^2 , compute

$$B := A \bmod m.$$

In general, one must obtain B as the remainder of an integer division. If m is a generalized Mersenne number, however, then B can be expressed as a sum or difference (mod m) of a small number of terms. To compute this expression, one can evaluate the integer sum or difference and reduce the result modulo m . The latter reduction can be accomplished by adding or subtracting a few copies of m .

The prime moduli p for each of the five example curves is a generalized Mersenne number.

Curve P-192:

The modulus for this curve is $p = 2^{192} - 2^{64} - 1$. Every integer A less than p^2 can be written

$$A = A_5 \cdot 2^{320} + A_4 \cdot 2^{256} + A_3 \cdot 2^{192} + A_2 \cdot 2^{128} + A_1 \cdot 2^{64} + A_0,$$

where each A_i is a 64-bit integer. The expression for B is

$$B := T + S_1 + S_2 + S_3 \bmod p,$$

where the 192-bit terms are given by

$$T = A_2 \cdot 2^{128} + A_1 \cdot 2^{64} + A_0$$

$$S_1 = A_3 \cdot 2^{64} + A_3$$

$$S_2 = A_4 \cdot 2^{128} + A_4 \cdot 2^{64}$$

$$S_3 = A_5 \cdot 2^{128} + A_5 \cdot 2^{64} + A_5.$$

Curve P-224:

The modulus for this curve is $p = 2^{224} - 2^{96} + 1$. Every integer A less than p^2 can be written

$$\begin{aligned} A = & A_{13} \cdot 2^{416} + A_{12} \cdot 2^{384} + A_{11} \cdot 2^{352} + A_{10} \cdot 2^{320} + \\ & A_9 \cdot 2^{288} + A_8 \cdot 2^{256} + A_7 \cdot 2^{224} + A_6 \cdot 2^{192} + A_5 \cdot 2^{160} + \\ & A_4 \cdot 2^{128} + A_3 \cdot 2^{96} + A_2 \cdot 2^{64} + A_1 \cdot 2^{32} + A_0, \end{aligned}$$

where each A_i is a 32-bit integer. As a concatenation of 32-bit words, this can be denoted by

$$A = (A_{13} \parallel A_{12} \parallel \cdots \parallel A_0).$$

The expression for B is

$$B := T + S_1 + S_2 - D_1 - D_2 \bmod p,$$

where the 224-bit terms are given by

$$\begin{aligned} T &= (A_6 \parallel A_5 \parallel A_4 \parallel A_3 \parallel A_2 \parallel A_1 \parallel A_0) \\ S_1 &= (A_{10} \parallel A_9 \parallel A_8 \parallel A_7 \parallel 0 \parallel 0 \parallel 0) \\ S_2 &= (0 \parallel A_{13} \parallel A_{12} \parallel A_{11} \parallel 0 \parallel 0 \parallel 0) \\ D_1 &= (A_{13} \parallel A_{12} \parallel A_{11} \parallel A_{10} \parallel A_9 \parallel A_8 \parallel A_7) \\ D_2 &= (0 \parallel 0 \parallel 0 \parallel 0 \parallel A_{13} \parallel A_{12} \parallel A_{11}). \end{aligned}$$

Curve P-256:

The modulus for this curve is $p = 2^{256} - 2^{224} + 2^{192} + 2^{96} - 1$. Every integer A less than p^2 can be written

$$\begin{aligned} A = & A_{15} \cdot 2^{480} + A_{14} \cdot 2^{448} + A_{13} \cdot 2^{416} + A_{12} \cdot 2^{384} + A_{11} \cdot 2^{352} + \\ & A_{10} \cdot 2^{320} + A_9 \cdot 2^{288} + A_8 \cdot 2^{256} + A_7 \cdot 2^{224} + A_6 \cdot 2^{192} + \\ & A_5 \cdot 2^{160} + A_4 \cdot 2^{128} + A_3 \cdot 2^{96} + A_2 \cdot 2^{64} + A_1 \cdot 2^{32} + A_0, \end{aligned}$$

where each A_i is a 32-bit integer. As a concatenation of 32-bit words, this can be denoted by

$$A = (A_{15} \parallel A_{14} \parallel \cdots \parallel A_0).$$

The expression for B is

$$B := T + 2S_1 + 2S_2 + S_3 + S_4 - D_1 - D_2 - D_3 - D_4 \pmod{p},$$

where the 256-bit terms are given by

$$\begin{aligned} T &= (A_7 \parallel A_6 \parallel A_5 \parallel A_4 \parallel A_3 \parallel A_2 \parallel A_1 \parallel A_0) \\ S_1 &= (A_{15} \parallel A_{14} \parallel A_{13} \parallel A_{12} \parallel A_{11} \parallel 0 \parallel 0 \parallel 0) \\ S_2 &= (0 \parallel A_{15} \parallel A_{14} \parallel A_{13} \parallel A_{12} \parallel 0 \parallel 0 \parallel 0) \\ S_3 &= (A_{15} \parallel A_{14} \parallel 0 \parallel 0 \parallel 0 \parallel A_{10} \parallel A_9 \parallel A_8) \\ S_4 &= (A_8 \parallel A_{13} \parallel A_{15} \parallel A_{14} \parallel A_{13} \parallel A_{11} \parallel A_{10} \parallel A_9) \\ D_1 &= (A_{10} \parallel A_8 \parallel 0 \parallel 0 \parallel 0 \parallel A_{13} \parallel A_{12} \parallel A_{11}) \\ D_2 &= (A_{11} \parallel A_9 \parallel 0 \parallel 0 \parallel A_{15} \parallel A_{14} \parallel A_{13} \parallel A_{12}) \\ D_3 &= (A_{12} \parallel 0 \parallel A_{10} \parallel A_9 \parallel A_8 \parallel A_{15} \parallel A_{14} \parallel A_{13}) \\ D_4 &= (A_{13} \parallel 0 \parallel A_{11} \parallel A_{10} \parallel A_9 \parallel 0 \parallel A_{15} \parallel A_{14}). \end{aligned}$$

Curve P-384:

The modulus for this curve is $p = 2^{384} - 2^{128} - 2^{96} + 2^{32} - 1$. Every integer A less than p^2 can be written

$$\begin{aligned} A = & A_{23} \cdot 2^{736} + A_{22} \cdot 2^{704} + A_{21} \cdot 2^{672} + A_{20} \cdot 2^{640} + A_{19} \cdot 2^{608} + \\ & A_{18} \cdot 2^{576} + A_{17} \cdot 2^{544} + A_{16} \cdot 2^{512} + A_{15} \cdot 2^{480} + A_{14} \cdot 2^{448} + \\ & A_{13} \cdot 2^{416} + A_{12} \cdot 2^{384} + A_{11} \cdot 2^{352} + A_{10} \cdot 2^{320} + A_9 \cdot 2^{288} + \\ & A_8 \cdot 2^{256} + A_7 \cdot 2^{224} + A_6 \cdot 2^{192} + A_5 \cdot 2^{160} + A_4 \cdot 2^{128} + \\ & A_3 \cdot 2^{96} + A_2 \cdot 2^{64} + A_1 \cdot 2^{32} + A_0, \end{aligned}$$

where each A_i is a 32-bit integer. As a concatenation of 32-bit words, this can be denoted by

$$A = (A_{23} \parallel A_{22} \parallel \cdots \parallel A_0).$$

The expression for B is

$$B := T + 2S_1 + S_2 + S_3 + S_4 + S_5 + S_6 - D_1 - D_2 - D_3 \pmod{p},$$

where the 384-bit terms are given by

$$\begin{aligned} T &= (A_{11} \parallel A_{10} \parallel A_9 \parallel A_8 \parallel A_7 \parallel A_6 \parallel A_5 \parallel A_4 \parallel A_3 \parallel A_2 \parallel A_1 \parallel A_0) \\ S_1 &= (0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel A_{23} \parallel A_{22} \parallel A_{21} \parallel 0 \parallel 0 \parallel 0 \parallel 0) \\ S_2 &= (A_{23} \parallel A_{22} \parallel A_{21} \parallel A_{20} \parallel A_{19} \parallel A_{18} \parallel A_{17} \parallel A_{16} \parallel A_{15} \parallel A_{14} \parallel A_{13} \parallel A_{12}) \\ S_3 &= (A_{20} \parallel A_{19} \parallel A_{18} \parallel A_{17} \parallel A_{16} \parallel A_{15} \parallel A_{14} \parallel A_{13} \parallel A_{12} \parallel A_{23} \parallel A_{22} \parallel A_{21}) \\ S_4 &= (A_{19} \parallel A_{18} \parallel A_{17} \parallel A_{16} \parallel A_{15} \parallel A_{14} \parallel A_{13} \parallel A_{12} \parallel A_{20} \parallel 0 \parallel A_{23} \parallel 0) \\ S_5 &= (0 \parallel 0 \parallel 0 \parallel 0 \parallel A_{23} \parallel A_{22} \parallel A_{21} \parallel A_{20} \parallel 0 \parallel 0 \parallel 0 \parallel 0) \\ S_6 &= (0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel A_{23} \parallel A_{22} \parallel A_{21} \parallel 0 \parallel 0 \parallel A_{20}) \\ D_1 &= (A_{22} \parallel A_{21} \parallel A_{20} \parallel A_{19} \parallel A_{18} \parallel A_{17} \parallel A_{16} \parallel A_{15} \parallel A_{14} \parallel A_{13} \parallel A_{12} \parallel A_{23}) \\ D_2 &= (0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel A_{23} \parallel A_{22} \parallel A_{21} \parallel A_{20} \parallel 0) \\ D_3 &= (0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel 0 \parallel A_{23} \parallel A_{23} \parallel 0 \parallel 0 \parallel 0). \end{aligned}$$

Curve P-521:

The modulus for this curve is $p = 2^{521} - 1$. Every integer A less than p^2 can be written

$$A = A_1 \cdot 2^{521} + A_0,$$

The expression for B is

$$B := A_0 + A_1 \bmod p.$$

APPENDIX 2: NORMAL BASES

The elements of $GF(2^m)$ are expressed in terms of the a type T normal basis² \mathcal{B} for $GF(2^m)$, for some T . Each element has a unique representation as a bit string

$$(a_0 a_1 \dots a_{m-1}).$$

The arithmetic operations are performed as follows.

Addition: addition of two elements is implemented by bitwise addition modulo 2. Thus, for example,

$$(1100111) + (1010010) = (0110101).$$

Squaring: if

$$\alpha = (a_0 a_1 \dots a_{m-1}),$$

then

$$\alpha^2 = (a_{m-1} a_0 a_1 \dots a_{m-2}).$$

Multiplication: to perform multiplication, one first constructs a function $F(\underline{u}, \underline{v})$ on inputs

$$\underline{u} = (u_0 u_1 \dots u_{m-1}) \quad \text{and} \quad \underline{v} = (v_0 v_1 \dots v_{m-1})$$

as follows.

1. Set $p \leftarrow Tm + 1$
2. Let u be an integer having order T modulo p

²It is assumed in this section that m is odd and T even, since this is the only case considered in this standard.

3. Compute the sequence $F(1), F(2), \dots, F(p-1)$ as follows:

3.1 Set $w \leftarrow 1$

3.2 For j from 0 to $T-1$ do

Set $n \leftarrow w$

For i from 0 to $m-1$ do

Set $F(n) \leftarrow i$

Set $n \leftarrow 2n \bmod p$

Set $w \leftarrow uw \bmod p$

4. Output the formula

$$F(\underline{u}, \underline{v}) := \sum_{k=1}^{p-2} u_{F(k+1)} v_{F(p-k)}.$$

This computation need only be performed once per basis.

Given the function F for \mathcal{B} , one computes the product

$$(c_0 c_1 \dots c_{m-1}) = (a_0 a_1 \dots a_{m-1}) \times (b_0 b_1 \dots b_{m-1})$$

as follows.

1. Set $(u_0 u_1 \dots u_{m-1}) \leftarrow (a_0 a_1 \dots a_{m-1})$

2. Set $(v_0 v_1 \dots v_{m-1}) \leftarrow (b_0 b_1 \dots b_{m-1})$

3. For k from 0 to $m-1$ do

3.1 Compute

$$c_k := F(\underline{u}, \underline{v})$$

3.2 Set $u \leftarrow \text{LeftShift}(u)$ and $v \leftarrow \text{LeftShift}(v)$, where **LeftShift** denotes the circular left shift operation.

4. Output $c := (c_0 c_1 \dots c_{m-1})$

EXAMPLE. For the type 4 normal basis for $GF(2^7)$, one has $p = 29$ and $u = 12$ or 17 . Thus the values of F are given by

$F(1) = 0$	$F(8) = 3$	$F(15) = 6$	$F(22) = 5$
$F(2) = 1$	$F(9) = 3$	$F(16) = 4$	$F(23) = 6$
$F(3) = 5$	$F(10) = 2$	$F(17) = 0$	$F(24) = 1$
$F(4) = 2$	$F(11) = 4$	$F(18) = 4$	$F(25) = 2$
$F(5) = 1$	$F(12) = 0$	$F(19) = 2$	$F(26) = 5$
$F(6) = 6$	$F(13) = 4$	$F(20) = 3$	$F(27) = 1$
$F(7) = 5$	$F(14) = 6$	$F(21) = 3$	$F(28) = 0$

Therefore

$$\begin{aligned}
 F(\underline{u}, \underline{v}) &= u_0 v_1 + u_1 (v_0 + v_2 + v_5 + v_6) + u_2 (v_1 + v_3 + v_4 + v_5) \\
 &\quad + u_3 (v_2 + v_5) + u_4 (v_2 + v_6) + u_5 (v_1 + v_2 + v_3 + v_6) \\
 &\quad + u_6 (v_1 + v_4 + v_5 + v_6).
 \end{aligned}$$

Thus, if

$$a = (1010111) \quad \text{and} \quad b = (1100001),$$

then

$$\begin{aligned}
 c_0 &= F((1010111), (1100001)) = 1, \\
 c_1 &= F((0101111), (1000011)) = 0, \\
 &\quad \vdots \\
 c_6 &= F((1101011), (1110000)) = 1,
 \end{aligned}$$

so that $c = ab = (1011001)$.

APPENDIX 3: SCALAR MULTIPLICATION ON KOBLITZ CURVES

This appendix describes a particularly efficient method of computing the scalar multiple nP on the Koblitz curve E_a over $GF(2^m)$.

The operation τ is defined by

$$\tau(x, y) = (x^2, y^2).$$

When the normal basis representation is used, then the operation τ is implemented by performing right circular shifts on the bit strings representing x and y .

Given m and a , define the following parameters:

- C is some integer greater than 5.
- $\mu := (-1)^{1-a}$
- For $i = 0$ and $i = 1$, define the sequence $s_i(m)$ by

$$s_i(0) = 0, \quad s_i(1) = 1 - i,$$

$$s_i(m) = \mu \cdot s_i(m - 1) - 2 s_i(m - 2) + (-1)^i.$$

- Define the sequence $V(m)$ by

$$V(0) = 2, \quad V(1) = \mu,$$

$$V(m) = \mu \cdot V(m - 1) - 2 V(m - 2).$$

For the example curves, the quantities $s_i(m)$ and $V(m)$ are as follows.

Curve K-163:

$$\begin{aligned}s_0(163) &= 2579386439110731650419537 \\s_1(163) &= -755360064476226375461594 \\V(163) &= -4845466632539410776804317\end{aligned}$$

Curve K-233:

$$\begin{aligned}s_0(233) &= -27859711741434429761757834964435883 \\s_1(233) &= -44192136247082304936052160908934886 \\V(233) &= -137381546011108235394987299651366779\end{aligned}$$

Curve K-283:

$$\begin{aligned}s_0(283) &= -665981532109049041108795536001591469280025 \\s_1(283) &= 1155860054909136775192281072591609913945968 \\V(283) &= 7777244870872830999287791970962823977569917\end{aligned}$$

Curve K-409:

$$\begin{aligned}s_0(409) &= -1830751045600238213781031719875646137859054248755686\backslash \\ &9338419259 \\s_1(409) &= -8893048526138304097196653241844212679626566100996606\backslash \\ &444816790 \\V(409) &= 1045728873731562592744768538704832073763879695768757\backslash \\ &5791173829\end{aligned}$$

Curve K-571:

$$s_0(571) = -373731944687646369242938589247611556714729396459613\backslash$$

1024123406420235241916729983261305

$s_1(571) = -319185770644641609953814595948959674131968912148564 \backslash$
65861056511758982848515832612248752

$V(571) = -148380926981691413899619140297051490364542574180493 \backslash$
936232912339534208516828973111459843

The following algorithm computes the scalar multiple nP on the Koblitz curve E_a over $GF(2^m)$. The average number of elliptic additions and subtractions is at most $\sim 1 + (m/3)$, and is at most $\sim m/3$ with probability at least $1 - 2^{5-C}$.

```
For  $i = 0$  to  $1$  do
     $n' \leftarrow \lfloor n / 2^{a-C+(m-9)/2} \rfloor$ 
     $g' \leftarrow s_i(m) \cdot n'$ 
     $h' \leftarrow \lfloor g' / 2^m \rfloor$ 
     $j' \leftarrow V(m) \cdot h'$ 
     $\ell' \leftarrow \text{Round}((g' + j') / 2^{(m+5)/2})$ 
     $\lambda_i \leftarrow \ell' / 2^C$ 
     $f_i \leftarrow \text{Round}(\lambda_i)$ 
     $\eta_i \leftarrow \lambda_i - f_i$ 
     $h_i \leftarrow 0$ 
 $\eta \leftarrow 2\eta_0 + \mu\eta_1$ 
If  $\eta \geq 1$ 
    then
        if  $\eta_0 - 3\mu\eta_1 < -1$ 
            then set  $h_1 \leftarrow \mu$ 
            else set  $h_0 \leftarrow 1$ 
```

```

else
    if  $\eta_0 + 4\mu\eta_1 \geq 2$ 
        then set  $h_1 \leftarrow \mu$ 
If  $\eta < -1$ 
    then
        if  $\eta_0 - 3\mu\eta_1 \geq 1$ 
            then set  $h_1 \leftarrow -\mu$ 
            else set  $h_0 \leftarrow -1$ 
        else
            if  $\eta_0 + 4\mu\eta_1 < -2$ 
                then set  $h_1 \leftarrow -\mu$ 
 $q_0 \leftarrow f_0 + h_0$ 
 $q_1 \leftarrow f_1 + h_1$ 
 $r_0 \leftarrow n - (s_0 + \mu s_1)q_0 - 2s_1q_1$ 
 $r_1 \leftarrow s_1q_0 - s_0q_1$ 
Set  $Q \leftarrow \mathcal{O}$ 
 $P_0 \leftarrow P$ 
While  $r_0 \neq 0$  or  $r_1 \neq 0$ 
    If  $r_0$  odd then
        set  $u \leftarrow 2 - (r_0 - 2r_1 \bmod 4)$ 
        set  $r_0 \leftarrow r_0 - u$ 
        if  $u = 1$  then set  $Q \leftarrow Q + P_0$ 
        if  $u = -1$  then set  $Q \leftarrow Q - P_0$ 
    Set  $P_0 \leftarrow \tau P_0$ 
    Set  $(r_0, r_1) \leftarrow (r_1 + \mu r_0/2, -r_0/2)$ 
EndWhile
Output  $Q$ 

```

APPENDIX 4: GENERATION OF
PSEUDO-RANDOM CURVES (PRIME CASE)

Let ℓ be the bit length of p , and define

$$v = \lfloor (\ell - 1)/160 \rfloor$$

$$w = \ell - 160v - 1$$

1. Choose an arbitrary 160-bit string s .
2. Compute $h := \text{SHA-1}(s)$.
3. Let h_0 be the bit string obtained by taking the w rightmost bits of h .
4. Let z be the integer whose binary expansion is given by the 160-bit string s .
5. For i from 1 to v do:
 - 5.1 Define the 160-bit string s_i to be binary expansion of the integer $(z + i) \bmod (2^{160})$.
 - 5.2 Compute $h_i := \text{SHA-1}(s_i)$.
6. Let h be the bit string obtained by the concatenation of h_0, h_1, \dots, h_v as follows:

$$h = h_0 \| h_1 \| \dots \| h_v.$$

7. Let c be the integer whose binary expansion is given by the bit string h .
8. If $c = 0$ or $4c + 27 \equiv 0 \pmod{p}$, then go to Step 1.
9. Choose integers $a, b \in GF(p)$ such that

$$cb^2 \equiv a^3 \pmod{p}.$$

(The simplest choice is $a = c$ and $b = c$. However, one may want to choose differently for performance reasons.)

10. Check that the elliptic curve E over $GF(p)$ given by $y^2 = x^3 + ax + b$ has suitable order. If not, go to Step 1.

APPENDIX 5: VERIFICATION OF CURVE
PSEUDO-RANDOMNESS (PRIME CASE)

Given the 160-bit seed value s , one can verify that the coefficient b was obtained from s via the cryptographic hash function SHA-1 as follows. Let ℓ be the bit length of p , and define

$$v = \lfloor (\ell - 1)/160 \rfloor$$

$$w = \ell - 160v - 1$$

1. Compute $h := \text{SHA-1}(s)$.
2. Let h_0 be the bit string obtained by taking the w rightmost bits of h .
3. Let z be the integer whose binary expansion is given by the 160-bit string s .
4. For i from 1 to v do
 - 4.1 Define the 160-bit string s_i to be binary expansion of the integer $(z + i) \bmod (2^{160})$
 - 4.2 Compute $h_i := \text{SHA-1}(s_i)$.
5. Let h be the bit string obtained by the concatenation of h_0, h_1, \dots, h_v as follows:

$$h = h_0 \| h_1 \| \dots \| h_v.$$

6. Let c be the integer whose binary expansion is given by the bit string h .
7. Verify that $b^2 c \equiv -27 \pmod{p}$.

APPENDIX 6: GENERATION OF
PSEUDO-RANDOM CURVES (BINARY CASE)

Let:

$$v = \lfloor (m - 1)/B \rfloor$$

$$w = m - Bv$$

1. Choose an arbitrary 160-bit string s .
2. Compute $h := \text{SHA-1}(s)$.
3. Let h_0 be the bit string obtained by taking the w rightmost bits of h .
4. Let z be the integer whose binary expansion is given by the 160-bit string s .
5. For i from 1 to v do:
 - 5.1 Define the 160-bit string s_i to be binary expansion of the integer $(z + i) \bmod (2^{160})$.
 - 5.2 Compute $h_i := \text{SHA-1}(s_i)$.
6. Let h be the bit string obtained by the concatenation of h_0, h_1, \dots, h_v as follows:

$$h = h_0 \| h_1 \| \dots \| h_v.$$

7. Let b be the element of $GF(2^m)$ whose binary expansion is given by the bit string h .
8. Choose an element a of $GF(2^m)$.
9. Check that the elliptic curve E over $GF(2^m)$ given by $y^2 + xy = x^3 + ax^2 + b$ has suitable order. If not, go to Step 1.

APPENDIX 7: VERIFICATION OF CURVE
PSEUDO-RANDOMNESS (BINARY CASE)

Given the 160-bit seed value s , one can verify that the coefficient b was obtained from s via the cryptographic hash function SHA-1 as follows. Define

$$v = \lfloor (m - 1)/160 \rfloor$$

$$w = m - 160v$$

1. Compute $h := \text{SHA-1}(s)$.
2. Let h_0 be the bit string obtained by taking the w rightmost bits of h .
3. Let z be the integer whose binary expansion is given by the 160-bit string s .
4. For i from 1 to v do
 - 4.1 Define the 160-bit string s_i to be binary expansion of the integer $(z + i) \bmod (2^{160})$
 - 4.2 Compute $h_i := \text{SHA-1}(s_i)$.
5. Let h be the bit string obtained by the concatenation of h_0, h_1, \dots, h_v as follows:

$$h = h_0 \| h_1 \| \dots \| h_v.$$

6. Let c be the element of $GF(2^m)$ which is represented by the bit string h .
7. Verify that $c = b$.

APPENDIX 8: POLYNOMIAL BASIS TO NORMAL BASIS CONVERSION

Suppose that α an element of the field $GF(2^m)$. Denote by \mathbf{p} the bit string representing α with respect to a given polynomial basis. It is desired to compute \mathbf{n} , the bit string representing α with respect to a given normal basis. This is done via the matrix computation

$$\mathbf{p} \Gamma = \mathbf{n},$$

where Γ is an m -by- m matrix with entries in $GF(2)$. The matrix Γ , which depends only on the bases, can be computed easily given its second-to-last row. The second-to-last row for each conversion is given in the table below.

Degree 163:

3 e173bfaf 3a86434d 883a2918 a489ddb d 69fe84e1

Degree 233:

0be 19b89595 28bbc490
038f4bc4 da8bdfc1 ca36bb05 853fd0ed 0ae200ce

Degree 283:

3347f17 521fdabc 62ec1551 acf156fb
0bceb855 f174d4c1 7807511c 9f745382 add53bc3

Degree 409:

0eb00f2 ea95fd6c 64024e7f
0b68b81f 5ff8a467 acc2b4c3 b9372843 6265c7ff
a06d896c ae3a7e31 e295ec30 3eb9f769 de78bef5

Degree 571:

```
7940ffa ef996513 4d59dcbf
e5bf239b e4fe4b41 05959c5d 4d942ffd 46ea35f3
e3cdb0e1 04a2aa01 cef30a3a 49478011 196bfb43
c55091b6 1174d7c0 8d0cdd61 3bf6748a bad972a4
```

Given the second-to-last row \mathbf{r} of Γ , the rest of the matrix is computed as follows. Let β be the element of $GF(2^m)$ whose representation with respect to the normal basis is \mathbf{r} . Then the rows of Γ , from top to bottom, are the bit strings representing the elements

$$\beta^{m-1}, \beta^{m-2}, \dots, \beta^2, \beta, 1$$

with respect to the normal basis. (Note that the element 1 is represented by the all-1 bit string.)

Alternatively, the matrix is the inverse of the matrix described in Appendix 9.

More details of these computations can be found in Annex A.7 of the IEEE P1363 standard.

APPENDIX 9: NORMAL BASIS TO POLYNOMIAL BASIS CONVERSION

Suppose that α an element of the field $GF(2^m)$. Denote by \mathbf{n} the bit string representing α with respect to a given normal basis. It is desired to compute \mathbf{p} , the bit string representing α with respect to a given polynomial basis. This is done via the matrix computation

$$\mathbf{n}\Gamma = \mathbf{p},$$

where Γ is an m -by- m matrix with entries in $GF(2)$. The matrix Γ , which depends only on the bases, can be computed easily given its top row. The top row for each conversion is given in the table below.

Degree 163:

7 15169c10 9c612e39 0d347c74 8342bcd3 b02a0bef

Degree 233:

149 9e398ac5 d79e3685
59b35ca4 9bb7305d a6c0390b cf9e2300 253203c9

Degree 283:

31e0ed7 91c3282d c5624a72 0818049d
053e8c7a b8663792 bc1d792e ba9867fc 7b317a99

Degree 409:

0dfa06b e206aa97 b7a41fff
b9b0c55f 8f048062 fbe8381b 4248adf9 2912ccc8
e3f91a24 e1cfb395 0532b988 971c2304 2e85708d

Degree 571:

```
452186b bf5840a0 bcf8c9f0
2a54efa0 4e813b43 c3d41496 06c4d27b 487bf107
393c8907 f79d9778 beb35ee8 7467d328 8274caeb
da6ce05a eb4ca5cf 3c3044bd 4372232f 2c1a27c4
```

Given the top row \mathbf{r} of Γ , the rest of the matrix is computed as follows. Let β be the element of $GF(2^m)$ whose representation with respect to the polynomial basis is \mathbf{r} . Then the rows of Γ , from top to bottom, are the bit strings representing the elements

$$\beta, \beta^2, \beta^{2^2}, \dots, \beta^{2^{m-1}}$$

with respect to the polynomial basis.

Alternatively, the matrix is the inverse of the matrix described in Appendix 8.

More details of these computations can be found in Annex A.7 of the IEEE P1363 standard.