System-Level Virtualization & OSCAR-V

Presented by

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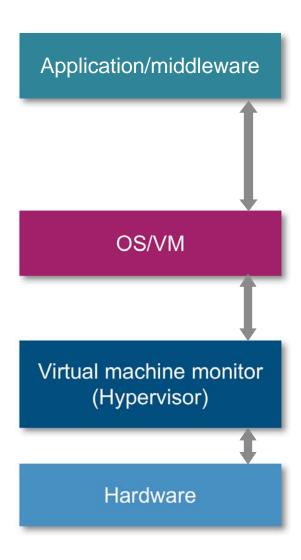
Virtualization technologies

Application/middleware

- Software component frameworks
 - Harness, Common Component Architecture
- Parallel programming languages and environments
 - PVM, MPI, Co-Array Fortran
- Serial programming languages and environments
 - C, POSIX (Processes, IPC, Threads)

OS/VM

- VMWare, Virtual PC, Virtual Server, and Qemu
- Hypervisor
 - Xen, Denali
- Hardware
 - OS Drivers, BIOS, Intel VT, AMD-V (Pacifica)





Emerging system-level virtualization

Hypervisors

- OS-level virtual machines (VMs)
- Paravirtualization for performance gain
 - Intercept and marshal privileged instructions issued by the guest machines
- Example: Xen + Linux
- HPC using virtualization
 - Example: Xen + Linux cluster + Infiniband (OSU/IBM)
 - Hypervisor (Host OS) bypass directly to IB



Why hypervisors in HPC?

Improved utilization

- Users with differing OS requirements can be easily satisfied, e.g.,
 Linux, Catamount, others in future
- Enable early access to petascale software environment on existing smaller systems
- Improved manageability
 - OS upgrades can be staged across VMs and thus minimize downtime
 - OS/RTE can be reconfigured and deployed on demand
- Improved reliability
 - Application-level software failures can be isolated to the VMs in which they occur
- Improved workload isolation, consolidation, and migration
 - Seamless transition between application development and deployment using petascale software environment on development systems
 - Proactive fault tolerance (preemptive migration) transparent to OS, runtime, and application

What about performance?

- Today hypervisors cost around 4–8% CPU time
- Improvements in hardware support by AMD and Intel will lessen this impact
- Proactive fault tolerance improves efficiency
 - Nonstop computing through preemptive measures
 - Significant reduction of checkpoint frequency
- Xen-like Catamount effort by Sandia/UNM to use Catamount as a HPC hypervisor



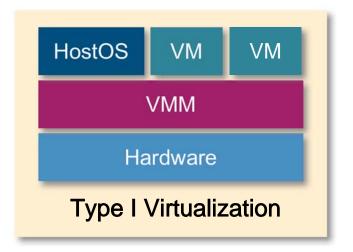
Virtual system environment

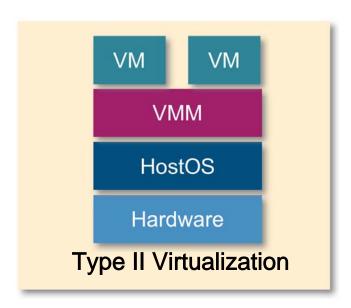
- Powerful abstraction concept that encapsulates OS, application runtime, and application
- Virtual parallel system instance running on a real HPC system using system-level virtualization
- Addressed key issues
 - Usability through virtual system management tools
 - Partitioning and reliability using adaptive runtime
 - Efficiency and reliability via proactive fault tolerance
 - Portability and efficiency through hypervisor + Linux/Catamount



System-level virtualization

- First research in the domain,
 Goldberg 73
 - Type-I virtualization
 - Type-II virtualization
- Xen created a new real interest
 - Performance (paravirtualization)
 - Open source
 - Linux based
- Interest for HPC
 - VMM bypass
 - Network communication optimization
 - Etc.







Virtual machines

- Basic terminology
 - Host OS: The OS running on a physical machine
 - Guest OS: The OS running on a virtual machine
- Today different approaches
 - Full virtualization: Run an unmodified OS
 - Paravirtualization: Modification of OS for performance
 - Emulation: Host OS and Guest OS can have different architecture
 - Hardware support: Intel-VT, AMD-V



Type-I: Architecture

- Device drivers typically not included in the hypervisor
- Couple hypervisor + Host OS
 - Host OS includes drivers (used by hypervisor)
 - VMs access hardware via the Host OS

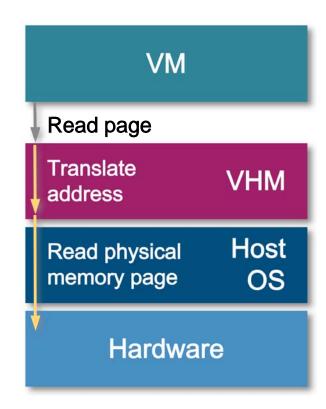
Domain 0 (Host OS)	Арр	Арр		Арр			
	App	App		Арр			
	Guest OS 1 (Domain 1)	Guest OS 2 (Domain 2)		Guest OS n (Domain n)			
Hypervisor (Xen)							
Hardware							

Source: Barney Maccabe, University of New Mexico



Type-II: Architecture

- Simpler model
 - Host OS and the hypervisor are "stacked"
 - No modifications to OSs
 - Provide a BIOS simulation
- Well suited for architecture emulation
 - Ex., PPC on x86_64
- Less efficient than type-I virtualization
 - Especially to paravirtualization





Why a hypervisor specifically for HPC?

- Networking
 - Bridges vs. zero copy (VMM bypass)
 - No RDMA support
- Memory: Important vs. minimal memory footprint
- Processor: Current solutions treat multicores as SMPs
- Tools: No tools available for the management of hundreds of VMs, hypervisors, and Host OSs



Three approaches

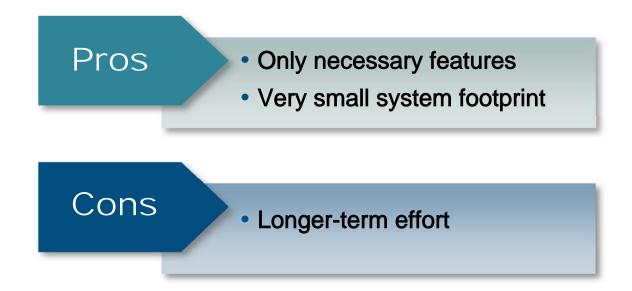
Investigate the development of an HPC hypervisor

- New hypervisor from scratch
- New hypervisor using the microkernel Catamount
- New hypervisor modifying and extending Xen



Hypervisor from scratch

- Develop a new hypervisor using GeekOS
- Current status: A minimal hypervisor has been developed supporting Intel-VT







Hypervisor based on Catamount

- Extend Catamount to
 - Be used as hypervisor
 - As Guest OS
- Current status: Catamount ported to XenoLinux

Pros

- Very small system footprint
- Provide the XT environment within the VMs

Cons

Still based on the Xen hypervisor





Xen-based hypervisor

- Remove unneeded Xen features
- Extend the hypervisor for adaptation (concept of modules)
- Current status
 - Paravirtualization supported
 - Working toward full virtualization (Intel-VT, AMD-V)
- Adaptation capability
- Designed FY 2007
- Implementation FY 2008

Pros

- Quick prototyping
- Compatibility with emerging architectures

Cons

No optimization (yet)



Reaping the benefit of virtualization: Proactive fault tolerance

Context

- Large-scale systems are often subject to failures due to the number of distributed components
- Checkpoint/restart does not scale very well
- Provide capabilities for proactive fault tolerance
 - Failure prediction
 - Migrate application away from faulty node
 - Without stopping application
 - No application code knowledge (or code modification)



Proactive fault tolerance (System and application resilience)

Modular framework

- Support virtualization: Xen, VMM-HPC
- Designed to support process-level checkpoint/restart and migration
- Proactive fault tolerance adaptation: Possible to implement new policies using our SDK

Policy simulator

- Ease the initial phase of study of new policies
- Results from simulator match experimental virtualization results



Management of virtualized environments

- Current issues similar to real systems
 - How to deploy a VM?
 - How to configure a VM?
 - How to deploy multiple VMs?
- Reduce complexity (hide technical details); a VM is just
 - An architecture
 - Some NICs
 - Some memory
 - Etc.



System management issues

Current solutions for virtual environments

- Image repository
- Solutions developed from scratch

Current solutions for standard HPC systems

- Mature system management software, e.g., OSCAR, Rocks
- System definition
- Deployment
- Management of VM images
- Management of Host OS
- Deployment of both Host OSs and VM images
- Distributed environment (clusters of VMs)



OSCAR-V

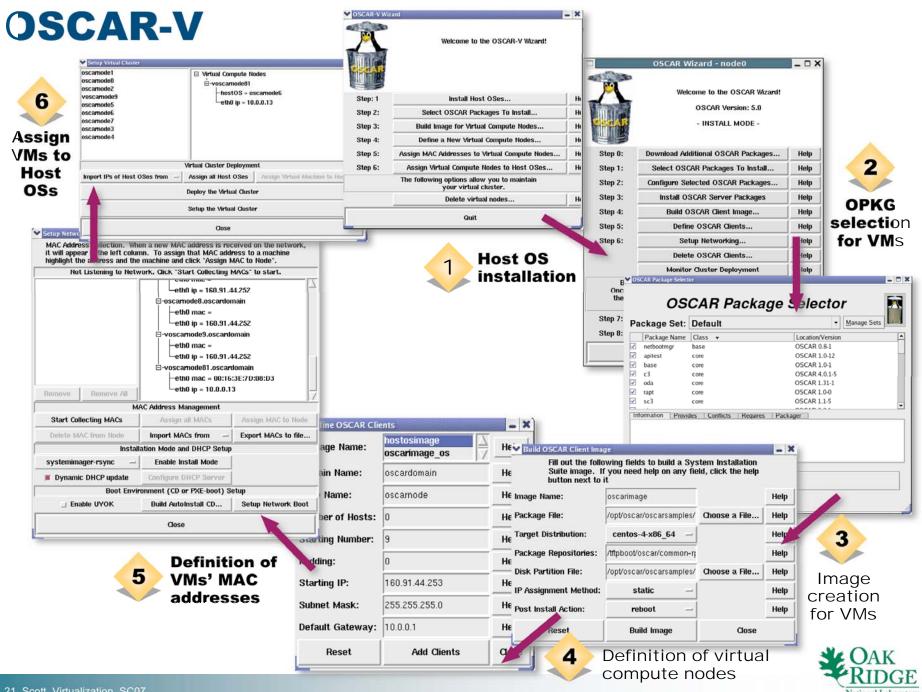
Enhancements to support virtual clusters

- OSCAR-core modifications
- Create OSCAR Packages for virtualization solutions
- Integrate scripts for automatic installation and configuration

Abstracts
differences in virtualization solutions

- Must provide abstraction layer and tools libv3m/v2m
- Enable easy switch between virtualization solutions
- High-level definition and management of VMs: Mem/cpu/etc., start/stop/pause





OSCAR-V: Description of steps

Initial setup

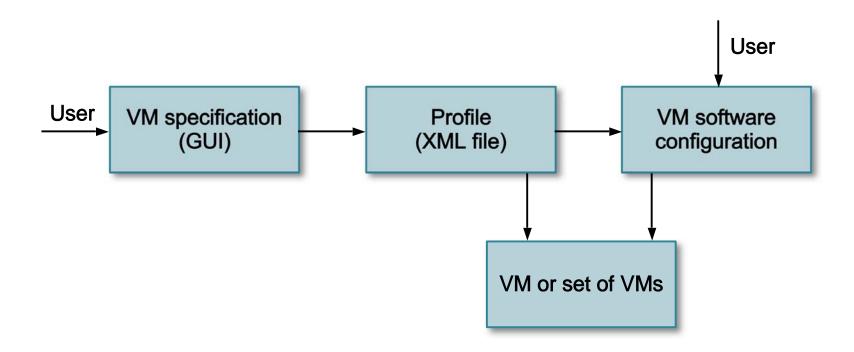
- 1. Install supported distro head node (host)
- 2. Download/set up OSCAR and OSCAR-V
 - OSCAR: Untar oscar-common, oscar-base, etc., and copy distro RPMs
 - OSCAR-V: Untar; run "make install"
- 3. Start Install Wizard
 - run "./oscarv \$network_interface" and follow setups



OSCAR-V: VM profile management

Concept of profiles

- VMs: A profile is memory, disk, OS, NICs, network config.
- Virtual distributed system: A profile is set of VM profiles





OSCAR-V: Virtual machine abstraction

Provide a simple, human-readable VM specification

```
<?xml version="1.0"?>
<!DOCTYPE profile PUBLIC "" "xen_vm.dtd">
cprofile>
    <name>test</name>
    <image size="500">/home/gvallee/vms/test.img</image>
    <nic1>
         <mac>00:02:03:04:05:06</mac>
    </nic1>
</profile>
```



OSCAR-V: V2M – virtual machine management

V2M (Virtual machine management command-line interface)

KVMs (GUI for Linux - KDE/Qt) Applications based on libv3m

High-level interface

V3M Front end

Virtualization abstraction

Qemu

Xen

VMWare

...

V3M Back ends



OSCAR-V: V3M - functionality

- Check the system (files/tools)
- Check the profile (validation)
- Create configuration scripts for VM management
- Provide simple interface for VM management
 - Boot, image management, status
- Switch to a new virtualization solution
 - Change only the "type"



OSCAR-V: V3M – supported features summary

Supported features	Xen (paravirtualization	Xen (full virtualization)	Qemu	VM ware
VM instantiation	Yes	Yes	Yes	Yes
VM image creation	Yes	Yes	Yes	No
Installation via CD-ROM	N/A	Yes	Yes	No
Installation via OSCAR	Yes	Yes	Yes	No
VM migration	Yes	Experimental	No	No
VM pause/unpause	Yes	Experimental	Experimental	Experimental
Virtual disk	Yes	Yes	Yes	Yes



OSCAR-V: Modifications for OSCAR-V

SystemConfigurator modification

- Used after the image copy on the remote node to do local configuration (IP, hostname, etc.)
- Goal: Support Xen specific GRUB entries

Title Xen system

Root (hd0,0)

Kernel /boot/xen.gz dom0_mem=131072

Module /boot/vmlinuz-2.6.12-dom0 root=/dev/sda1 ro

Module /boot/initrd.img-2.6.12

- Add a new option to specify "module" options
- Integrated in SystemConfigurator trunk

kernel_picker modification

- Allow one to include a specific kernel within an image
- Set up a specific SystemConfigurator configuration file
- Add a new option to specify "module" options



OSCAR-V: Image management

Host OS

- OSCAR Packages (OPKG) are available
 - Xen case: Xen hypervisor, Xen kernels (dom0, domU), Xen tools
- Use the unmodified OPKG/OPD mechanism
 - Automatically add software components
 - Automatically set up the virtualization solution
- Current limitation
 - Only REHL, CentOS, Fedora Core are currently supported

Virtual machines

- One OSCAR Package is available
 - Automatically includes the kernel (optional)
 - Automatically sets up the environment
- OSCAR can be used to define VMs
 - Set up the number of VMs
 - MAC addresses
 - IPs

Virtual machines may be deployed



OSCAR-V: Summary

- Capability to create image for Host OSs
 - Minimal image
 - Take benefit of OSCAR features for the deployment
 - Automatic configuration of system-level virtualization solutions
 - Complete networking tools for virtualization solutions
- Capability to create images for VMs
 - May be based on any OSCAR-supported distributions: Mandriva,
 SuSE, Debian, FC, Red Hat EL, etc.
 - Leverage the default OSCAR configuration for compute nodes



OSCAR-V: Current status

- Stabilization for a public release
 - OSCAR-V 1.0
 - Support of Xen full virtualization and paravirtualization
 - CentOS 4.4 x86_64 and x86
 - OSCAR modifications for OSCAR-V (still ongoing)
 - Road map
 - OSCAR-V 2.0: Add support of QEMU and KVM, CentOS 5 (x86_64, x86)
 - Google Summer of Code
 - VM monitoring via V2M
 - Support of other Linux distributions and architecture
 - Stabilize VM migration support (first prototype unsuitable for a public release)

Resources

- V2M/libv3m: http://www.csm.ornl.gov/srt/v2m.html
- OSCAR-V: http://www.csm.ornl.gov/srt/oscarv.html
- OSCAR: http://oscar.openclustergroup.org



OSCAR-V Collaborations: VM deployment on demand

- OSCAR-V does not allow for the automatic deployment of VMs at job submission time
- Integration of Dynamic Virtual Clusters (DVCs)
 - Moab extensions for VM deployment during job submission
 - Use OSCAR images; deployment-based on DVC
 - Collaboration with ASU (Dan Stanzione)



Contacts regarding System-Level Virtualization & OSCAR-V

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