# Estimates for Green's functions of Schrödinger operators; also, a pure mathematician's adventures in wavelet applications

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Consider the operator  $\mathcal{L} = \mathcal{L}_{\alpha} = (-\triangle)^{\alpha/2} - q$ , and the equation

$$\mathcal{L}u = (-\triangle)^{\alpha/2}u - qu = \varphi \quad \text{in } \Omega \subseteq \mathbb{R}^n,$$
 
$$u = 0 \quad \text{on } \partial\Omega,$$

where  $q \in L^1_{loc}(\Omega)$ . We could have  $\Omega = \mathbb{R}^n, 0 < \alpha < n, \ \Omega = \text{half-plane, or e.g., } \Omega = \text{bounded}$ Lipschitz domain, with  $0 < \alpha \leq 2$ .

Remark: Kalton and Verbitsky studied the existence of solutions u>0 to

$$(-\triangle)^{\alpha/2}u - qu^s = \varphi,$$

with s > 1. Curiously, their results did not apply to the linear case s = 1.

Main Equation:

$$\mathcal{L}u = (-\triangle)^{\alpha/2}u - qu = \varphi$$

If  $\alpha = 2$ , then  $\mathcal{L}$  is the time-independent Schrödinger operator  $-\triangle - q$ .

Remark: Case  $\alpha \neq 2$  is of interest in probability, where  $(-\triangle)^{\alpha/2}$  corresponds to  $\alpha$ -stable Lévy processes in the same way that  $-\triangle$  corresponds to Brownian motion.

Joint work with Igor Verbitsky

Main Equation

$$\mathcal{L}u = (-\triangle)^{\alpha/2}u - qu = \varphi \quad \text{in } \Omega \subseteq \mathbb{R}^n$$

Let  $G = G^{(\alpha)}$  be the Green's operator for  $(-\triangle)^{\alpha/2}$  on  $\Omega$ ,

$$G^{(\alpha)}(f)(x) = \int_{\Omega} G^{(\alpha)}(x, y) f(y) \, dy,$$

and we assume  $G^{(\alpha)}(x,y) \geq 0$ . If  $\Omega = \mathbb{R}^n$  then  $G^{(\alpha)}$  is the Riesz potential  $I^{(\alpha)}$  with kernel  $c_n|x-y|^{\alpha-n}$ ; for  $\alpha=2$  and  $n\geq 3$ , G is the Newtonian potential.

We apply G to both sides of the Main Equation to obtain

$$u - G(qu) = G(\varphi) \equiv f$$

Supposing  $q \geq 0$ , let

$$T(u)(x) = G(qu)(x) = \int_{\Omega} G(x, y)u(y)q(y) dy$$

(if q is not non-negative, we let T(u) = G(|q|u) and use this to obtain upper bounds). Then we have u - T(u) = f, or (I - T)(u) = f, so the formal solution is

$$u = (I - T)^{-1}(f) = \sum_{j=0}^{\infty} T^{j}(f).$$

If  $f \in L^2(|q(y)|dy)$  and

$$||T||_{L^2(|q(y)|dy)\to L^2(|q(y)|dy)} < 1,$$

then  $\sum_{j=0}^{\infty} T^j f$  converges in  $L^2(\omega)$  to a solution. Our goal is pointwise estimates. So

$$Tf(x) = \int_{\Omega} G(x, y) f(y) |q(y)| dy$$

$$= \int_{\Omega} G(x,y) f(y) d\omega(y),$$

for  $d\omega(y) = |q(y)| dy$ . Then

$$T^{j}f(x) = \int_{\Omega} G_{j}(x, y) f(y) d\omega(y),$$

where

$$G_1(x,y) = G(x,y),$$

and, for j > 1, inductively define

$$G_j(x,y) = \int_{\Omega} G(x,z)G_{j-1}(z,y) d\omega(z).$$

Let  $V(x,y) = \sum_{j=1}^{\infty} G_j(x,y)$ , so that

$$u(x) = f(x) + \sum_{j=1}^{\infty} T^{j} f(x)$$

$$= f(x) + \int_{\Omega} V(x, y) f(y) d\omega(y).$$

Our goal is to estimate V, the minimal Green's function for  $\mathcal{L}$ .

Theorem A (lower bound). Let  $q \ge 0$ . Then there exist  $c_1, c_2 > 0$  such that

$$V(x,y) \ge c_1 G(x,y) e^{c_2 G_2(x,y)/G(x,y)}$$
.

Remark: Theorem A is relatively easy. The more interesting result of the paper is that under a certain smallness condition on q, we obtain upper bounds of the same form.

Remark: In pretty general circumstances, there is a formula

$$V(x,y) = G(x,y)\mathbb{E}_y^x \left[ e^{\frac{1}{2} \int_0^{\zeta} q(X_s) \, ds} \right],$$

where  $X_t$  is Brownian motion, if  $\alpha=2$ , or an  $\alpha$ -stable symmetric process, if  $0<\alpha<2$ , conditioned to start at x and end at y, and  $\zeta$  is its lifetime. The quantity  $\mathbb{E}^x_y\left[e^{\frac{1}{2}\int_0^\zeta q(X_s)\,ds}\right]$  is called the conditional Feynman-Kac gauge. Our results give estimates for the conditional gauge.

Theorem A follows from a general result about "quasi-metric kernels." Let  $(\Omega, \omega)$  be a  $\sigma$ -finite measure space. A function

$$K: \Omega \times \Omega \longrightarrow (0, \infty]$$

is a quasi-metric kernel on  $\Omega$  if K(x,y)=K(y,x) for all  $x,y\in\Omega$ ,  $K(x,y)<\infty$  if  $x\neq y$ , and there exists  $\kappa>1/2$  such that d(x,y)=1/K(x,y) satisfies

$$d(x,y) \le \kappa(d(x,z) + d(z,y)), \quad x,y,z \in \Omega,$$

for some  $\kappa \geq 1/2$  (we don't require d(x,x) = 0).

Theorem A': Let K be a quasi-metric kernel on a  $\sigma$ -finite measure space  $(\Omega, \omega)$ . Let  $K_1 = K$  and inductively define

$$K_j(x,y) = \int_{\Omega} K(x,z) K_{j-1}(z,y) d\omega(z).$$

Then there exists  $c_2$  depending only on  $\kappa$  such that

$$V(x,y) \equiv \sum_{j=1}^{\infty} K_j(x,y) \ge K(x,y)e^{c_2K_2(x,y)/K(x,y)}$$

Sometimes Theorem A' implies Theorem A directly, with K=G. E.g., for  $\Omega=\mathbb{R}^n$ ,  $0<\alpha< n$ , then  $G(x,y)=c_n|x-y|^{\alpha-n}$  is quasi-metric. However, for domains  $\Omega$ , G may not be quasi-metric. But, for very general domains (including all bounded Lipschitz domains), there exists a function m>0 on  $\Omega$  such that

$$H(x,y) = \frac{G(x,y)}{m(x)m(y)}$$

is a quasi-metric kernel. To get Theorem A, apply Theorem A' with K=H and the measure  $d\nu=m^2d\omega$ , noting that  $H_2/H=G_2/G$ .

Now let's consider the upper estimate for V(x,y). To see what is appropriate, recall the equation

$$u = T(u) + f,$$

where  $T(u)(x) = \int_{\Omega} G(x,y)u(y)d\omega(y)$ . If there exists  $f \geq 0$  such that there is a solution u > 0 to u = T(u) + f, then  $T(u)(x) \leq u(x)$  for all x. Then, by Schur's Lemma,

$$||T||_{L^2(\omega)\to L^2(\omega)} \le 1.$$

If we test the norm on  $\chi_E$ , we obtain

$$\int_{E\times E} G(x,y)d\omega(x)d\omega(y) \le \omega(E)$$

for all measurable  $E \subset \Omega$ .

The first condition is invariant: if H(x,y) = G(x,y)/(m(x)m(y)),  $d\nu = m^2 d\omega$ , and  $S(u)(x) = \int_{\Omega} H(x,y)u(y)d\nu(y)$ , then

$$||S||_{L^2(\nu)\to L^2(\nu)} = ||T||_{L^2(\omega)\to L^2(\omega)}.$$

The second condition is not invariant. Define  $\|\omega\|$  to be the smallest constant C such that

$$\int_{E\times E} G(x,y)m(x)m(y)\,d\omega(y) \le C\int_E m^2(z)\,d\omega(z)$$

for all measurable  $E \subseteq \Omega$ , where m is as above.

We define  $\|\omega\|_*$  similarly, except that we only consider balls B instead of general sets E.

Theorem B (upper bound) Let  $\Omega$  and  $\alpha$  be as above. Then there exists  $\epsilon > 0$  such that if either

(i) 
$$||T||_{L^2(\omega)\to L^2(\omega)} < \epsilon$$
,

(ii) 
$$\|\omega\| < \epsilon$$
, or

(iii)  $\omega$  is a doubling measure and  $\|\omega\|_* < \epsilon$ ,

then there exist  $c_3, c_4 > 0$  such that

$$V(x,y) \le c_3 G(x,y) e^{c_4 G_2(x,y)/G(x,y)}$$
.

As in Theorem A, the result follows from an abstract result for quasi-metric kernels.

Remark: There is interest in when q is sufficiently mild that  $V(x,y) \approx G(x,y)$ . We always have  $V(x,y) \geq G(x,y)$ , and by above, when (i), (ii), or (iii) holds, we obtain  $V \approx G$  if  $G_2 \leq cG$ . This holds e.g., under the Kato condition.

Probably the main interest in our results is in the case where V is not equivalent to G.

Example: Let  $\Omega = \mathbb{R}^n$ ,  $0 < \alpha < n$ , and let  $q = \frac{A}{|x|^{\alpha}}$ , for some constant A, so

$$\mathcal{L} = (-\triangle)^{\alpha/2} - \frac{A}{|x|^{\alpha}}.$$

Then there exists  $\epsilon > 0$  such that for  $0 \le A < \epsilon$ ,

$$c_1 \frac{\left(\max\left\{\left|\frac{x}{y}\right|,\left|\frac{y}{x}\right|\right\}\right)^{c_2}}{|x-y|^{n-\alpha}} \le V(x,y) \le c_3 \frac{\left(\max\left\{\left|\frac{x}{y}\right|,\left|\frac{y}{x}\right|\right\}\right)^{c_4}}{|x-y|^{n-\alpha}},$$

for some  $c_1, c_2, c_3, c_4 > 0$ .

Proof: After some elementary computations, obtain

$$G_2(x,y)/G(x,y) \approx 1 + \log\left(\max\left\{\left|\frac{x}{y}\right|,\left|\frac{y}{x}\right|\right\}\right).$$

Apply Theorems A and B. Here  $\omega$  is doubling, and one can check the condition  $\|\omega\|_* \leq CA^2$ .

Note that we don't get sharp powers, and we require the smallness condition on q. However, our results work for very general  $\Omega$  and for a range of  $\alpha$ .

In the literature, q is often assumed bounded, or very nice. From our results, we can see what singularities of q are feasible, and the general form of V.

Of course, estimates on V yield solvability results for the original equation  $\mathcal{L}_{\alpha}u = \varphi$ , because  $u(x) = G(\varphi)(x) + \int_{\Omega} V(x,y)G(\varphi)(x)q(y) dy$ .

Comments on proof of abstract theorem on quasi-metric kernels:

Actually show

$$V(x,y) \approx \int_{d(x,y)}^{\infty} \frac{e^{c(G_t(x) + G_t(y))}}{t^2} dt \approx K e^{c_1 K_2/K},$$

where

$$G_t(x) = \int_0^t \frac{\omega(B_r(x))}{r^2} dr.$$

$$V(x,y) \approx \int_{d(x,y)}^{\infty} \frac{e^{c(G_t(x)+G_t(y))}}{t^2} dt$$

For the lower estimate, prove inductively that

$$K_j(x,y) \ge c^{j-1} \int_{d(x,y)}^{\infty} \frac{G_t(y)^{j-1}}{(j-1)!t^2} dt,$$

using an integration by parts. Then sum on j, obtaining  $e^{cG_t(y)}$ , use symmetry to get  $e^{cG_t(x)}$ , average, and use inequality between arithmetic and geometric means.

$$V(x,y) \approx \int_{d(x,y)}^{\infty} \frac{e^{c(G_t(x) + G_t(y))}}{t^2} dt$$

For the upper estimate, inductively prove

$$K_j(x,y) \le c_1 \left(\frac{c_2}{\beta}\right)^{j-1} \int_{d(x,y)}^{\infty} \frac{e^{\beta(G_t(x) + G_t(y))}}{t^2} dt.$$

Then for  $\beta$  large enough, sum on j. For the induction, need an estimate of the form

$$\int_{B_t(x)} e^{\beta G_t} d\omega \le c\omega(B_{2t}(x)).$$

We get this from

$$\int_{B_t(x)} G_t^m d\omega \le m! C^m \|\omega\|^m \omega(B_{2t}(x)).$$

We need  $\|\omega\| < \epsilon$  to sum on m.

## **Transient Signal Detection**

(Daniel Wagner Associates, 1990-1991)

Have noisy environment generating random discrete noise signals

$$n = (n(0), n(1), n(2), ..., n(N-1)).$$

Assume n is a jointly Gaussian random variable.

### Want to detect a given prototype signal

$$s = (s(0), s(1), s(2), ..., s(N-1)).$$

### Suppose we receive signal

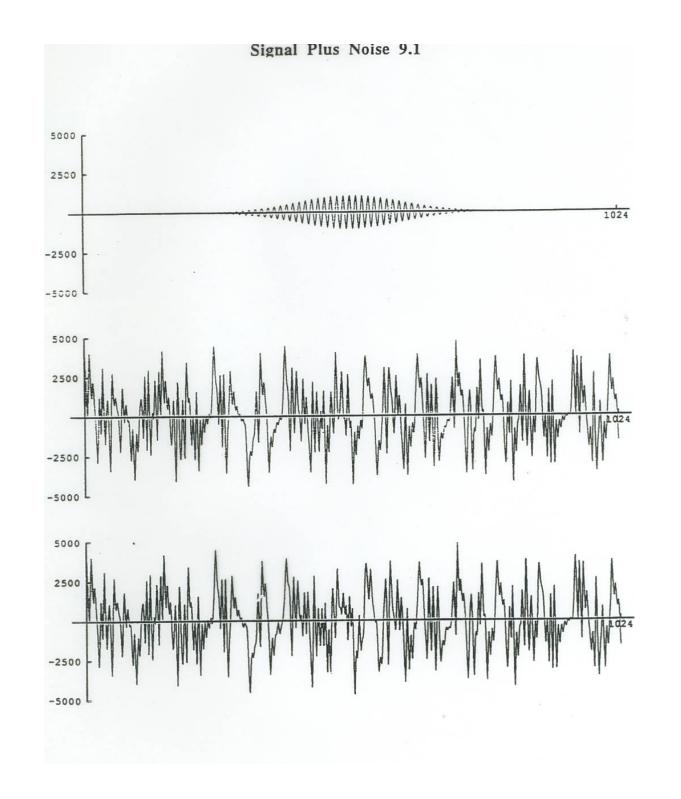
$$x = (x(0), x(1), x(2), ..., x(N-1)).$$

There are two possibilities:

 $H_0: x = n$  (null hypothesis: s not present in x)

 $H_1: x = s + n$  (s is present in x).

Need: hypothesis test.



A "decision criterion" assigns to each x a conclusion, either  $H_0$  or  $H_1$ .

There are two types of mistakes:

False alarm: Accept  $H_1$  when  $H_0$  is true

Detection failure: Accept  $H_0$  when  $H_1$  is true

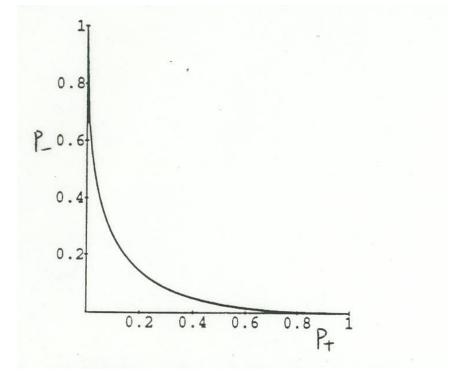
Let  $P_+ = \text{Prob}(\text{False alarm})$ 

and  $P_{-} = Prob(Detection failure)$ .

Want to be able to assign  $P_+$ .

A "test" assigns to each  $P_+ \in [0,1]$  a decision criterion having that value of  $P_+$ .

To evaluate a test, plot  $P_{-}$  as a function of  $P_{+}$ :



This is called the ROC (Receiver Operating Characteristic) curve for the test. One test is clearly superior to another if its ROC curve is everywhere lower.

A well-known, relatively elementary statistics result states that there is an optimal test for this problem, i.e., a test with lower ROC curve than any other test. This test is usually called the "matched filter" test.

### Matched Filter Test

Let  $R = (R_{i,j})_{i,j=0,1,2,3,...,N-1}$  be the noise correlation matrix, i.e.,

$$R_{i,j} = E(n(i)n(j)).$$

Let 
$$\lambda = \langle R^{-1}s, x \rangle = \sum_{i=0}^{N-1} (R^{-1}s)(i)x(i)$$
.

Matched filter test: Accept  $H_1 \iff \lambda \geq \beta$ .

Choice of parameter  $\beta$  determines  $P_+$ , hence a point on the ROC curve.

Problem: May have very large number of prototype signals s we want to test for. May not be able to compute all the test statistics

$$\lambda = \langle R^{-1}s, x \rangle = \sum_{i=0}^{N-1} (R^{-1}s)(i)x(i).$$

in real time.

Plan: Compress the test. Select a subset M of 0, 1, 2, ..., N-1, and compute

$$\lambda_M = \sum_{i \in M} (R^{-1}s)(i)x(i).$$

If most of the information is contained in a small number of terms, and those terms are included in the sum by the choice of M, we can get a good approximation to the optimal test with much less computation. But the information may not be concentrated in a few terms in the standard basis.

Idea: change basis for the compressed test. Apply any orthogonal change of basis to R, s, and x, to get  $\widehat{R}, \widehat{s}$ , and  $\widehat{x}$ . Then can compute  $\lambda$  in new basis:

$$\lambda = \langle R^{-1}s, x \rangle = \langle \widehat{R}^{-1}\widehat{s}, \widehat{x} \rangle.$$

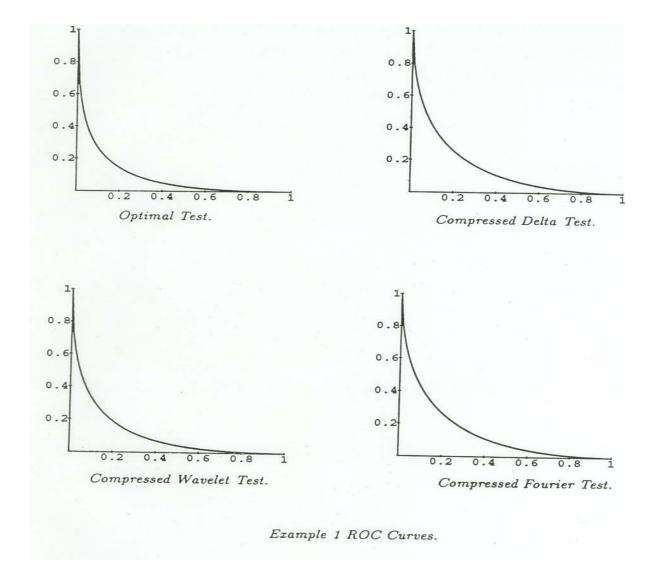
So full test is basis invariant (must be, since optimal). But compression of optimal test is not basis invariant.

Heuristic: If prototype signal is localized in space (a transient signal) and has definite frequency characteristics, and/or if the noise has definite frequency characteristics, then a wavelet basis should do a better job of compressing the optimal test.

### Examples: N = 64, Cardinality of M = 4

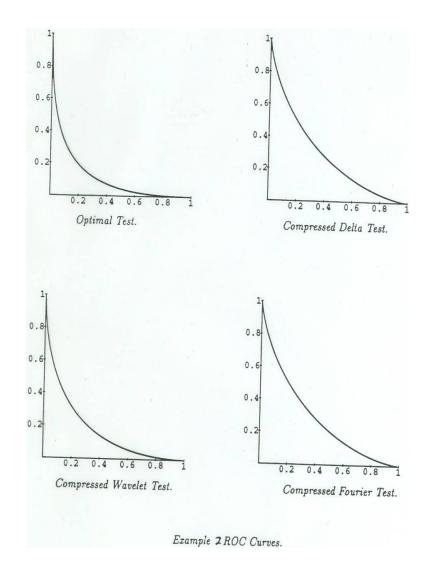
1.)  $s(n) = .8\cos(n\pi/6)e^{-(n-31)^2/32}$ , white

noise



Examples: N = 64, Cardinality of M = 4

2.)  $s(n) = .8e^{-(n-31)^2/4}$ , correlated noise



# Car Rattles and the Shift-Invariant Discrete Wavelet Transform

2001 Michigan State University

Ford Representative: David Scholl

Masters in Industrial Mathematics Team:

Joerg Enders

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Faculty Advisor: Michael Frazier

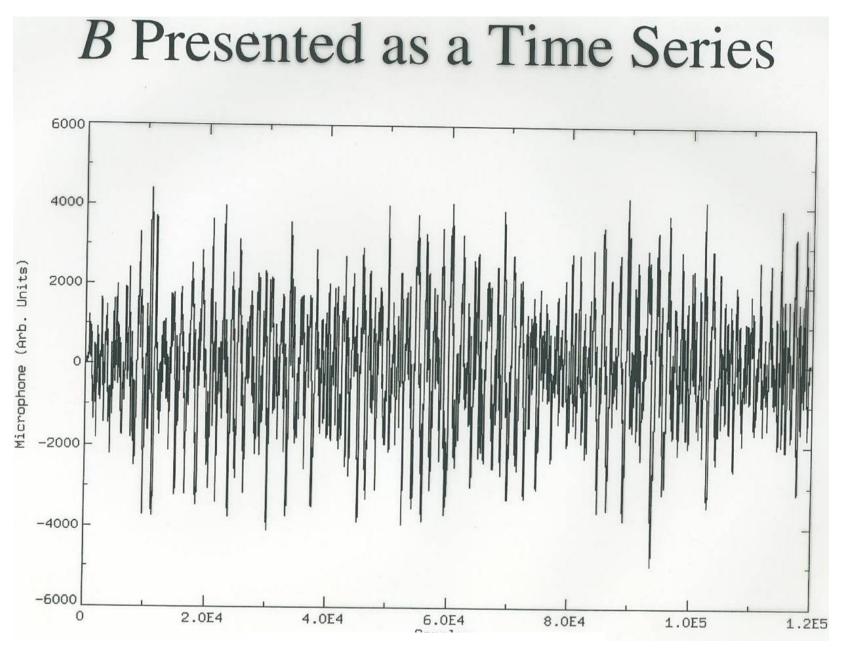
Project: Analyze and diagnose car rattles electronically

Example: Creak 4

Example: Clatter



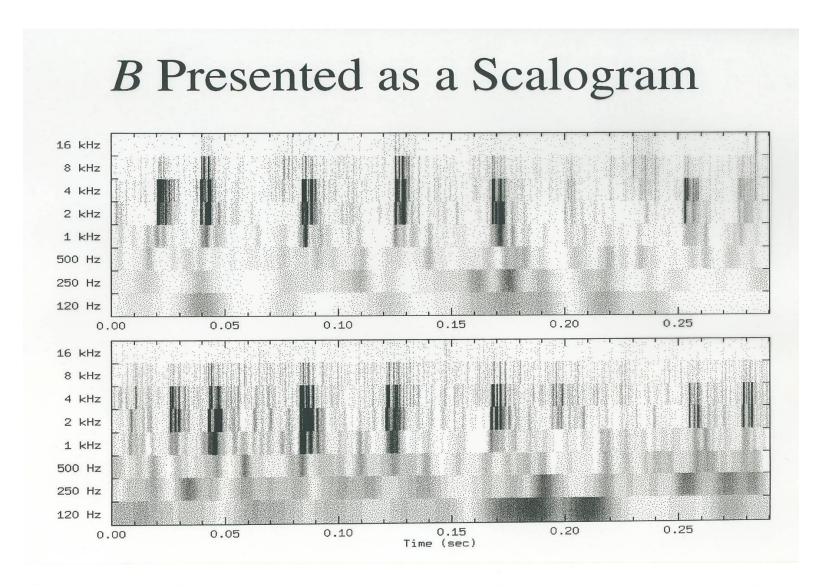
What is right representation to allow extraction of main features from the noise?



Time series: unclear

# B Presented as a Spectral Density 1.0E4 (Arb. Units) 1.0E2 Energy Spectral Density ( 100 1000 10000 Frequency (Hz)

Frequency Representation (spectrogram): unclear



Discrete Wavelet Representation (Top): Better

Shift-Invariant Discrete Wavelet Representation (Bottom): Even Better

If z is a vector of length N, wavelet transform of z is a vector of length N, and the wavelet transform is an invertible linear map. In fact, the inverse can be computed rapidly via convolutions, in O(N) steps.

The Shift-Invariant Discrete Wavelet Transform (SIDWT) of z is obtained (roughly) by averaging the wavelet transform of z over all N translations of z. The SIDWT can be computed in  $O(N\log^2 N)$  steps. However, the output is a vector of length  $N\log_2 N$ . Thus the SIDWT is a linear map from a lower-dimensional space into a much larger dimensional space. Hence it is not invertible.

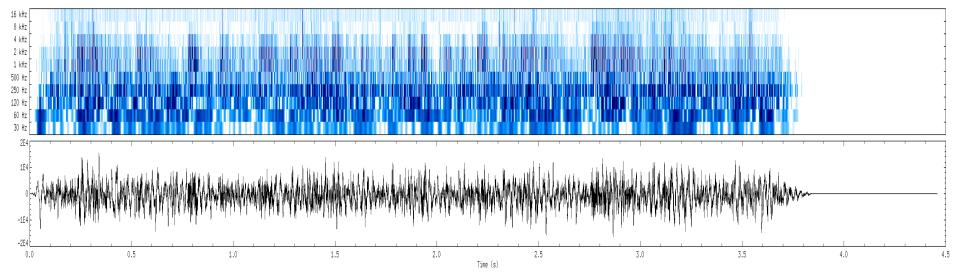
In linear algebra, in this situation, one learns to use the pseudoinverse. If  $T: \mathbb{R}^n \to \mathbb{R}^m$  is linear and 1-1, where n < m, define the pseudoinverse  $S: \mathbb{R}^m \to \mathbb{R}^n$  as follows. For a point w in  $\mathbb{R}^m$ , find its orthogonal projection u on the range of T, and define S(w) to be the unique  $z \in \mathbb{R}^n$  such that T(z) = u.

Problem: Formula for S is  $S = (T^*T)^{-1}T^*$ . Here T is a matrix of size  $N \log_2 N \times N$ , so  $T^*T$  is  $N \times N$ , where here typically  $N \approx 10^5$  or  $10^6$ . So the matrix is too large to invert rapidly and accurately.

My group figured out that the pseudoinverse of the SIDWT is also computable fast via convolutions, in fact in  $O(N \log^2 N)$  steps. This allowed Dave Scholl to continue his examples.

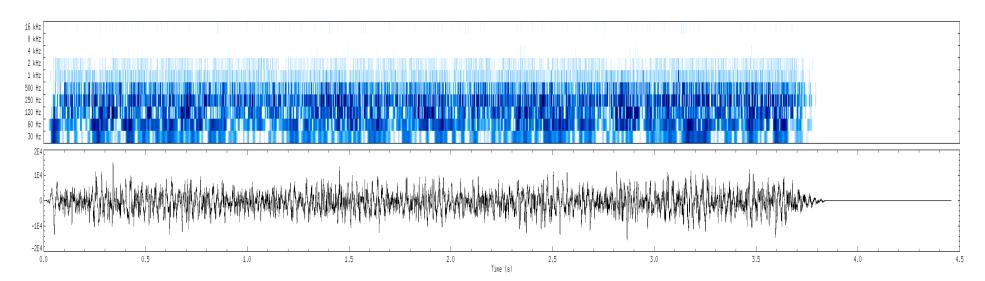
Creak: Original





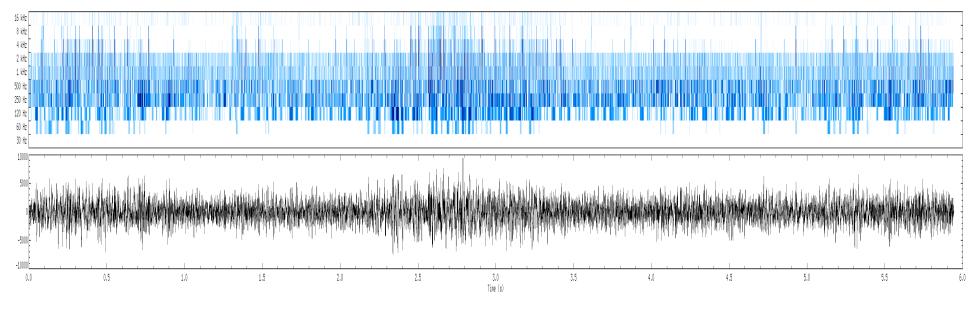
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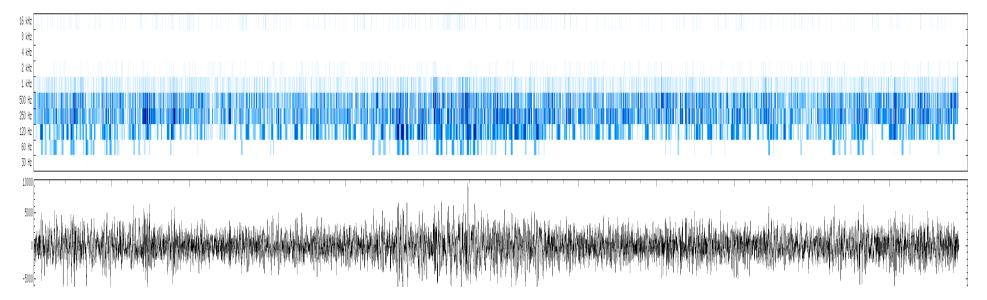
Clatter: Original





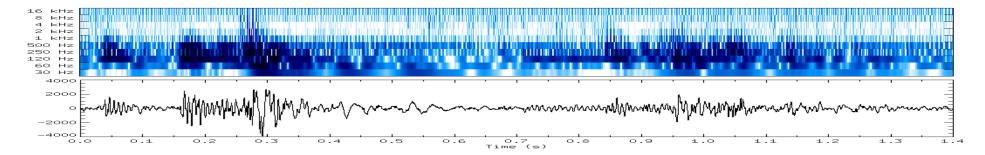
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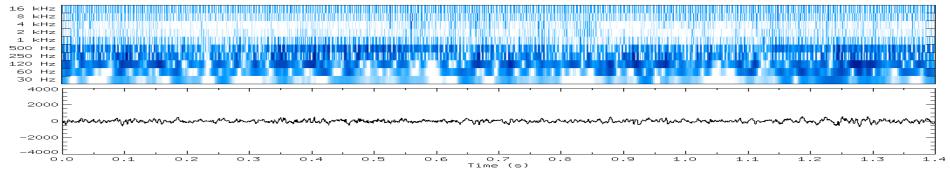
### Windshield Wiper Reversal Thud: Original





#### Windshield Wiper Reversal Thud: Background





#### Windshield Wiper Reversal Thud: Thud



