

Towards the Final Generation of Dense Linear Algebra Libraries

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Current FLAME Team

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n Industry

Hewlett-Packard National Instruments IBM
Intel NEC Solutions (America), Inc

Support

- n National Science Foundation
 - n Modest Funding through 2006
- n Hewlett-Packard
 - n Equipment donations
- n Unrestricted grants
 - n Dr. James Truchard (National Instruments)
 - n NEC Solutions (America), Inc

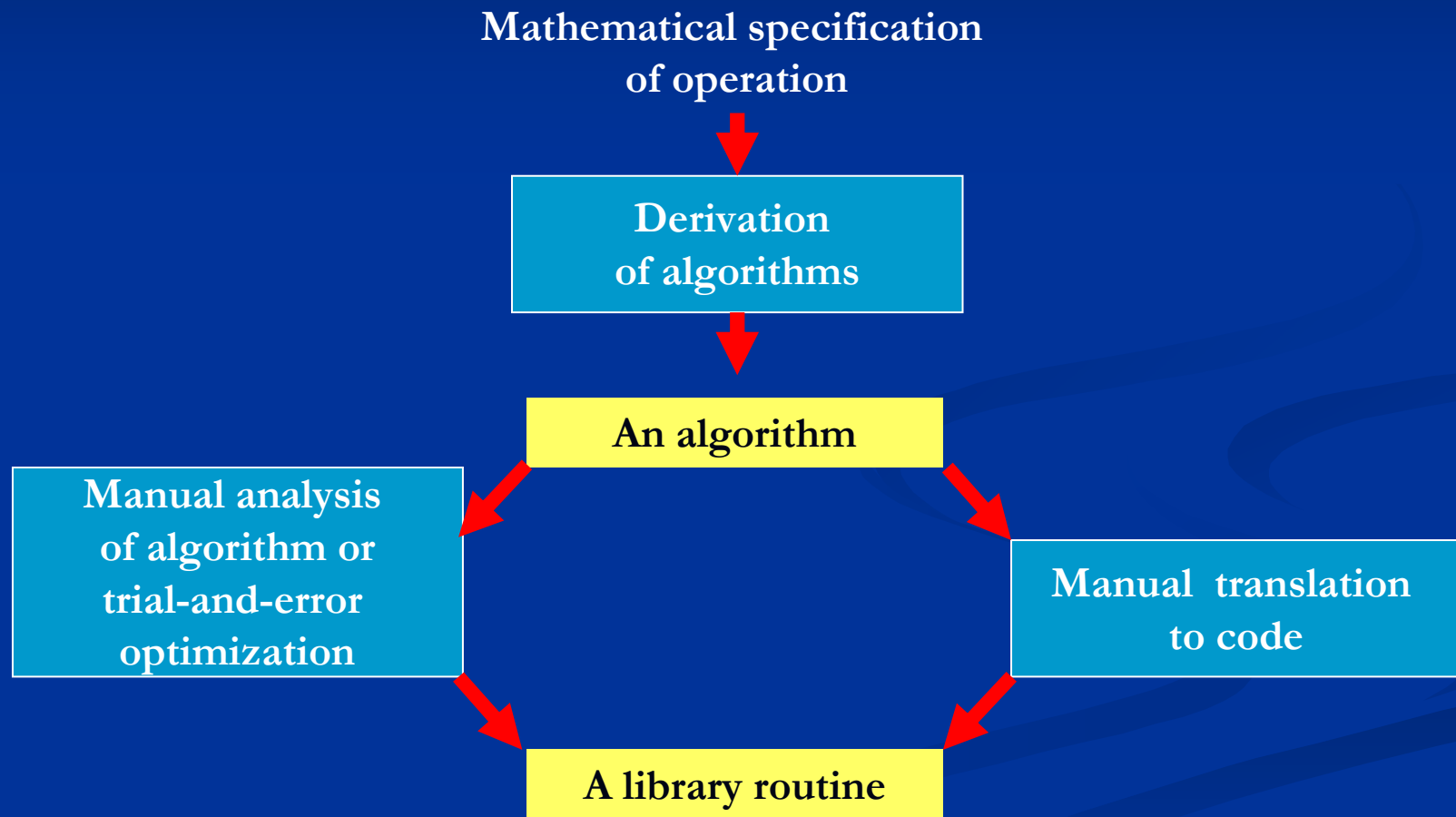
Motivation

- n Developing dense linear algebra libraries
 - n Traditional approach:
 - n Evolve from existing libraries
 - n Ask the question: What added functionality is needed?
 - n Reactive to current needs
 - n Always a “Next Generation” library
 - n The Big Question:
 - n Can we build a “Final Generation” library?
 - n Proactive to (as of yet undefined) future needs

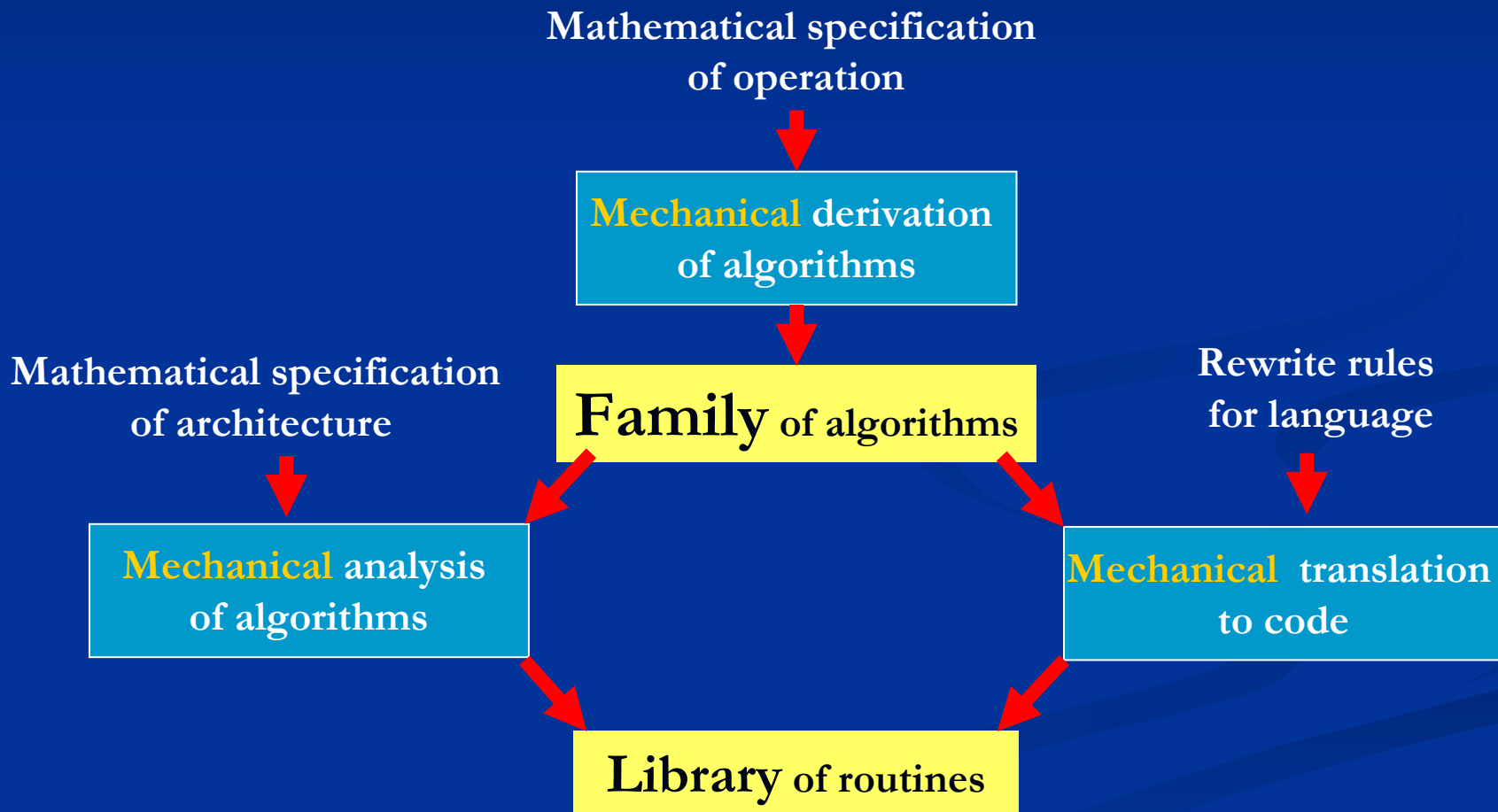
Motivation

- n Properties of a Final Generation Library
 - n Forward compatible to
 - n New architectures
 - n New languages
 - n New datastructures
 - n New operations

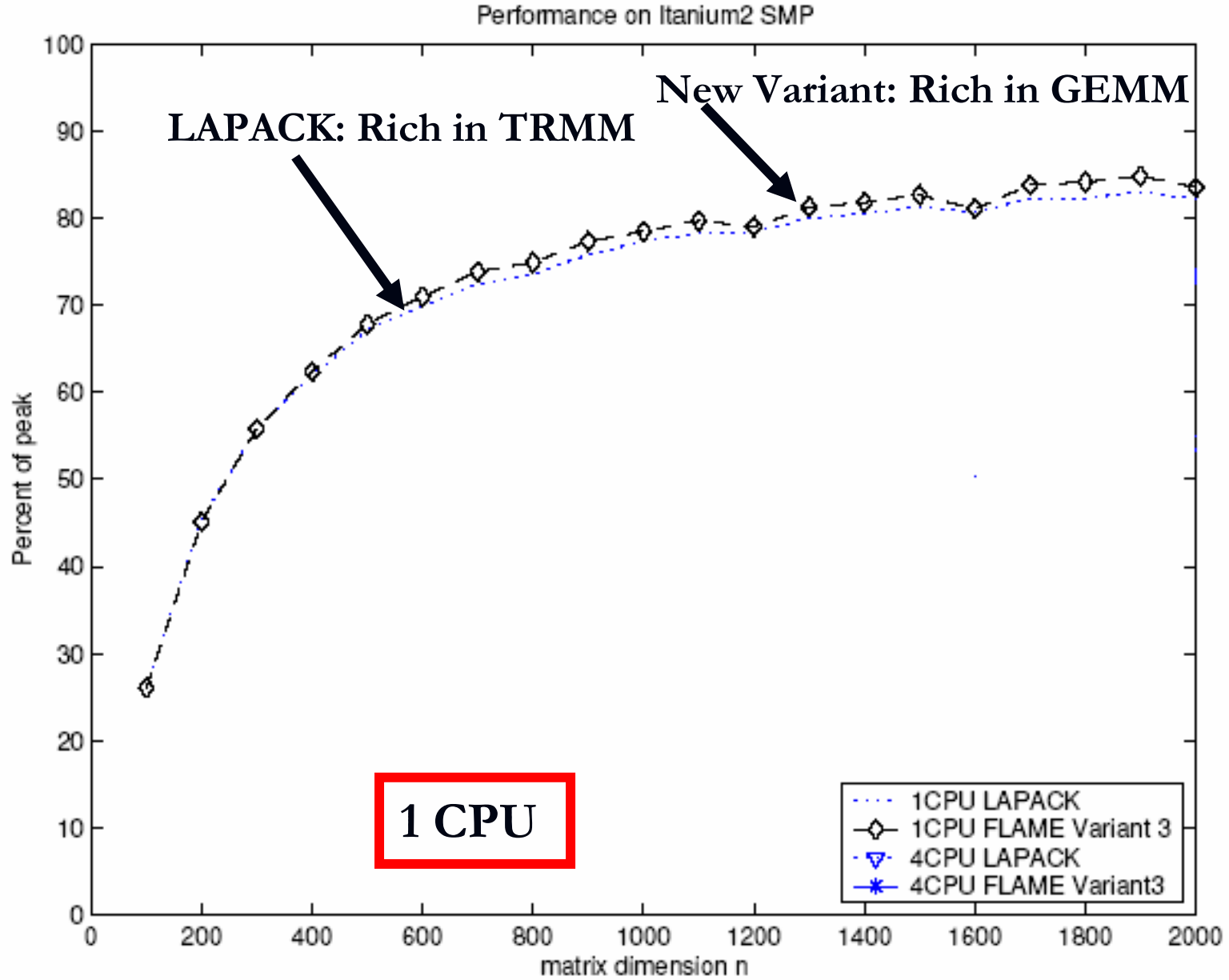
What do we do by hand?

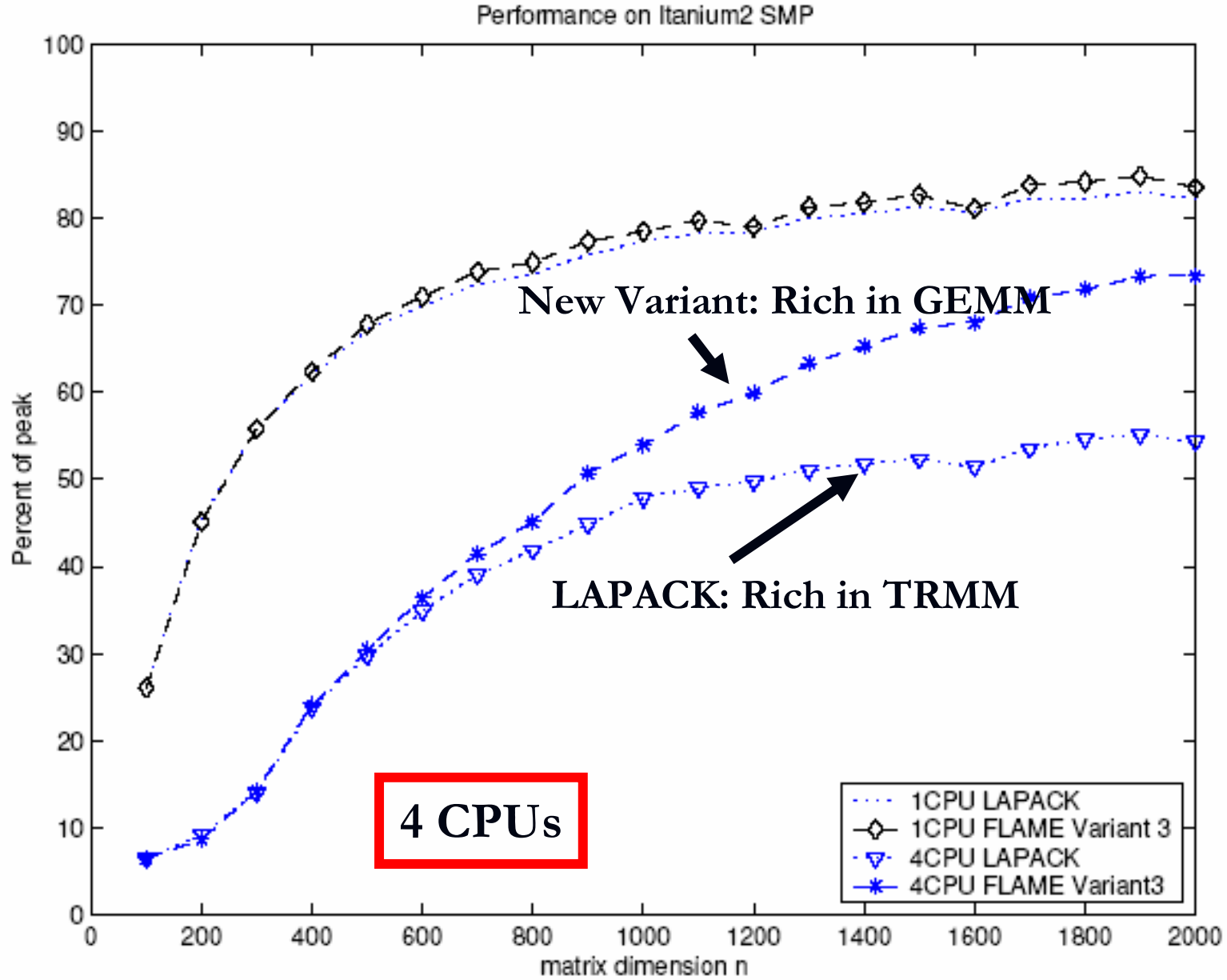


The Final Generation



Why a family of
algorithms?

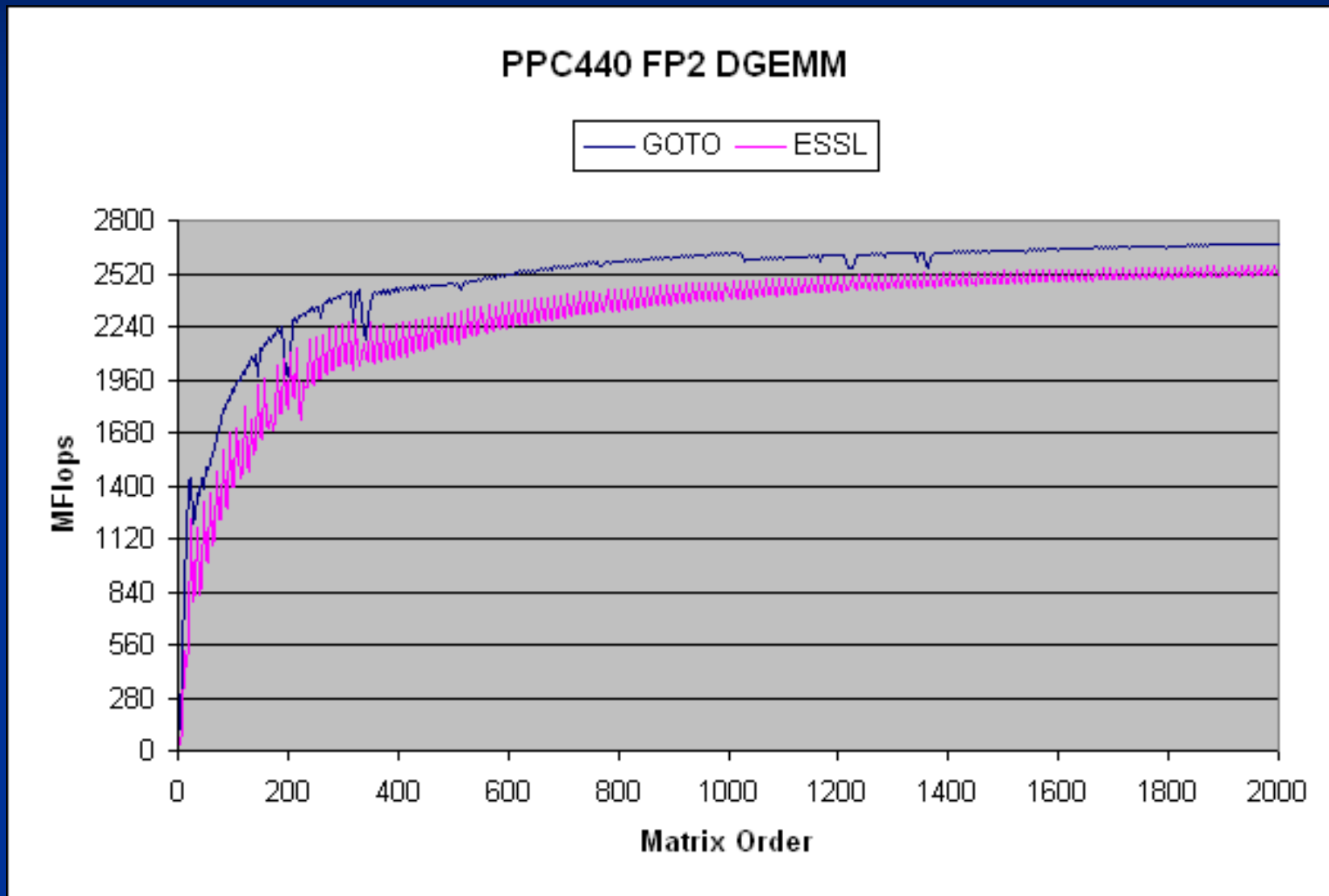




A Note on Performance

- n The FLAME team is recognized for performance:
 - n GOTO BLAS by Kazushige Goto:
 - n Fastest for essentially all platforms
 - n FLAME:
 - n Outperforms LAPACK
 - n PLAPACK:
 - n Outperforms ScaLAPACK for all major operations
- n I will show few performance graph

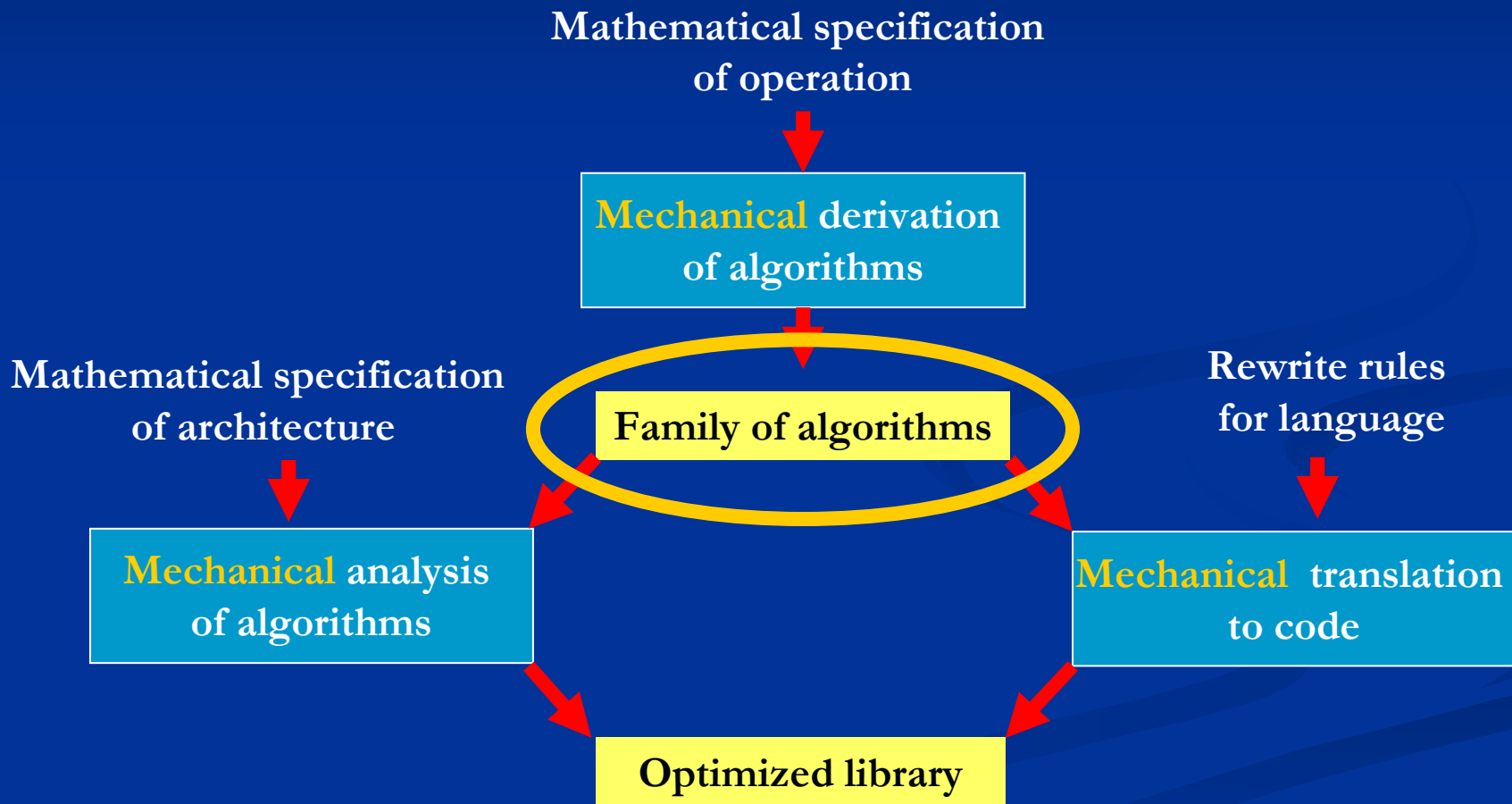
Kazushige Goto's BLAS



Overview

- n New Notation for Expressing Algorithms
- n APIs for Representing Algorithms in Code
- n Mechanical Derivation of Algorithms
- n Mechanical Analysis of Algorithms
- n Addressing Future Challenges
- n Conclusion

The Final Generation

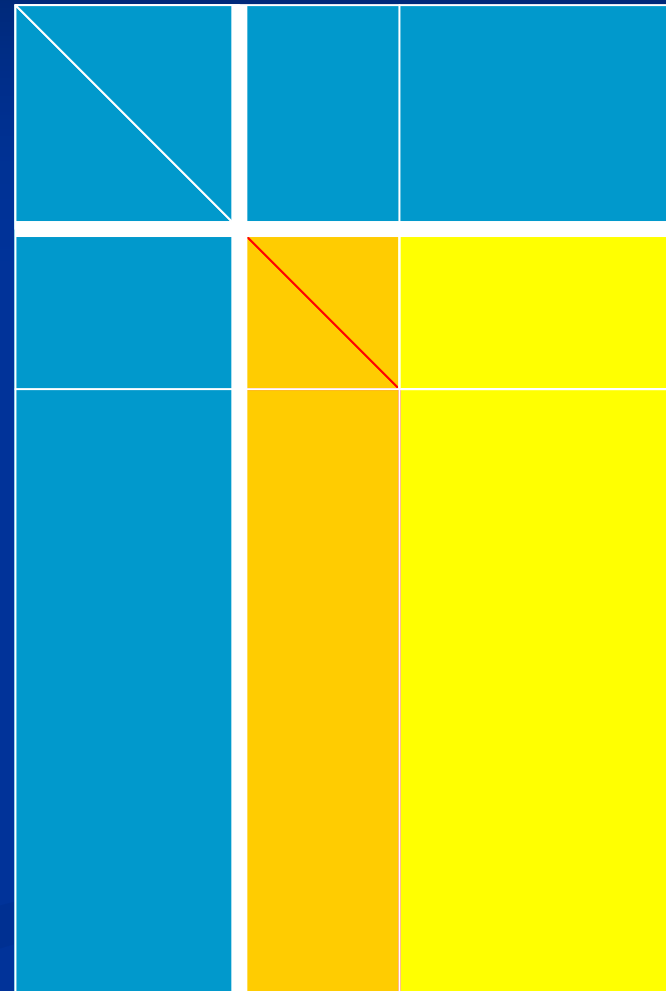


Step 1:

**Change the Notation for
Expressing Algorithms**

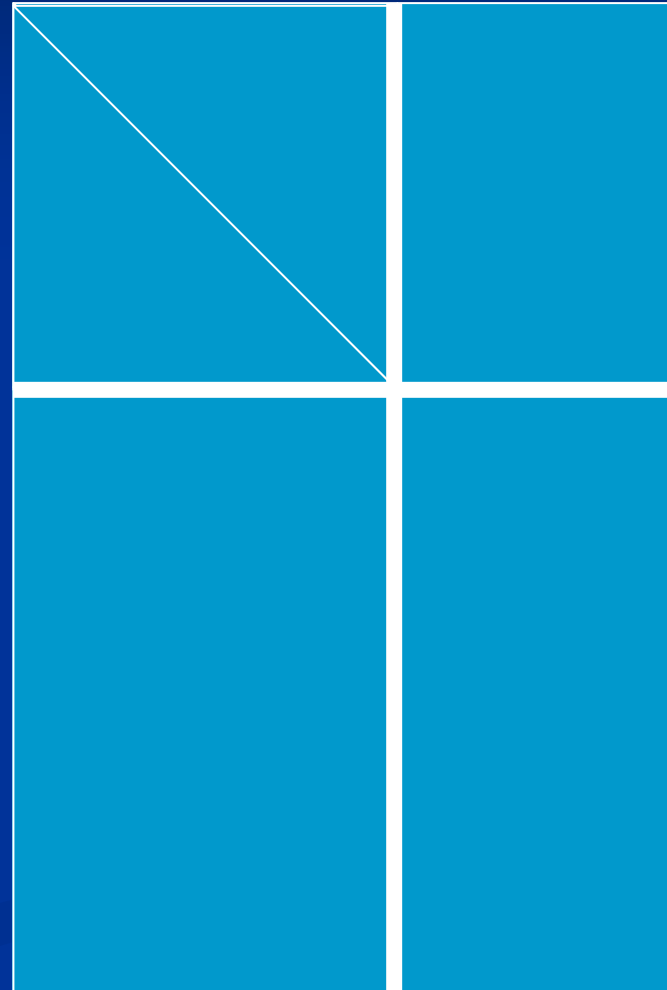
Example: QR factorization via Householder Transformations

- n Blocked Algorithm:
 - n Factor current panel
 - n Form compact WY transform
 - n Update rest of matrix

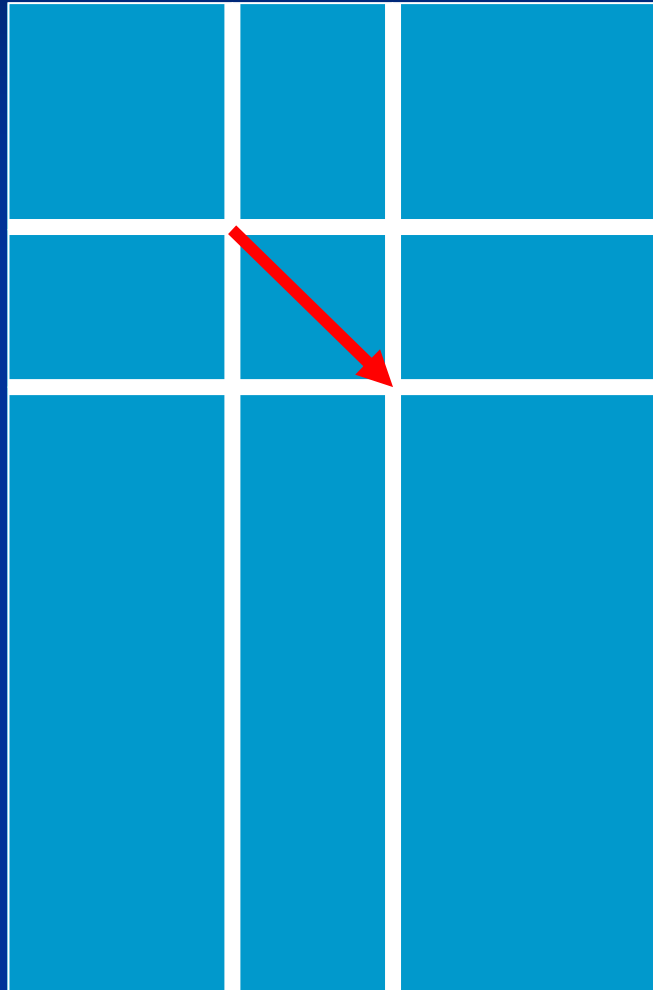


QR factorization via Householder transformations

- n Blocked Algorithm:
 - n Factor current panel
 - n Form compact WY transform
 - n Update rest of matrix
 - n Move forward

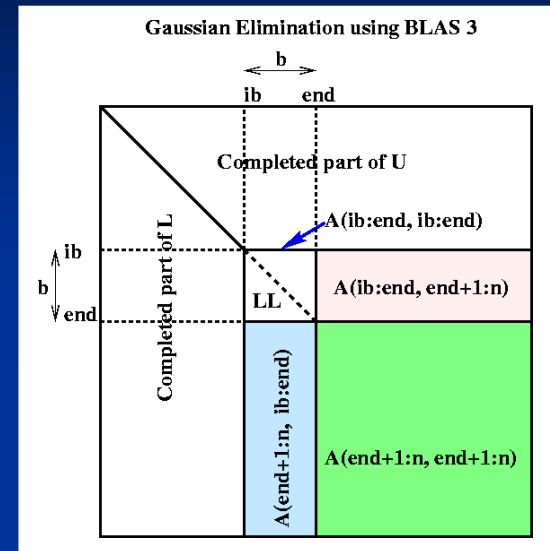
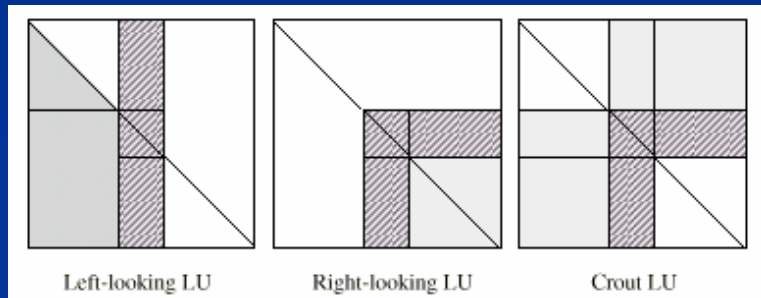


Capturing movement through matrices

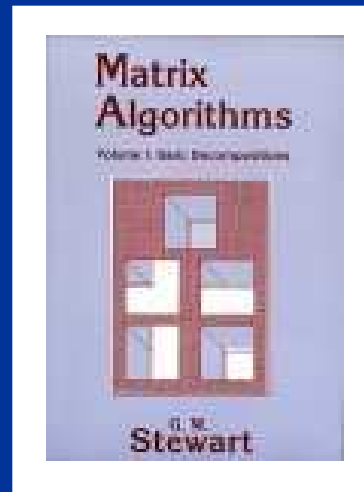


This picture has been around:

- n In a typical talk on LAPACK:



- n Pete Stewart's recent book:



Can the picture be the
algorithm?

Algorithm: $[A, s] := \text{QRBLK}(A)$

Partition $A \rightarrow \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$ and $t \rightarrow \left(\begin{array}{c} s_T \\ \hline s_B \end{array} \right)$

where A_{TL} is 0×0 and s_T has 0 elements

while $n(A_{BR}) \neq 0$ **do**

Determine block size k

Repartition

$\left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right) \rightarrow \left(\begin{array}{c|c|c} A_{00} & A_{01} & A_{02} \\ \hline A_{10} & A_{11} & A_{12} \\ \hline A_{20} & A_{21} & A_{22} \end{array} \right)$ and $\left(\begin{array}{c} s_T \\ \hline s_B \end{array} \right) \rightarrow \left(\begin{array}{c} s_0 \\ \hline s_1 \\ \hline s_2 \end{array} \right)$

where A_{11} is $k \times k$ and s_1 has k elements

$\left[\left(\begin{array}{c} A_{11} \\ \hline A_{21} \end{array} \right), s_1 \right] := \left[\left(\begin{array}{c} \{U \setminus R\}_{11} \\ \hline U_{21} \end{array} \right), s_1 \right] = \text{QRUNB} \left(\left(\begin{array}{c} A_{11} \\ \hline A_{21} \end{array} \right) \right)$

Compute S_1 from $\left[\left(\begin{array}{c} U_{11} \\ \hline U_{21} \end{array} \right), s_1 \right]$

Update

$\left(\begin{array}{c} A_{12} \\ \hline A_{22} \end{array} \right) := \left(I + \left(\begin{array}{c} U_{11} \\ \hline U_{21} \end{array} \right) S_1 \left(\begin{array}{c} U_{11} \\ \hline U_{21} \end{array} \right)^T \right)^T \left(\begin{array}{c} A_{12} \\ \hline A_{22} \end{array} \right)$

Continue with

$\left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right) \leftarrow \left(\begin{array}{c|c|c} A_{00} & A_{01} & A_{02} \\ \hline A_{10} & A_{11} & A_{12} \\ \hline A_{20} & A_{21} & A_{22} \end{array} \right)$ and $\left(\begin{array}{c} s_T \\ \hline s_B \end{array} \right) \leftarrow \left(\begin{array}{c} s_0 \\ \hline s_1 \\ \hline s_2 \end{array} \right)$

endwhile

Partition $A \rightarrow \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$

where A_{TL} is 0×0

while $n(A_{BR}) \neq 0$ do

Determine block size b

Repartition

$$\left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right) \rightarrow \left(\begin{array}{c|c|c} A_{00} & A_{01} & A_{02} \\ \hline A_{10} & A_{11} & A_{12} \\ \hline A_{20} & A_{21} & A_{22} \end{array} \right)$$

where A_{11} is $b \times b$

$$\left[\left(\begin{array}{c} A_{11} \\ A_{21} \end{array} \right), s_1 \right] := \left[\left(\begin{array}{c} \{U \setminus R\}_{11} \\ U_{21} \end{array} \right), s_1 \right] = \text{QR} \left(\begin{array}{c} A_{11} \\ A_{21} \end{array} \right)$$

Compute S_1 from $\left[\left(\begin{array}{c} U_{11} \\ U_{21} \end{array} \right), s_1 \right]$

Update

$$\left(\begin{array}{c} A_{12} \\ A_{22} \end{array} \right) := \left(I + \left(\begin{array}{c} U_{11} \\ U_{21} \end{array} \right) S_1 \left(\begin{array}{c} U_{11} \\ U_{21} \end{array} \right)^T \right)^T \left(\begin{array}{c} A_{12} \\ A_{22} \end{array} \right)$$

Continue with

$$\left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right) \leftarrow \left(\begin{array}{c|c|c} A_{00} & A_{01} & A_{02} \\ \hline A_{10} & A_{11} & A_{12} \\ \hline A_{20} & A_{21} & A_{22} \end{array} \right)$$

endwhile

Partition $A \rightarrow \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right)$
 where A_{TL} is 0×0

while $n(A_{BR}) \neq 0$ do

Determine block size b

Repartition

$$\left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right) \rightarrow \left(\begin{array}{c|c|c} A_{00} & A_{01} & A_{02} \\ \hline A_{10} & A_{11} & A_{12} \\ \hline A_{20} & A_{21} & A_{22} \end{array} \right)$$

where A_{11} is $b \times b$

$$\left[\left(\begin{array}{c} A_{11} \\ \hline A_{21} \end{array} \right), s_1 \right] := \left[\left(\begin{array}{c} \{U \setminus R\}_{11} \\ \hline U_{21} \end{array} \right), s_1 \right] = \text{QR} \left(\begin{array}{c} A_{11} \\ \hline A_{21} \end{array} \right)$$

Compute S_1 from $\left[\left(\begin{array}{c} U_{11} \\ \hline U_{21} \end{array} \right), s_1 \right]$

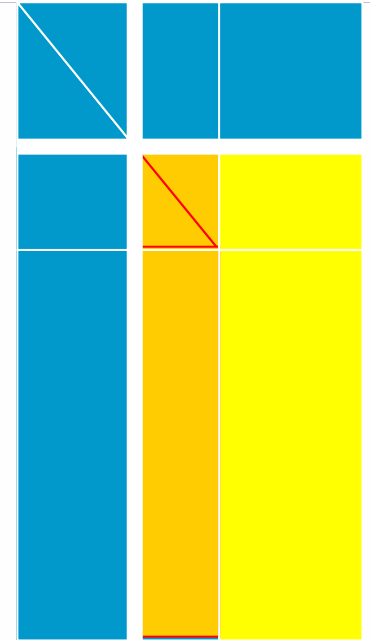
Update

$$\left(\begin{array}{c} A_{12} \\ \hline A_{22} \end{array} \right) := \left(I + \left(\begin{array}{c} U_{11} \\ \hline U_{21} \end{array} \right) S_1 \left(\begin{array}{c} U_{11} \\ \hline U_{21} \end{array} \right)^T \right)^T \left(\begin{array}{c} A_{12} \\ \hline A_{22} \end{array} \right)$$

Continue with

$$\left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right) \leftarrow \left(\begin{array}{c|c|c} A_{00} & A_{01} & A_{02} \\ \hline A_{10} & A_{11} & A_{12} \\ \hline A_{20} & A_{21} & A_{22} \end{array} \right)$$

endwhile



Algorithm: $[A, t] := QR(A)$

```
[ ATL, ATR, ...  
  ABL, ABR ] = FLA_Part_2x2( A, 0, 0, 'FLA_TL' );
```

```
while ( size( ATL, 2 ) ~= size( A, 2 ) )
```

```
  b = min( size( ABR, 1 ), nb_alg );
```

```
  [ A00, A01, A02, ...  
    A10, A11, A12, ...  
    A20, A21, A22 ] = FLA_Repart_2x2_to_3x3( ATL, ATR, ...  
                                              ABL, ABR, ...
```

```
  [ U1, s1 ] = QR_unb_var1( [ A11          b, b, 'FLA_BR' );  
                          A21 ], s1 );
```

```
  [ A11, ...  
    A21 ] = FLA_Part_2x1( U1, b, 'FLA_TOP' );
```

```
  S1 = Accum_S( U1, s1 );
```

```
  U11 = trilu( A11 );
```

```
  U21 = A21;
```

```
  W12 = S1' * ( U11' * A12 + U21' * A22 )
```

```
  A12 = A12 - U11 * W12;
```

```
  A22 = A22 - U21 * W12;
```

Update

$$\begin{pmatrix} A_{12} \\ A_{22} \end{pmatrix} := \left(I + \begin{pmatrix} U_{11} \\ U_{21} \end{pmatrix} S_1 \begin{pmatrix} U_{11} \\ U_{21} \end{pmatrix}^T \right)^T \begin{pmatrix} A_{12} \\ A_{22} \end{pmatrix}$$

```
  [ ATL, ATR, ...
```

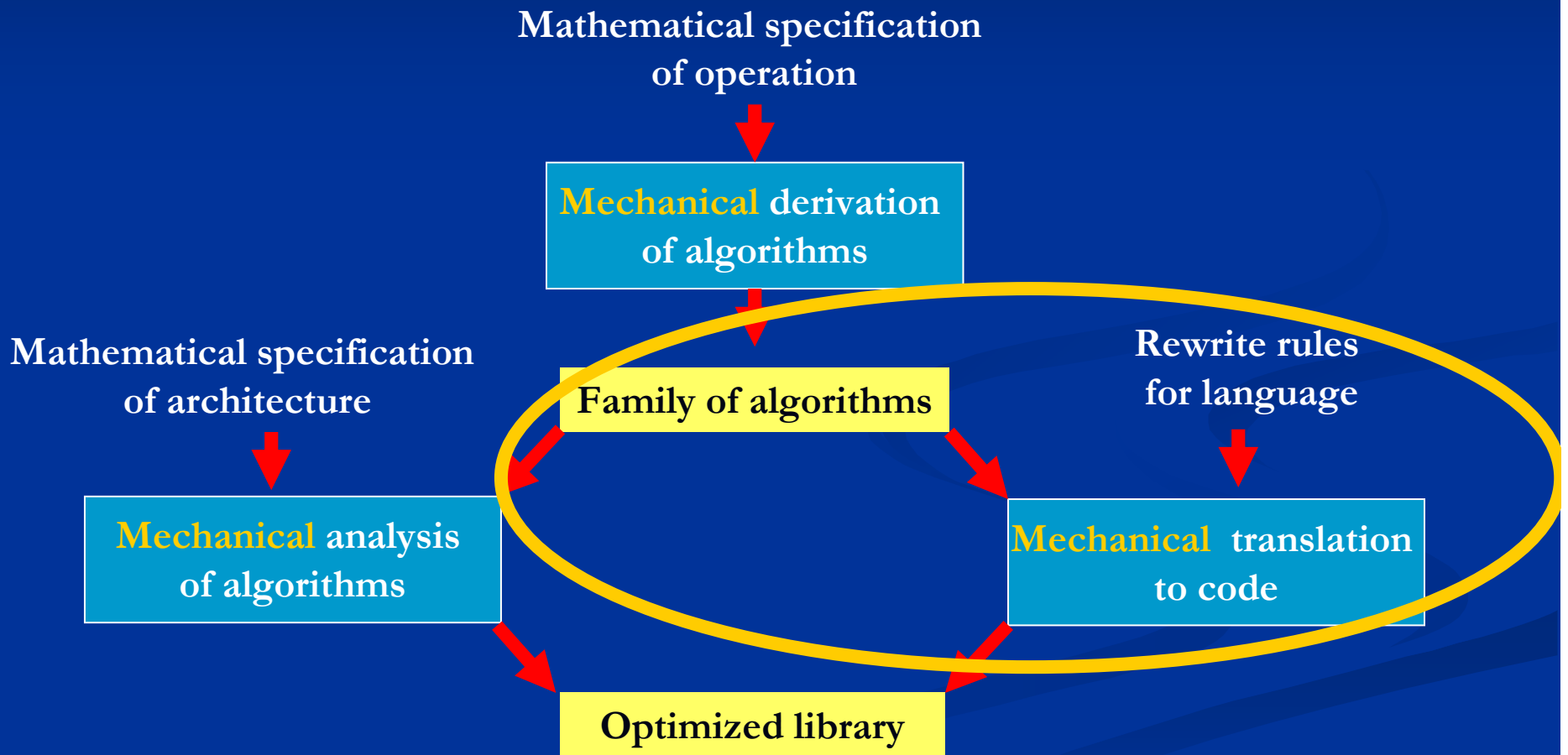
```
    ABL, ABR ] = FLA_Cont_with_3x3_to_2x2( A00, A01, A02, ...  
                                           A10, A11, A12, ...  
                                           A20, A21, A22, 'FLA_TL' );
```

```
end
```


Step 2:

**APIs for Representing
Algorithms in Code**

The Final Generation



LAPACK API

```

DO 10 I = 1, K - NX, NB
    IB = MIN( K-I-1, NB )
*
*      Compute the QR factorization of the current block
*      A(i:m,i:i+ib-1)
*
    CALL DGEQR2( M-I-1, IB, A( I, I ), LDA, TAU( I ), WORK,
$           IINFO )
    IF( I+IB.LE.N ) THEN
*
*      Form the triangular factor of the block reflector
*      H = H(i) H(i+1) . . . H(i+ib-1)
*
    CALL DLARFT( 'Forward', 'Columnwise', M-I-1, IB,
$           A( I, I ), LDA, TAU( I ), WORK, LDWORK )
*
*      Apply H' to A(i:m,i+ib:n) from the left
*
    CALL DLARFB( 'Left', 'Transpose', 'Forward',
$           'Columnwise', M-I-1, N-I-IB-1, IB,
$           A( I, I ), LDA, WORK, LDWORK, A( I, I+IB ),
$           LDA, WORK( IB+1 ), LDWORK )
        END IF
10    CONTINUE

```

```

DO 10 I = 1, K - NX, NB
    IB = MIN( K-I-1, NB )
*
*      Compute the QR factorization of the current block
*      A(i:m,i:i+ib-1)
*
    CALL DGEQR2( M-I-1, IB, A( I, I ), LDA, TAU( I ), WORK,
$              IINFO )
    IF( I+IB.LE.N ) THEN
*
*      Form the triangular factor of the block reflector
*      H = H(i) H(i+1) . . . H(i+ib-1)
*
    CALL DLARFT( 'Forward', 'Columnwise', M-I-1, IB,
$              A( I, I ), LDA, TAU( I ), WORK, LDWORK )
*
*      Apply H' to A(i:m,i+ib:n) from the left
*
    CALL DLARFB( 'Left', 'Transpose', 'Forward',
$              'Columnwise', M-I-1, N-I-IB-1, IB,
$              A( I, I ), LDA, WORK, LDWORK, A( I, I+IB ),
$              LDA, WORK( IB+1 ), LDWORK )
        END IF
10    CONTINUE

```

Warning: I introduced an error!

Step 2:

**New APIs that capture
the algorithm in code**

FLAME@lab

(FLAME/MATLAB API)

```
[ ATL, ATR, ...
  ABL, ABR ] = FLA_Part_2x2( A, 0, 0, 'FLA_TL' );
```

```
while ( size( ATL, 2 ) ~= size( A, 2 ) )
```

```
  b = min( size( ABR, 1 ), nb_alg );
```

```
  [ A00, A01, A02, ...
    A10, A11, A12, ...
    A20, A21, A22 ] = FLA_Repart_2x2_to_3x3( ATL, ATR, ...
                                              ABL, ABR, ...
                                              b, b, 'FLA_BR' );
```

```
  [ U1, s1 ] = QR_unb_var1( [ A11
                             A21 ], s1 );
```

```
  [ A11, ...
    A21 ] = FLA_Part_2x1( U1, b, 'FLA_TOP' );
```

```
  S1 = Accum_S( U1, s1 );
```

```
  U11 = trilu( A11 );
```

```
  U21 = A21;
```

```
  W12 = S1' * ( U11' * A12 + U21' * A22 )
```

```
  A12 = A12 - U11 * W12;
```

```
  A22 = A22 - U21 * W12;
```

Update

$$\begin{pmatrix} A_{12} \\ A_{22} \end{pmatrix} := \left(I + \begin{pmatrix} U_{11} \\ U_{21} \end{pmatrix} S_1 \begin{pmatrix} U_{11} \\ U_{21} \end{pmatrix}^T \right)^T \begin{pmatrix} A_{12} \\ A_{22} \end{pmatrix}$$

```
  [ ATL, ATR, ...
```

```
    ABL, ABR ] = FLA_Cont_with_3x3_to_2x2( A00, A01, A02, ...
                                           A10, A11, A12, ...
                                           A20, A21, A22, 'FLA_TL' );
```

```
end
```



```

[ ATL, ATR, ...
  ABL, ABR ] = FLA_Part_2x2( A, 0, 0, 'FLA_TL' );
[ sT, ...
  sB ] = FLA_Part_2x1( s, 0, 'FLA_TOP' );

while ( size( ATL, 2 ) ~= size( A, 2 ) )
  b = min( size( ABR, 1 ), nb_alg );
  [ A00, A01, A02, ...
    A10, A11, A12, ...
    A20, A21, A22 ] = FLA_Repart_2x2_to_3x3( ATL, ATR, ...
                                             ABL, ABR, b, b, 'FLA_BR' );

  [ s0, ...
    s1, ...
    s2 ] = FLA_Repart_2x1_to_3x1( sT, ...
                                  sB, b, 'FLA_BOTTOM' );

  %-----%
  [ U1, s1 ] = QR_unb_var1( [ A11
                            A21 ], s1 );

  [ A11, ...
    A21 ] = FLA_Part_2x1( U1, b, 'FLA TOP' );
  s1 = Accum_S( U1, s1 );
  % Update rest of matrix
  U11 = trilu( A11 );
  U21 = A21;
  W12 = s1' * ( U11' * A12 + U21' * A22 );
  A12 = A12 - U11 * W12;
  A22 = A22 - U21 * W12;

  %-----%
  [ ATL, ATR, ...
    ABL, ABR ] = FLA_Cont_with_3x3_to_2x2( A00, A01, A02, ...
                                             A10, A11, A12, ...
                                             A20, A21, A22, 'FLA_TL' );

  [ sT, ...
    sB ] = FLA_Cont_with_3x1_to_2x1( s0, ...
                                     s1, ...
                                     s2, 'FLA_TOP' );
end

```

```

[ ATL, ATR, ...
  ABL, ABR ] = FLA_Part_2x2( A, 0, 0, 'FLA_TL' );
[ sT, ...
  sB ] = FLA_Part_2x1( s, 0, 'FLA_TOP' );

while ( size( ATL, 2 ) ~= size( A, 2 ) )
  b = min( size( ABR, 1 ), nb_alg );
  [ A00, A01, A02, ...
    A10, A11, A12, ...
    A20, A21, A22 ] = FLA_Repart_2x2_to_3x3( ATL, ATR, ...
                                             ABL, ABR, b, b, 'FLA_BR' );

  [ s0, ...
    s1, ...
    s2 ] = FLA_Repart_2x1_to_3x1( sT, ...
                                  sB, b, 'FLA_BOTTOM' );

%-----%
  [ U1, s1 ] = QR_unb_var1( [ A11
                            A21 ], s1 );

  [ A11, ...
    A21 ] = FLA_Part_2x1( U1, b, 'FLA_TOP' );
  S1 = Accum_S( U1, s1 );
          % Update rest of matrix
  U11 = trilu( A11 );
  U21 = A21;
  W12 = s1' * ( U11' * A12 + U21' * A22 );
  A12 = A12 - U11 * W12;
  A22 = A22 - U21 * W12;

%-----%
  [ ATL, ATR, ...
    ABL, ABR ] = FLA_Cont_with_3x3_to_2x2( A00, A01, A02, ...
                                             A10, A11, A12, ...
                                             A20, A21, A22, 'FLA_TL' );

  [ sT, ...
    sB ] = FLA_Cont_with_3x1_to_2x1( s0, ...
                                     s1, ...
                                     s2, 'FLA_TOP' );
end

```

```

[ ATL, ATR, ...
  ABL, ABR ] = FLA_Part_2x2( A, 0, 0, 'FLA_TL' );

while ( size( ATL, 2 ) ~= size( A, 2 ) )
  b = min( size( ABR, 1 ), nb_alg );
  [ A00, A01, A02, ...
    A10, A11, A12, ...
    A20, A21, A22 ] = FLA_Repart_2x2_to_3x3( ATL, ATR, ...
                                              ABL, ABR, b, b, 'FLA_BR' );

%-----%
[ U1, s1 ] = QR_unb_var1( [ A11
                          A21 ], s1 );

[ A11, ...
  A21 ] = FLA_Part_2x1( U1, b, 'FLA_TOP' );
S1 = Accum_S( U1, s1 );
U11 = trilu( A11 );
U21 = A21;
W12 = S1' * ( U11' * A12 + U21' * A22 );
A12 = A12 - U11 * W12;
A22 = A22 - U21 * W12;
%-----%

[ ATL, ATR, ...
  ABL, ABR ] = FLA_Cont_with_3x3_to_2x2( A00, A01, A02, ...
                                         A10, A11, A12, ...
                                         A20, A21, A22, 'FLA_TL' );

```

} update

end

```
function [ S ] = Accum_S( U, s )
```

```
U = trilu( U );           % U = lower unit trapezoidal part of U  
s = ones( size( s ) ) ./ s; % Set each element of s to its inverse  
S = inv( triu( U' * U, 1 ) + diag( s ) );  
  
return
```

T. Joffrain, T. M. Low, E. Quintana-Orti, R. van de Geijn, and F. Van Zee, On Accumulating Householder Transformations. *TOMS*, to be revised.

X. Sun. Aggregations of elementary transformations. Tech Report. DUKE-TR-1996-03, **1996**.

C. Puglisi. Modification of the Householder method based on the compact WY representation. *SISC*, 18, 723-726, **1992**

H.F. Walker. Implementation of the GMRES method using Householder transformations. *SISC*, 9, 1, 152-163, **1988**

APIs for C

FLAME/C

QR factorization

```

FLA_Part_2x2( A,      &ATL, &ATR,
              &ABL, &ABR,      0, 0, FLA_TL );

while ( FLA_Obj_width ( ATL ) != FLA_Obj_width ( A ) ){
  b = min( FLA_Obj_width( ABR ), nb_alg );

  FLA_Repart_2x2_to_3x3( ATL, /**/ ATR,      &A00, /**/ &A01, &A02,
                        /* ***** */ /* ***** */
                        &A10, /**/ &A11, &A12,
                        ABL, /**/ ABR,      &A20, /**/ &A21, &A22,
                        b, b, FLA_BR );

  /*-----*/

  FLA_QR_unb( A11,
             A21, s1 );

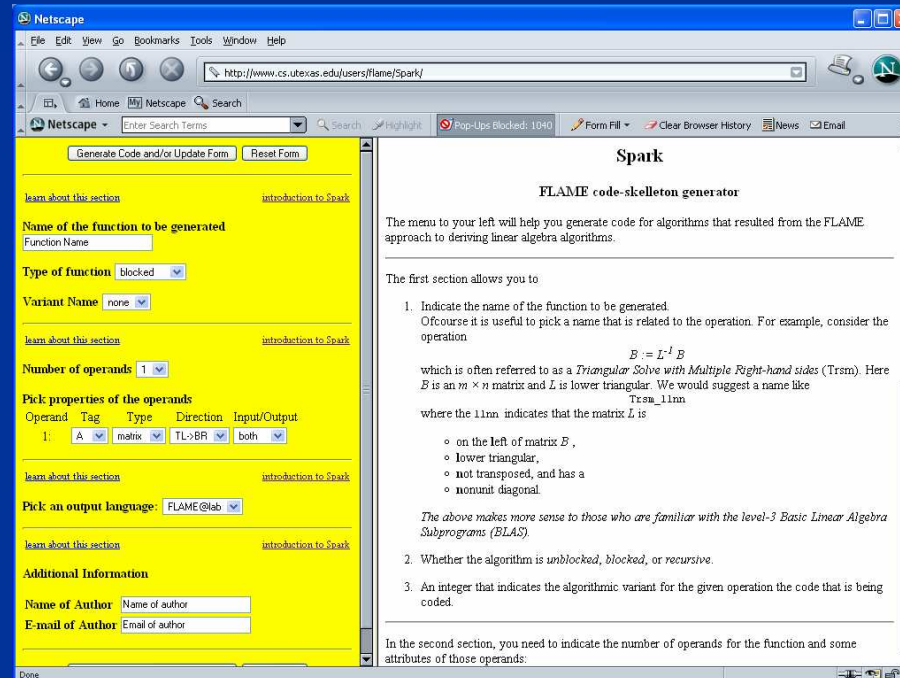
  FLA_Accum_S( A11,
             A21, s1, S1 );

  FLA_Apply_blk_transform( FLA_LEFT, FLA_TRANSPOSE, A11, S1, A12,
                          A21,      A22 );

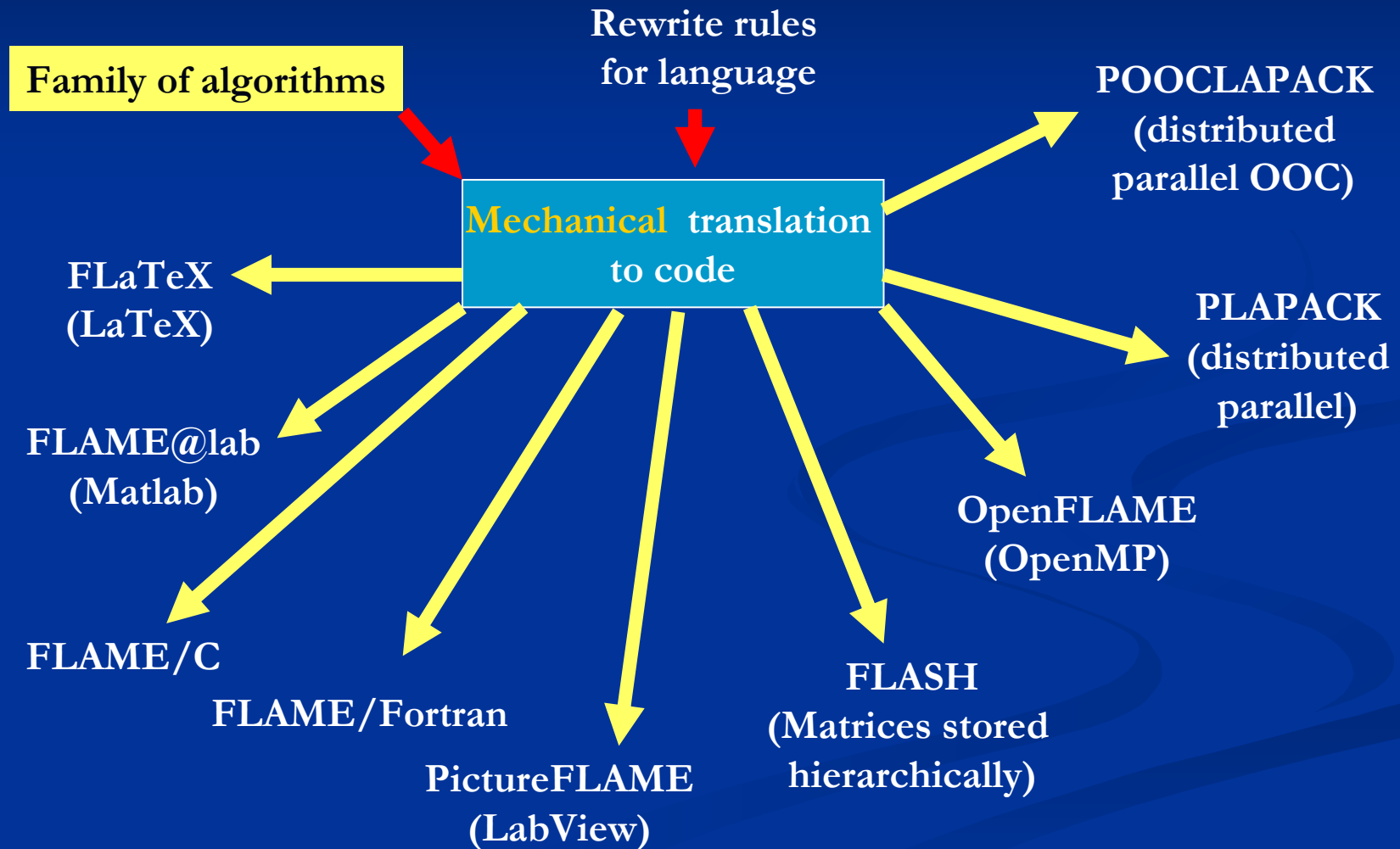
  /*-----*/
  FLA_Cont_with_3x3_to_2x2( &ATL, /**/ &ATR,      A00, A01, /**/ A02,
                            A10, A11, /**/ A12,
                            /* ***** */ /* ***** */
                            &ABL, /**/ &ABR,      A20, A21, /**/ A22,
                            FLA_TL );
}

```

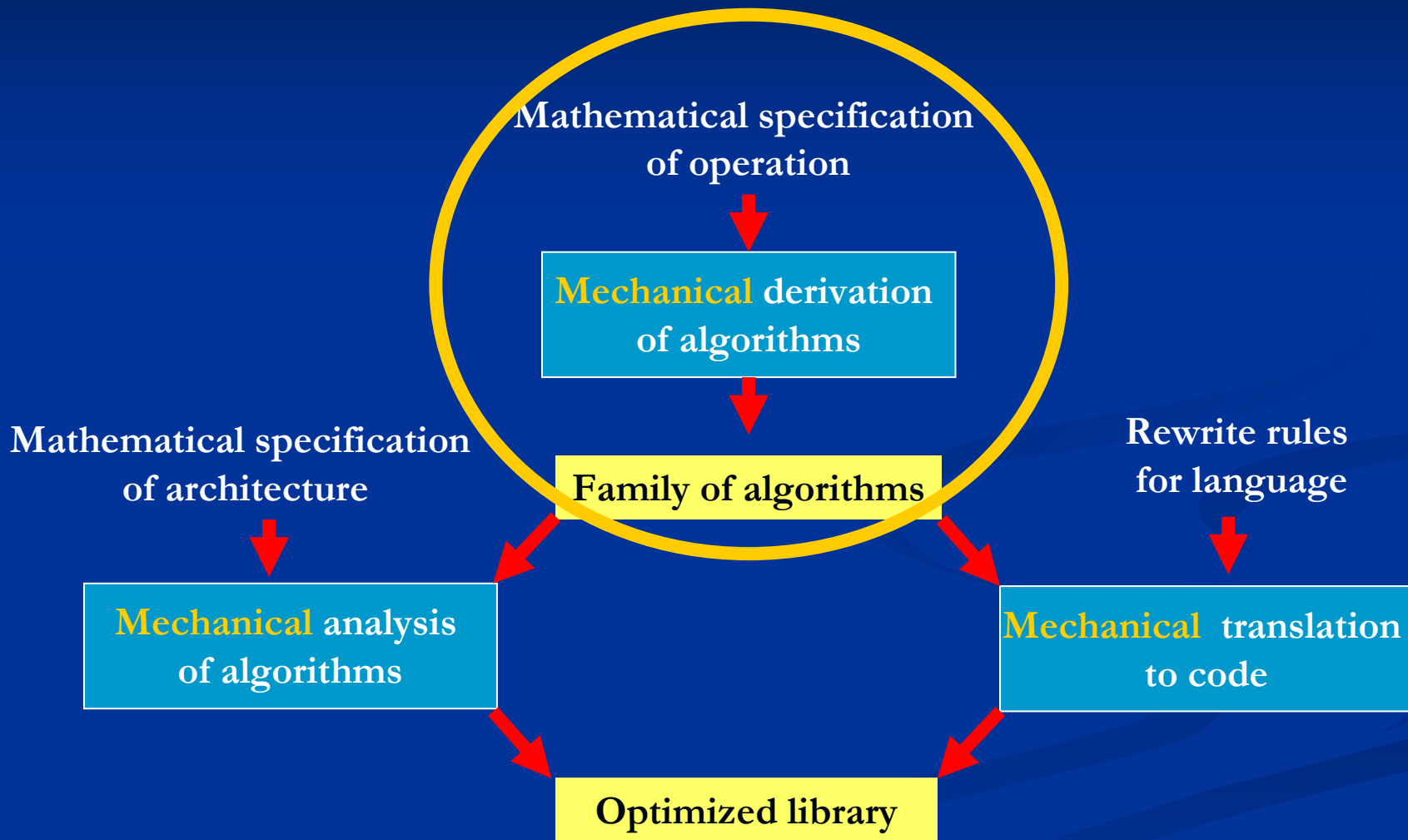
Spark: A Tool for Generating Representations



The Final Generation



The Final Generation



Some Wisdom from the Past

- n The only effective way to raise the confidence level of a program significantly is to give a convincing proof of its correctness. But one should not first make the program and then prove its correctness, because then the requirement of providing the proof would only increase the poor programmer's burden. On the contrary: the programmer should let correctness proof and program grow hand in hand. (E.W. Dijkstra: "The Humble Programmer," 1972 Turing Award lecture, in *ACM Turing Award Lectures: The First Twenty Years, 1966-1985*, ACM Press, New York, 1987.)

What else does the new notation buy us?

- n State of the matrix at the top of the loop

Right-looking algorithm

R_{TL}	R_{TR}
U_{TL}	R_{TR}
U_{BL}	Updated A_{BR}

Left-looking algorithm

R_{TL}	Original
U_{TL}	A_{TR}
U_{BL}	Original A_{BR}

Key insight

- n Given the state that is to be maintained, an algorithm can be systematically derived.
- n The method is sufficiently systematic that it can be made mechanical.

A simpler example: TRSM

n $B := U^{-1} B$ where U is upper triangular

$$\begin{bmatrix} B_T \\ \hline B_B \end{bmatrix} := \begin{bmatrix} U_{TL}^{-1} (B_T - U_{TR} U_{BR}^{-1} B_B) \\ \hline U_{BR}^{-1} B_B \end{bmatrix}$$

Mathematical specification
of operation



Mechanical derivation
of algorithms



Family of algorithms

$$U_{TL}^{-1} (B_T - U_{TR} U_{BR}^{-1} B_B)$$

$$U_{BR}^{-1} B_B$$

MATHEMATICA[®]5

Family of algorithms

Mechanical Derivation

Mathematical specification
of operation



Mechanical derivation
of algorithms



MATHEMATICA⁵
Family of algorithms

Switch to Demo

`Notation[$\Xi \left(\begin{array}{c} \underline{A}_-, \underline{B}_-, \underline{C}_- \\ \underline{D}_-, \underline{E}_-, \underline{F}_- \end{array} \right) \Rightarrow \text{coupledSylv}[\underline{A}_-, \underline{B}_-, \underline{C}_-, \underline{D}_-, \underline{E}_-, \underline{F}_-]$`

PMEs

1x2

2x1

2x2

```
worksheet[coupledSylv,  
{ {"A", "UpperTriangular", "TL"}, {"B", "UpperTriangular", "BR"}, {"C", "Overwrite", "TR"},  
  {"D", "UpperTriangular", "TL"}, {"E", "UpperTriangular", "BR"}, {"F", "TR", "Overwrite"} }]
```

2x1

2x2

PME 2x2

Loop Inv 1

$$\left(\begin{array}{c|c} C_{TL} & C_{TR} \\ \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket & C_{BR} \end{array} \right)$$

$$\left(\begin{array}{c|c} F_{TL} & F_{TR} \\ \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket & F_{BR} \end{array} \right)$$



State at top of loop

Loop Inv 2

$$\left(\begin{array}{c|c} -A_{TR} \cdot \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + C_{TL} & C_{TR} \\ \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket & C_{BR} \end{array} \right)$$

$$\left(\begin{array}{c|c} F_{TL} & F_{TR} \\ \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket & F_{BR} \end{array} \right)$$

Loop Inv 4

$$\begin{pmatrix} -A_{TR} \cdot \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + C_{TL} & C_{TR} \\ \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket & C_{BR} \\ -D_{TR} \cdot \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + F_{TL} & F_{TR} \\ \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket & F_{BR} \end{pmatrix}$$

```

In[156]:= coupledSylv4[
  {{aTL_, aTR_},
   {0, aBR_}},
  {{bTL_, bTR_},
   {0, bBR_}},
  {{cTL_, cTR_},
   {cBL_, cBR_}},

  {{dTL_, dTR_},
   {0, dBR_}},
  {{eTL_, eTR_},
   {0, eBR_}},
  {{fTL_, fTR_},
   {fBL_, fBR_}}] :=

Module[{BL, BL1, BL2},
  BL = coupledSylv[aBR, bTL, cBL, dBR, eTL, fBL];
  {BL1, BL2} = assignParts[BL, 2];

  {
    {{cTL - prod[aTR, BL1], cTR},
     {BL1, cBR}},
  }

```



**Mathematic specification
of state at top of loop**

```
{cBL_, cBR_}},

{{dTL_, dTR_},
 {0, dBR_}},
{{eTL_, eTR_},
 {0, eBR_}},
{{fTL_, fTR_},
 {fBL_, fBR_}}] :=
```

```
Module[{BL, BL1, BL2},
 BL = coupledSylv[aBR, bTL, cBL, dBR, eTL, fBL];
 {BL1, BL2} = assignParts[BL, 2];
```


```
{
 {{cTL - prod[aTR, BL1], cTR},
 {BL1, cBR}},

 {{fTL - prod[dTR, BL1], fTR},
 {BL2, fBR}}
}
]
```

```
In[177]:= Map[myMatrixForm, coupledSylv4[mA, mB, mC, mD, mE, mF]] // ColumnForm;
```

```
worksheet[coupledSylv4,
 {"A", "UpperTriangular", "TL"}, {"B", "UpperTriangular", "BR"}, {"C", "Overwrite", "TR"},
 {"D", "UpperTriangular", "TL"}, {"E", "UpperTriangular", "BR"}, {"F", "TR", "Overwrite"}]
```

Function that generates algorithm



Loop Inv 5

$$\begin{pmatrix} \left(\sum \left(\frac{A_{TL}, B_{TL}, -A_{TR} \cdot \left(\sum \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + C_{TL}}{D_{TL}, E_{TL}, -D_{TR} \cdot \left(\sum \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + F_{TL}} \right) \right) \llbracket 1 \rrbracket & C_{TR} \\ \left(\sum \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket & C_{BR} \end{pmatrix}$$

Operation: [C F]=coupledSylv4(A B C D E F)

Partition

$$A \rightarrow \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \quad B \rightarrow \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix} \quad C \rightarrow \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} \quad D \rightarrow \begin{pmatrix} D_{11} & D_{12} \\ D_{21} & D_{22} \end{pmatrix} \quad E \rightarrow \begin{pmatrix} E_{11} & E_{12} \\ E_{21} & E_{22} \end{pmatrix} \quad F \rightarrow \begin{pmatrix} F_{11} & F_{12} \\ F_{21} & F_{22} \end{pmatrix}$$

Loop Invariant (c):

$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} -A_{11} \cdot \left(\begin{matrix} \left(\frac{A_{22} \cdot F_{11} \cdot C_{11}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{11} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{12}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{21}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{21} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{22}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \end{matrix} \right) \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{11}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{11} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{12}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{21}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{21} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{22}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \end{pmatrix}$$

Matrix C

Separation

$$\left[\begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \rightarrow \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix}, \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix} \rightarrow \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}, \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} \rightarrow \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix}, \begin{pmatrix} D_{11} & D_{12} \\ D_{21} & D_{22} \end{pmatrix} \rightarrow \begin{pmatrix} D_{11} & D_{12} \\ D_{21} & D_{22} \end{pmatrix}, \begin{pmatrix} E_{11} & E_{12} \\ E_{21} & E_{22} \end{pmatrix} \rightarrow \begin{pmatrix} E_{11} & E_{12} \\ E_{21} & E_{22} \end{pmatrix}, \begin{pmatrix} F_{11} & F_{12} \\ F_{21} & F_{22} \end{pmatrix} \rightarrow \begin{pmatrix} F_{11} & F_{12} \\ F_{21} & F_{22} \end{pmatrix} \right]$$

Loop Invariant (c) before the update:

$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} -A_{11} \cdot \left(\begin{matrix} \left(\frac{A_{22} \cdot F_{11} \cdot C_{11}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{11} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{12}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{21}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{21} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{22}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \end{matrix} \right) \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{11}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{11} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{12}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{21}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{21} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{22}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \end{pmatrix}$$

Algorithm with intermediate states generated by system

```

C11 := -A11 * C10 + C00
C12 := -A11 * C11 - A12 * C21 + C01
C21 := (A11 * F11 * C10) / (B22 * F11 * F21)
C22 := (A11 * F11 * F10 * F12 * A12 * C21 + C11) / (B22 * F11 * F21)
C21 := (A12 * F11 * F10 * F12 * C21) / (B22 * F11 * F21)
F11 := -B11 * C10 + F00
F12 := -B11 * C11 - B12 * C21 + F01
F21 := (B11 * F11 * C10) / (B22 * F11 * F21)
F22 := (B11 * F11 * F10 * F12 * B12 * C21) / (B22 * F11 * F21)
F21 := (B12 * F11 * F10 * F12 * F21) / (B22 * F11 * F21)
    
```

Continue with

$$\left[\begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \rightarrow \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix}, \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix} \rightarrow \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}, \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} \rightarrow \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix}, \begin{pmatrix} D_{11} & D_{12} \\ D_{21} & D_{22} \end{pmatrix} \rightarrow \begin{pmatrix} D_{11} & D_{12} \\ D_{21} & D_{22} \end{pmatrix}, \begin{pmatrix} E_{11} & E_{12} \\ E_{21} & E_{22} \end{pmatrix} \rightarrow \begin{pmatrix} E_{11} & E_{12} \\ E_{21} & E_{22} \end{pmatrix}, \begin{pmatrix} F_{11} & F_{12} \\ F_{21} & F_{22} \end{pmatrix} \rightarrow \begin{pmatrix} F_{11} & F_{12} \\ F_{21} & F_{22} \end{pmatrix} \right]$$

Loop Invariant (c) after the update:

$$-A_{11} \cdot \left(\begin{matrix} \left(\frac{A_{22} \cdot F_{11} \cdot C_{11}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{11} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{12}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{21}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{21} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{22}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \end{matrix} \right) \quad D_{10} + A_{12} \cdot \left(\begin{matrix} \left(\frac{A_{22} \cdot F_{11} \cdot C_{11}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{11} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{12}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{21}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} + C_{21} \\ \left(\frac{A_{22} \cdot F_{11} \cdot C_{22}}{B_{22} \cdot F_{11} \cdot F_{21}} \right) D_{10} \end{matrix} \right)$$

$$\left(\begin{array}{c|cc} F_{20} & (F_{21} & F_{22}) \end{array} \right) \quad \left(\Xi \left(\frac{A_{22}, B_{00}, \hat{C}_{20}}{D_{22}, E_{00}, \hat{F}_{20}} \right) \right) [[2]] \quad \left(\hat{F}_{21} \right)$$

$$\begin{aligned} C_{00} & := -A_{01} \cdot C_{10} + C_{00} \\ C_{01} & := -A_{01} \cdot C_{11} - A_{02} \cdot C_{21} + C_{01} \\ C_{10} & := \left(\Xi \left(\frac{A_{11}, B_{00}, C_{10}}{D_{11}, E_{00}, F_{10}} \right) \right) [[1]] \\ C_{11} & := \left(\Xi \left(\frac{A_{11}, B_{11}, -F_{10} \cdot B_{01} - A_{12} \cdot C_{21} + C_{11}}{D_{11}, E_{11}, -F_{10} \cdot E_{01} - D_{12} \cdot C_{21} + F_{11}} \right) \right) [[1]] \\ C_{21} & := \left(\Xi \left(\frac{A_{22}, B_{11}, -F_{20} \cdot B_{01} + C_{21}}{D_{22}, E_{11}, -F_{20} \cdot E_{01} + F_{21}} \right) \right) [[1]] \\ F_{00} & := -D_{01} \cdot C_{10} + F_{00} \\ F_{01} & := -D_{01} \cdot C_{11} - D_{02} \cdot C_{21} + F_{01} \\ F_{10} & := \left(\Xi \left(\frac{A_{11}, B_{00}, C_{10}}{D_{11}, E_{00}, F_{10}} \right) \right) [[2]] \\ F_{11} & := \left(\Xi \left(\frac{A_{11}, B_{11}, -F_{10} \cdot B_{01} - A_{12} \cdot C_{21} + C_{11}}{D_{11}, E_{11}, -F_{10} \cdot E_{01} - D_{12} \cdot C_{21} + F_{11}} \right) \right) [[2]] \\ F_{21} & := \left(\Xi \left(\frac{A_{22}, B_{11}, -F_{20} \cdot B_{01} + C_{21}}{D_{22}, E_{11}, -F_{20} \cdot E_{01} + F_{21}} \right) \right) [[2]] \end{aligned}$$

Continue with

Update to be performed

$$\left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket$$

$$\left(\Xi \left(\frac{A_{BR}, B_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot B_{TR} + C_{BR}}{D_{BR}, E_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot E_{TR} + F_{BR}} \right) \right) \llbracket 2 \rrbracket$$

Loop Inv 49

$$\left(\Xi \left(\frac{A_{TL}, B_{TL}, -A_{TR} \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + C_{TL}}{D_{TL}, E_{TL}, -D_{TR} \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + F_{TL}} \right) \right) \llbracket 1 \rrbracket - \left(\Xi \left(\frac{A_{TL}, B_{TL}, -A_{TR} \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + C_{TL}}{D_{TL}, E_{TL}, -D_{TR} \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + F_{TL}} \right) \right) \llbracket 2 \rrbracket \cdot B_{TR} - A_{TR} \cdot \left(\Xi \left(\frac{A_{BR}, B_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot B_{TR} + C_{BR}}{D_{BR}, E_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot E_{TR} + F_{BR}} \right) \right) \llbracket 1 \rrbracket + \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket \left(\Xi \left(\frac{A_{BR}, B_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot B_{TR} + C_{BR}}{D_{BR}, E_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot E_{TR} + F_{BR}} \right) \right) \llbracket 1 \rrbracket$$

$$\left(\Xi \left(\frac{A_{TL}, B_{TL}, -A_{TR} \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + C_{TL}}{D_{TL}, E_{TL}, -D_{TR} \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + F_{TL}} \right) \right) \llbracket 2 \rrbracket - \left(\Xi \left(\frac{A_{TL}, B_{TL}, -A_{TR} \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + C_{TL}}{D_{TL}, E_{TL}, -D_{TR} \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 1 \rrbracket + F_{TL}} \right) \right) \llbracket 2 \rrbracket \cdot E_{TR} - D_{TR} \cdot \left(\Xi \left(\frac{A_{BR}, B_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot B_{TR} + C_{BR}}{D_{BR}, E_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot E_{TR} + F_{BR}} \right) \right) \llbracket 1 \rrbracket + \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \left(\Xi \left(\frac{A_{BR}, B_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot B_{TR} + C_{BR}}{D_{BR}, E_{BR}, - \left(\Xi \left(\frac{A_{BR}, B_{TL}, C_{BL}}{D_{BR}, E_{TL}, F_{BL}} \right) \right) \llbracket 2 \rrbracket \cdot E_{TR} + F_{BR}} \right) \right) \llbracket 2 \rrbracket$$

worksheet [coupledSylv,

- { "A", "UpperTriangular", "TL" }, { "B", "UpperTriangular", "BR" }, { "C", "Overwrite", "TR" },
- { "D", "UpperTriangular", "TL" }, { "E", "UpperTriangular", "BR" }, { "F", "TR", "Overwrite" } }

Scope of Methodology

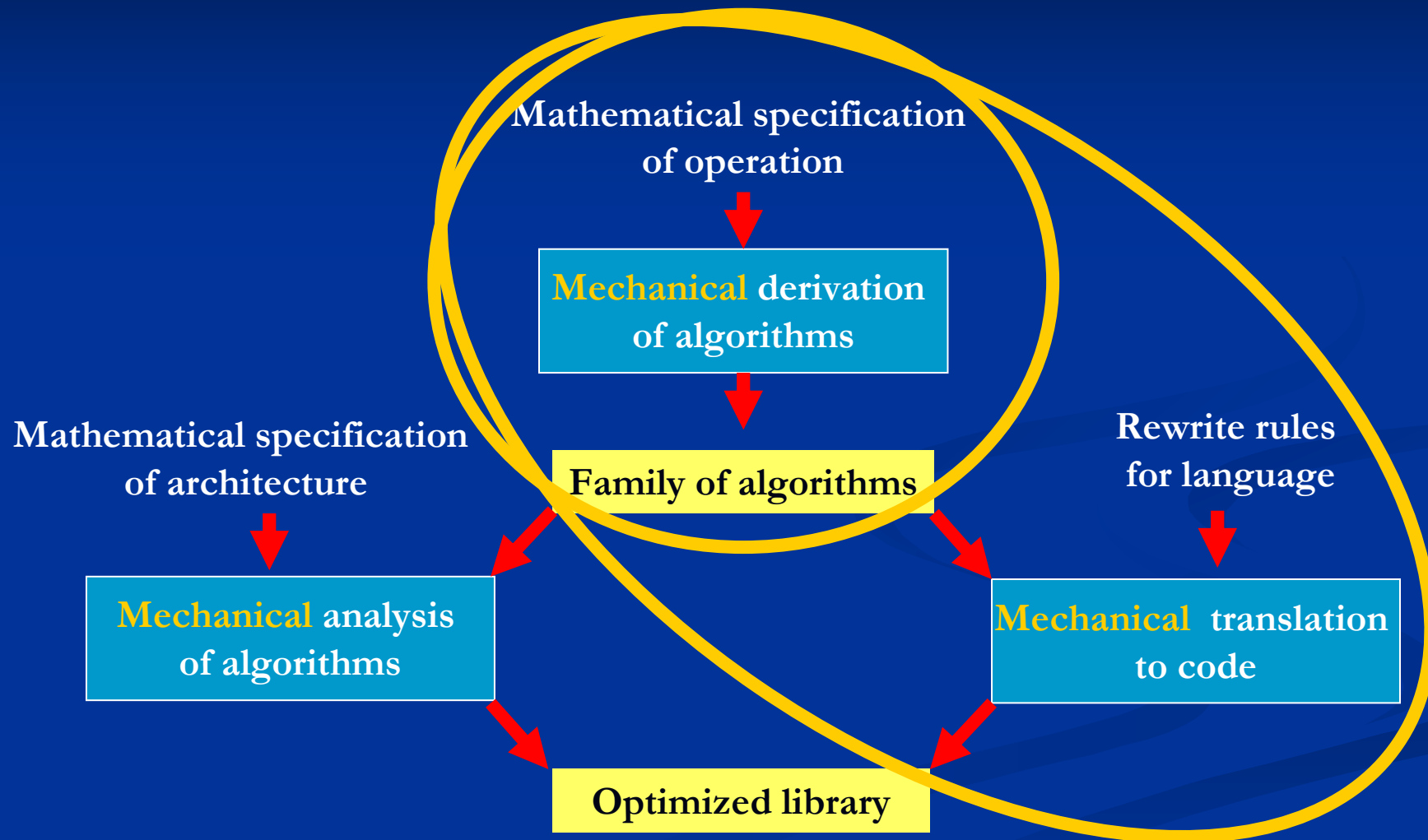
- n Triangular Generalized Sylvester Equation

$$A X + Y B = C$$

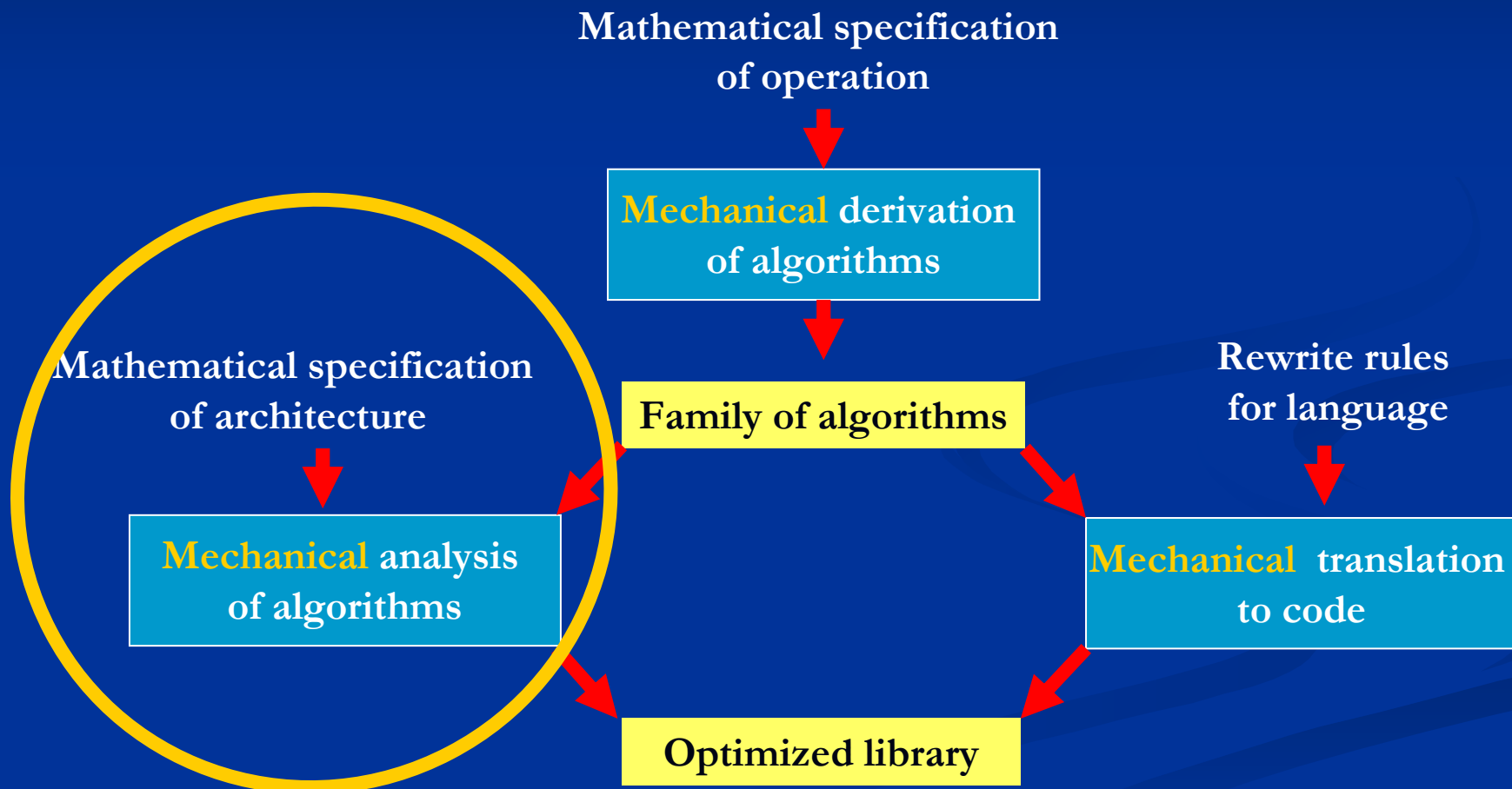
$$D X + Y E = F$$

- n A, B, D, E triangular
- n X and Y overwrite C and F
- n 57 algorithmic variants...
- n Blocked and unblocked each
- n Compose recursively for optimal performance...

The Final Generation



The Final Generation



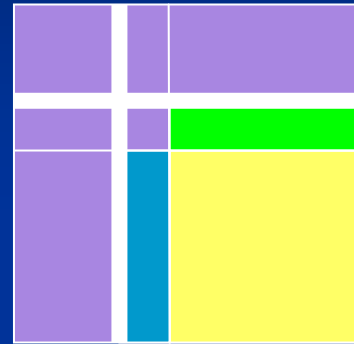
Mechanical Cost Analysis

Dissertation of
John Gunnels
(for parallel distributed)

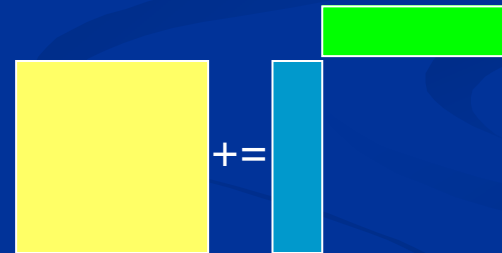
Mechanical Cost Analysis

Key: Vertical Integration

n Solvers (FLAME)



n BLAS (FLAME)



n Low level kernels (K.Goto)



Systematic
~~Mechanical~~ Stability
Analysis

Paolo Bientinesi
(in progress)

Annotated Algorithm: $\kappa := x^T y$ $\tilde{\kappa} = x^T \Delta y$ Step

Partition

$$x \rightarrow \left(\begin{array}{c} x_T \\ x_B \end{array} \right), y \rightarrow \left(\begin{array}{c} y_T \\ y_B \end{array} \right) \quad \Delta \rightarrow \left(\begin{array}{c|c} \Delta_T & 0 \\ \hline 0 & \Delta_B \end{array} \right)$$

where ...

$$\left\{ \tilde{\kappa} = \left[x_T^T y_T \right] \right\} \quad \left\{ \tilde{\kappa} = x_T^T \Delta_T y_T \right\}$$

while $m(x_B) > 0$ do $\{k = m(x_T)\}$

Repartition

$$\left(\begin{array}{c} x_T \\ x_B \end{array} \right) \rightarrow \left(\begin{array}{c} x_0 \\ \chi_1 \\ x_2 \end{array} \right), \left(\begin{array}{c} y_T \\ y_B \end{array} \right) \rightarrow \left(\begin{array}{c} y_0 \\ \psi_1 \\ y_2 \end{array} \right) \quad \left(\begin{array}{c|c} \Delta_T & 0 \\ \hline 0 & \Delta_B \end{array} \right) \rightarrow \left(\begin{array}{c|c|c} \Delta_0 & 0 & 0 \\ \hline 0 & \delta_1 & 0 \\ \hline 0 & 0 & \Delta_2 \end{array} \right)$$

where ...

$$\left\{ \tilde{\kappa} = \left[x_0^T y_0 \right] \right\} \quad \left\{ \tilde{\kappa} = x_0^T \Delta_0^{\{k\}} y_0 \right\}$$

$$\tilde{\kappa} := [\tilde{\kappa} + \chi_1 \psi_1]$$

$$\tilde{\kappa} := \left(\tilde{\kappa} + \chi_1 \psi_1 (1 + \epsilon_*) \right) (1 + \epsilon_+)$$

$$= x_0^T \Delta_0^{\{k\}} (1 + \epsilon_+) y_0 + \chi_1 (1 + \epsilon_*) (1 + \epsilon_+) \psi_1$$

$$\Delta_0 := \Delta_0 (1 + \epsilon_+)$$

$$\delta_1 := (1 + \epsilon_+) (1 + \epsilon_*)$$

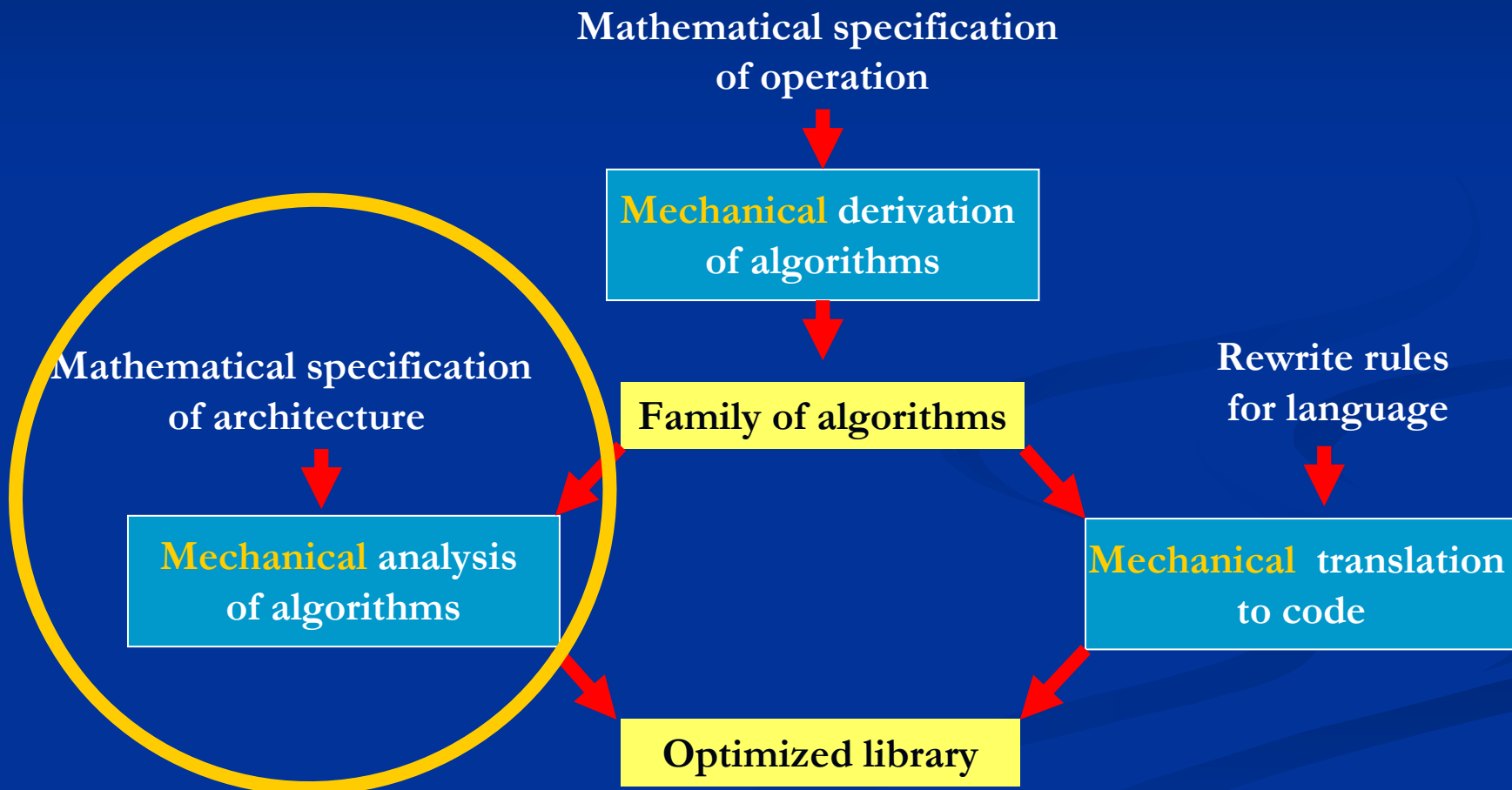
$$\left\{ \tilde{\kappa} = \left[x_0^T y_0 + \chi_1 \psi_1 \right] \right\} \quad \left\{ \tilde{\kappa} = \left(\begin{array}{c} x_0 \\ \chi_1 \end{array} \right)^T \left[\begin{array}{c|c} \Delta_0^{\{k\}} & \\ \hline & \delta_1 \end{array} \right] \left(\begin{array}{c} y_0 \\ \psi_1 \end{array} \right) = \right. \\ \left. = x_0^T \Delta_0^{\{k\}} y_0 + \chi_1 \delta_1 \psi_1 \right\}$$

Continue with ...

$$\left(\begin{array}{c|c} \Delta_T & 0 \\ \hline 0 & \Delta_B \end{array} \right) \leftarrow \left(\begin{array}{c|c|c} \Delta_0 & 0 & 0 \\ \hline 0 & \delta_1 & 0 \\ \hline 0 & 0 & \Delta_2 \end{array} \right)$$

enddo 5b

The Final Generation

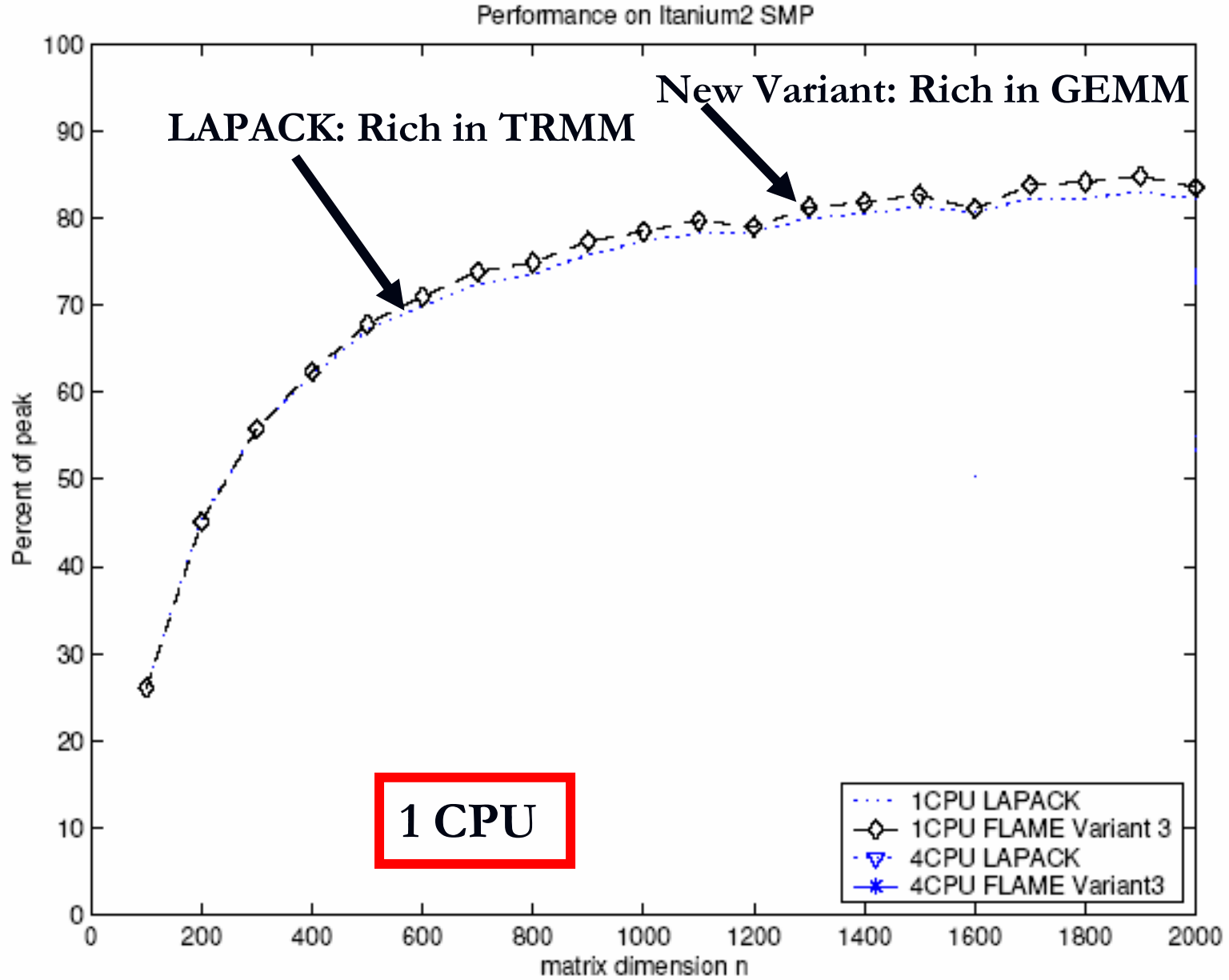


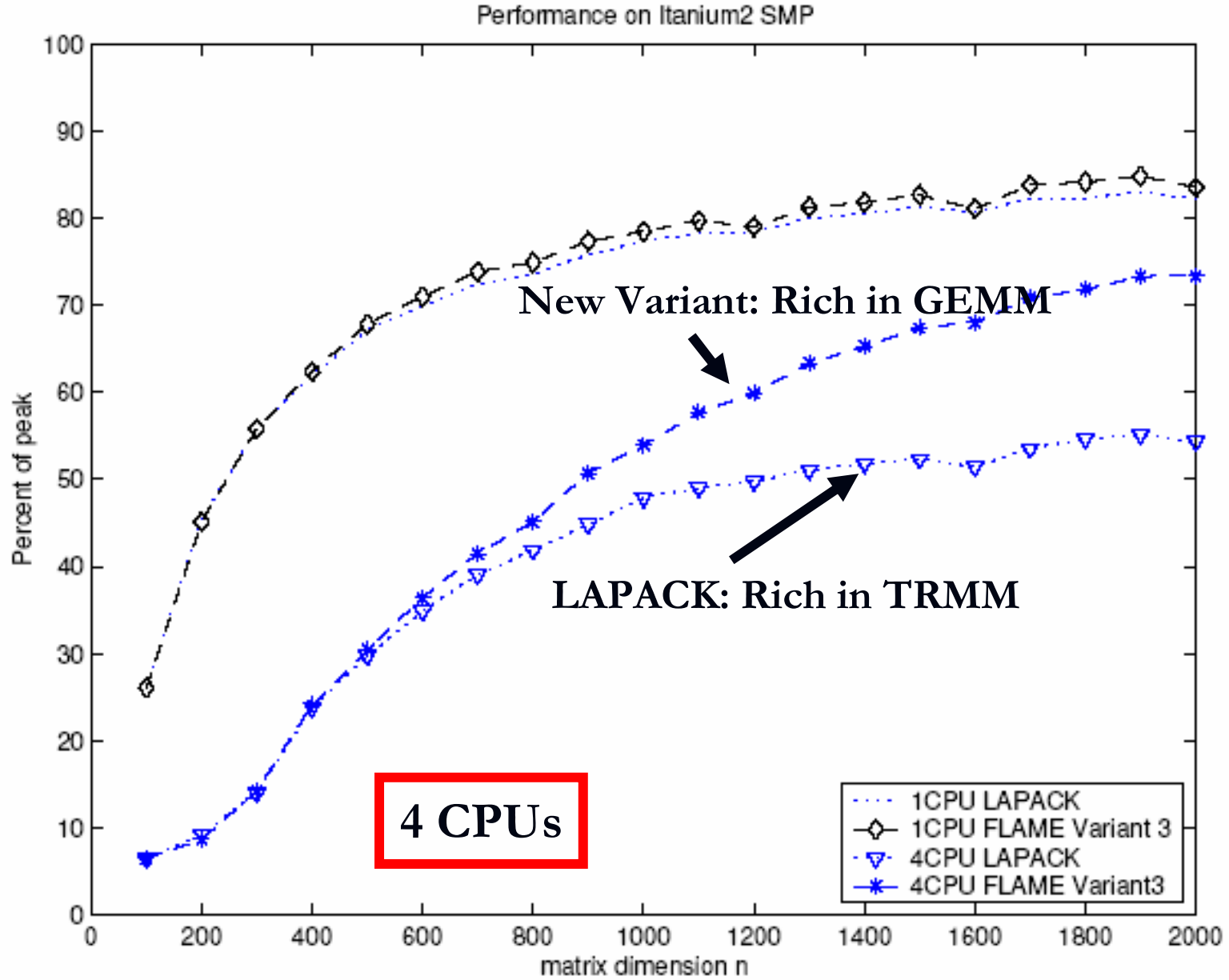
Future Challenges: Multicore Processors

- n Processors with two cores available now
- n All HP users will have to cope
- n 32-128 cores per processor in 10 years?
- n 32 processor SMP x 32 cores = 1024 way SMP parallelism
- n Virtually no literature on SCALABLE linear algebra libraries for SMPs

Shared Memory Parallelism

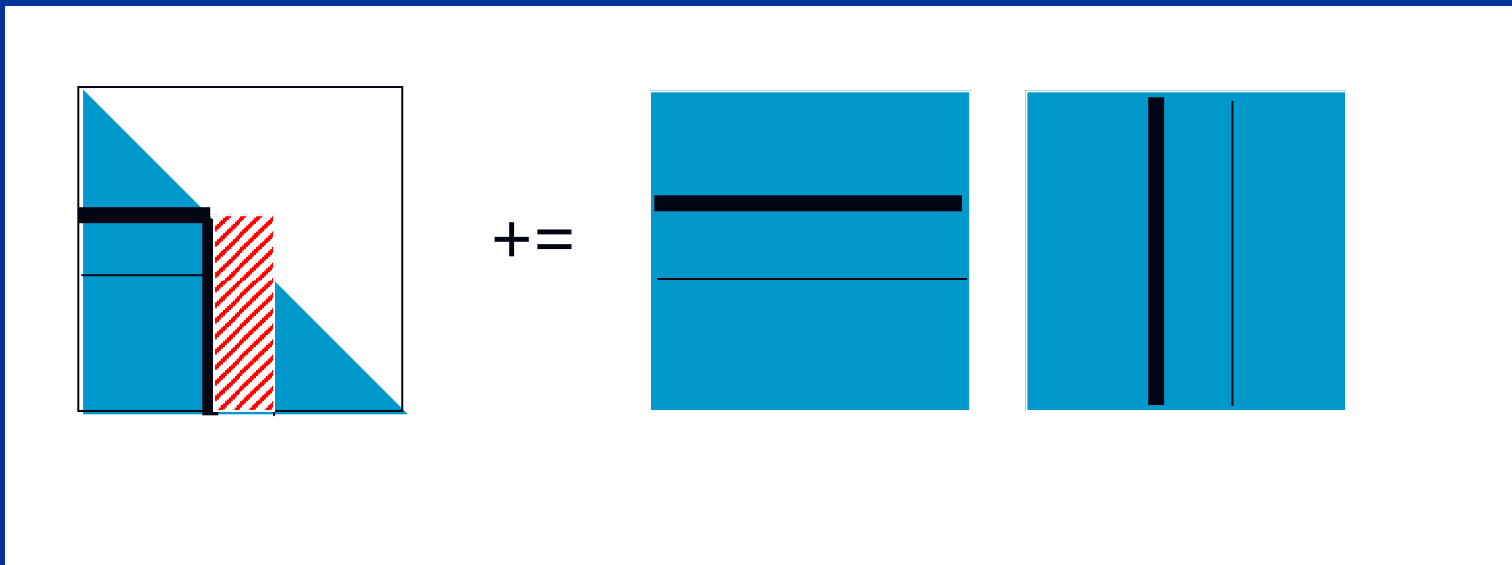
- n Traditional approach:
 - n Only from multithreaded BLAS
- n Observation:
 - n Better speedup if parallelism is exposed at a higher level
 - n Scalability requires 2D work distribution
- n FLAME approach:
 - n Choose the best algorithmic variant
 - n Work queuing (e.g., OpenMP task queues)





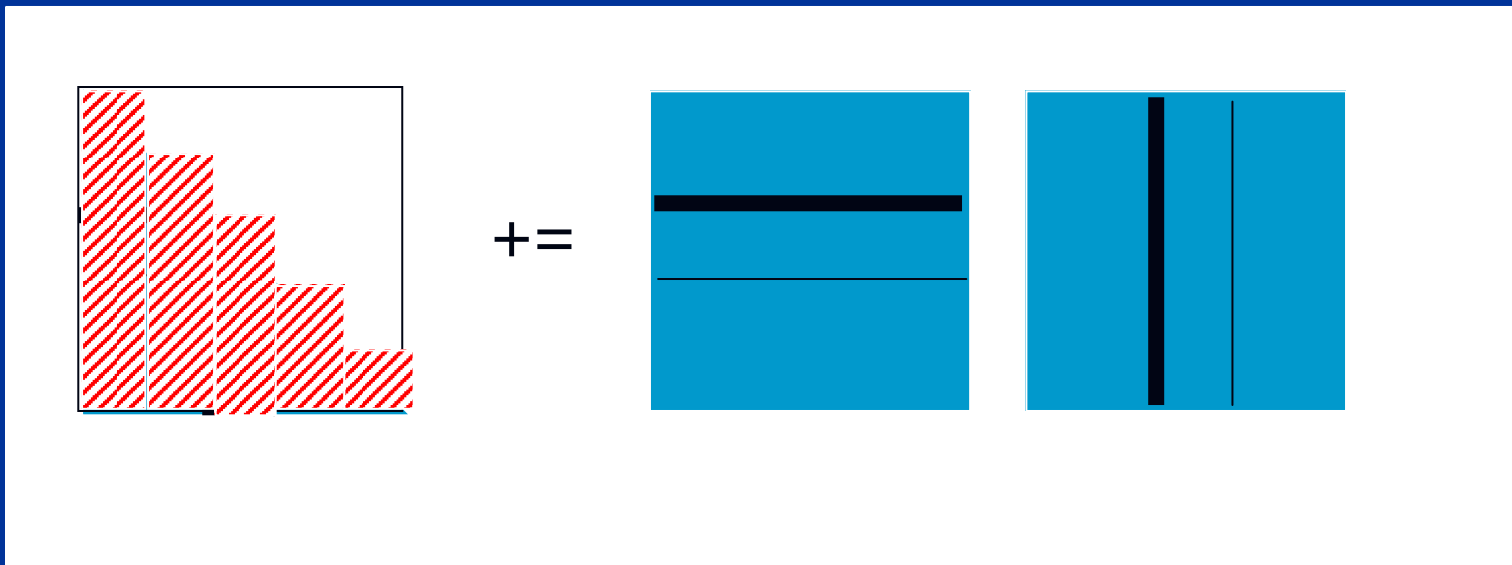
Work queuing

- n Simple example:
 - n Symmetric rank-k update



Work queuing

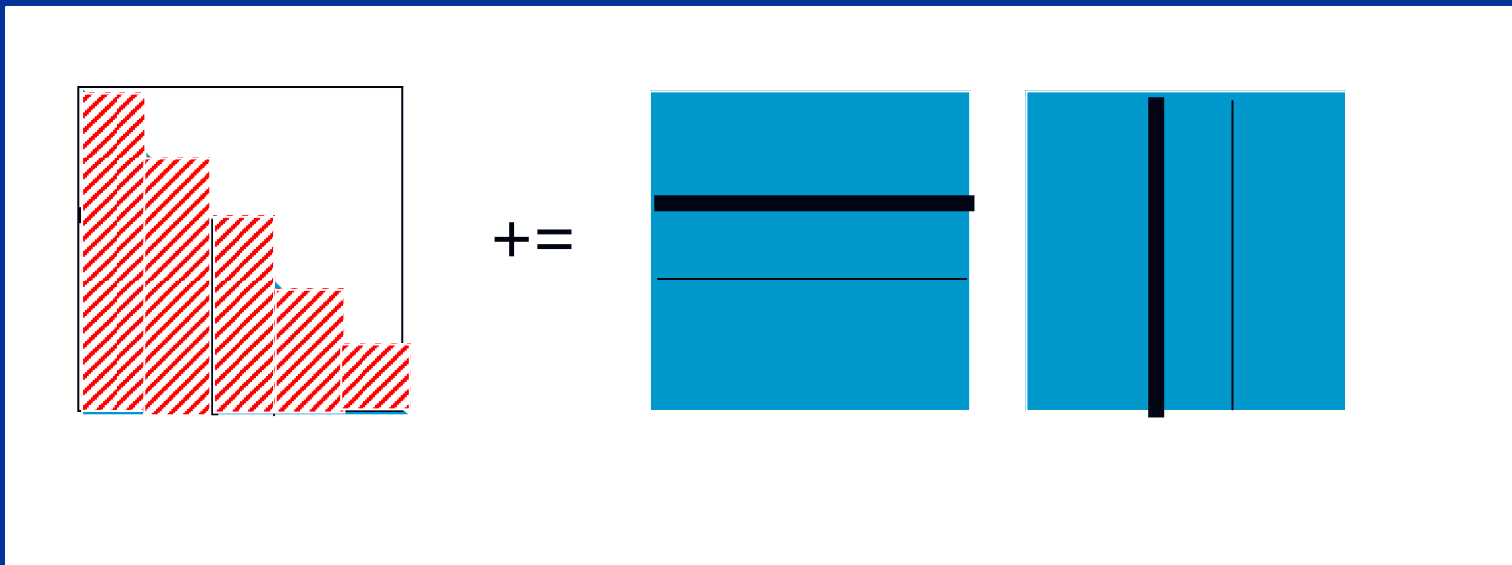
- n Simple example:
 - n Symmetric rank-k update



Shared Memory Parallelism

n Example:

n Symmetric rank-k update




```

while ( FLA_Obj_length( CTL ) < FLA_Obj_length( C ) ){
  b = min( FLA_Obj_length( CBR ), nb_alg );

  FLA_Repart_2x2_to_3x3( CTL, /**/ CTR,   &C00, /**/ &C01, &C02,
                        /***/ /***/
                        &C10, /**/ &C11, &C12,
                        CBL, /**/ CBR,   &C20, /**/ &C21, &C22,
                        b, b, FLA_BR );
  FLA_Repart_2x1_to_3x1( AT,           &A0,
                        /* ** */      /* ** */
                        &A1,
                        AB,           &A2,   b, FLA_BOTTOM );
  /*-----*/

  FLA_Gemm( FLA_NO_TRANSPOSE, FLA_TRANSPOSE, ONE, A0, A1, ONE, C10 );
  FLA_Syrk( FLA_LOWER_TRIANGULAR, FLA_NO_TRANSPOSE, ONE, A1, ONE, C11 );

  /*-----*/
  FLA_Cont_with_3x3_to_2x2( &CTL, /**/ &CTR, C00, C01, /**/ C02,
                           C10, C11, /**/ C12,
                           /***/ /***/
                           &CBL, /**/ &CBR, C20, C21, /**/ C22,
                           FLA_TL );
  FLA_Cont_with_3x1_to_2x1( &AT,           A0,
                           A1,
                           /* ** */      /* ** */
                           &AB,           A2,   FLA_TOP );
}

```

```

while ( FLA_Obj_length( CTL ) < FLA_Obj_length( C ) ){
  b = min( FLA_Obj_length( CBR ), nb_alg );

  FLA_Repart_2x2_to_3x3( CTL, /**/ CTR,   &C00, /**/ &C01, &C02,
                        /***/ /***/
                        &C10, /**/ &C11, &C12,
                        CBL, /**/ CBR,   &C20, /**/ &C21, &C22,
                        b, b, FLA_BR );
  FLA_Repart_2x1_to_3x1( AT,           &A0,
                        /* ** */       /* ** */
                        AB,             &A1,
                                       &A2,   b, FLA_BOTTOM );
  /*-----*/
  FLA_Gemm( FLA_NO_TRANSPOSE, FLA_TRANSPOSE, ONE, A0, A1, ONE, C10 );
  FLA_Syrk( FLA_LOWER_TRIANGULAR, FLA_NO_TRANSPOSE, ONE, A1, ONE, C11 );
  /*-----*/
  FLA_Cont_with_3x3_to_2x2( &CTL, /**/ &CTR, C00, C01, /**/ C02,
                           C10, C11, /**/ C12,
                           /***/ /***/
                           &CBL, /**/ &CBR, C20, C21, /**/ C22,
                           FLA_TL );
  FLA_Cont_with_3x1_to_2x1( &AT,           A0,
                           /* ** */       A1,
                           &AB,           /* ** */
                                       A2,   FLA_TOP );
}

```

```

#pragma intel omp parallel taskq
{
while ( FLA_Obj_length( CTL ) < FLA_Obj_length( C ) ){
  b = min( FLA_Obj_length( CBR ), nb_alg );

  FLA_Repart_2x2_to_3x3( CTL, /**/ CTR,   &C00, /**/ &C01, &C02,
                        /***/ /***/
                        &C10, /**/ &C11, &C12,
                        CBL, /**/ CBR,   &C20, /**/ &C21, &C22,
                        b, b, FLA_BR );

  FLA_Repart_2x1_to_3x1( AT,           &A0,
                        /* ** */      /* ** */
                        AB,           &A1,
                                       &A2,   b, FLA_BOTTOM );
  /*-----*/
#pragma intel omp task captureprivate( A0, A1, C10, C11 )
{
  FLA_Gemm( FLA_NO_TRANSPOSE, FLA_TRANSPOSE, ONE, A0, A1, ONE, C10 );
  FLA_Syrk( FLA_LOWER_TRIANGULAR, FLA_NO_TRANSPOSE, ONE, A1, ONE, C11 );
} /* end task */
  /*-----*/
  FLA_Cont_with_3x3_to_2x2( &CTL, /**/ &CTR, C00, C01, /**/ C02,
                           C10, C11, /**/ C12,
                           /***/ /***/
                           &CBL, /**/ &CBR, C20, C21, /**/ C22,
                           FLA_TL );

  FLA_Cont_with_3x1_to_2x1( &AT,           A0,
                           A1,
                           /* ** */      /* ** */
                           &AB,           A2,   FLA_TOP );
}
} /* end of taskq */

```

```
#pragma intel omp parallel taskq
```

```
{  
while ( FLA_Obj_length( AT ) < FLA_Obj_length( A ) ){  
    b = min( FLA_Obj_length( AB ), nb_alg );  
  
    FLA_Repart_2x1_to_3x1( AT,                &A0,  
                          /* ** */          /* ** */  
                          AB,                &A1,  
                          &A2,                b, FLA_BOTTOM );  
    FLA_Repart_2x2_to_3x3( CTL, /**/ CTR,     &C00, /**/ &C01, &C02,  
                          /* ***** */ /* ***** */  
                          &C10, /**/ &C11, &C12,  
                          CBL, /**/ CBR,     &C20, /**/ &C21, &C22, b, b, FLA_BR );
```

```
/*-----*/
```

```
#pragma intel omp task captureprivate(A2, A1, C11, C21)
```

```
{  
    FLA_Gemm( FLA_NO_TRANSPOSE, FLA_TRANSPOSE, ONE, A2, A1, ONE, C21 );  
    FLA_Syrk( FLA_LOWER_TRIANGULAR, FLA_NO_TRANSPOSE, ONE, A1, ONE, C11 );
```

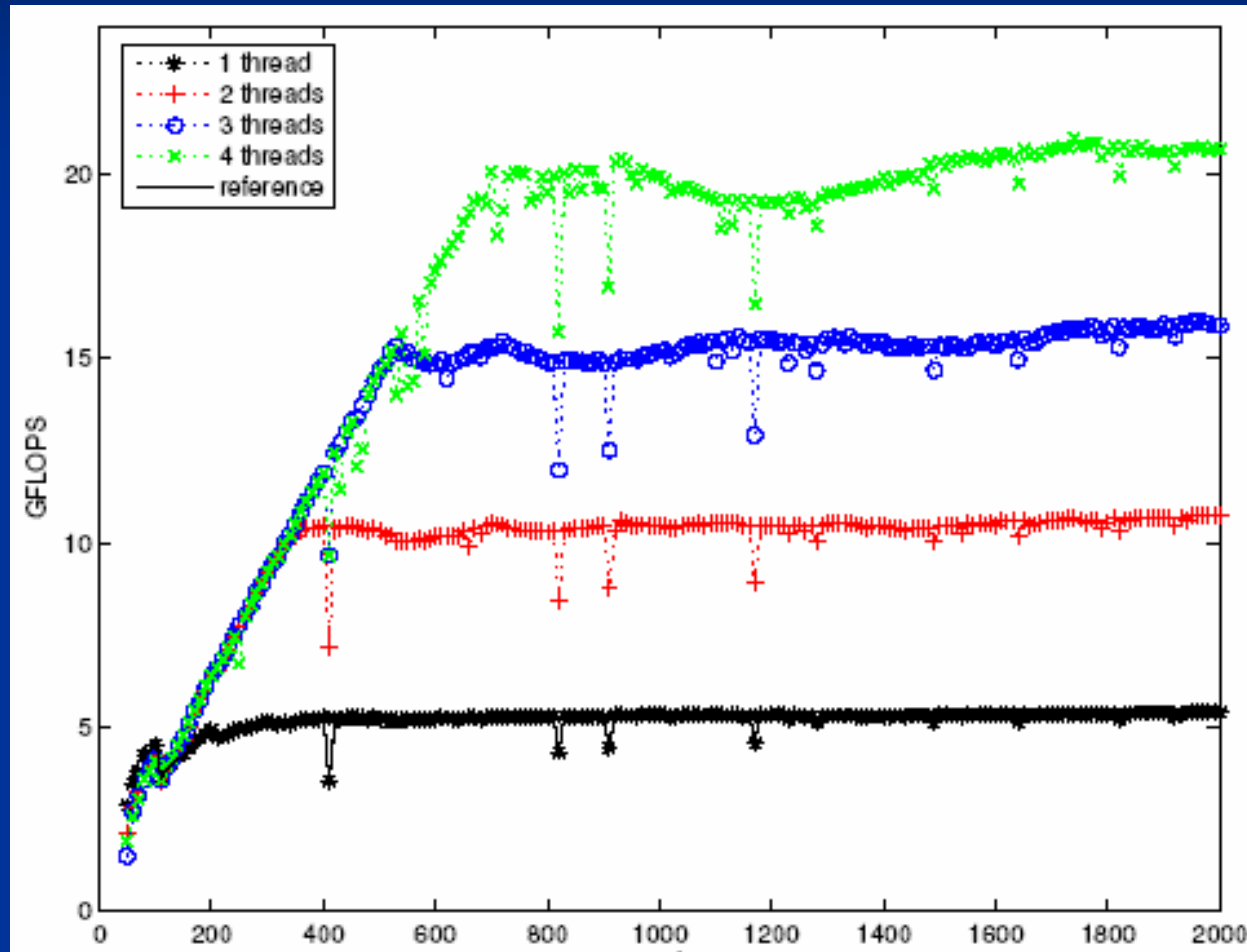
```
}
```

```
/*-----*/
```

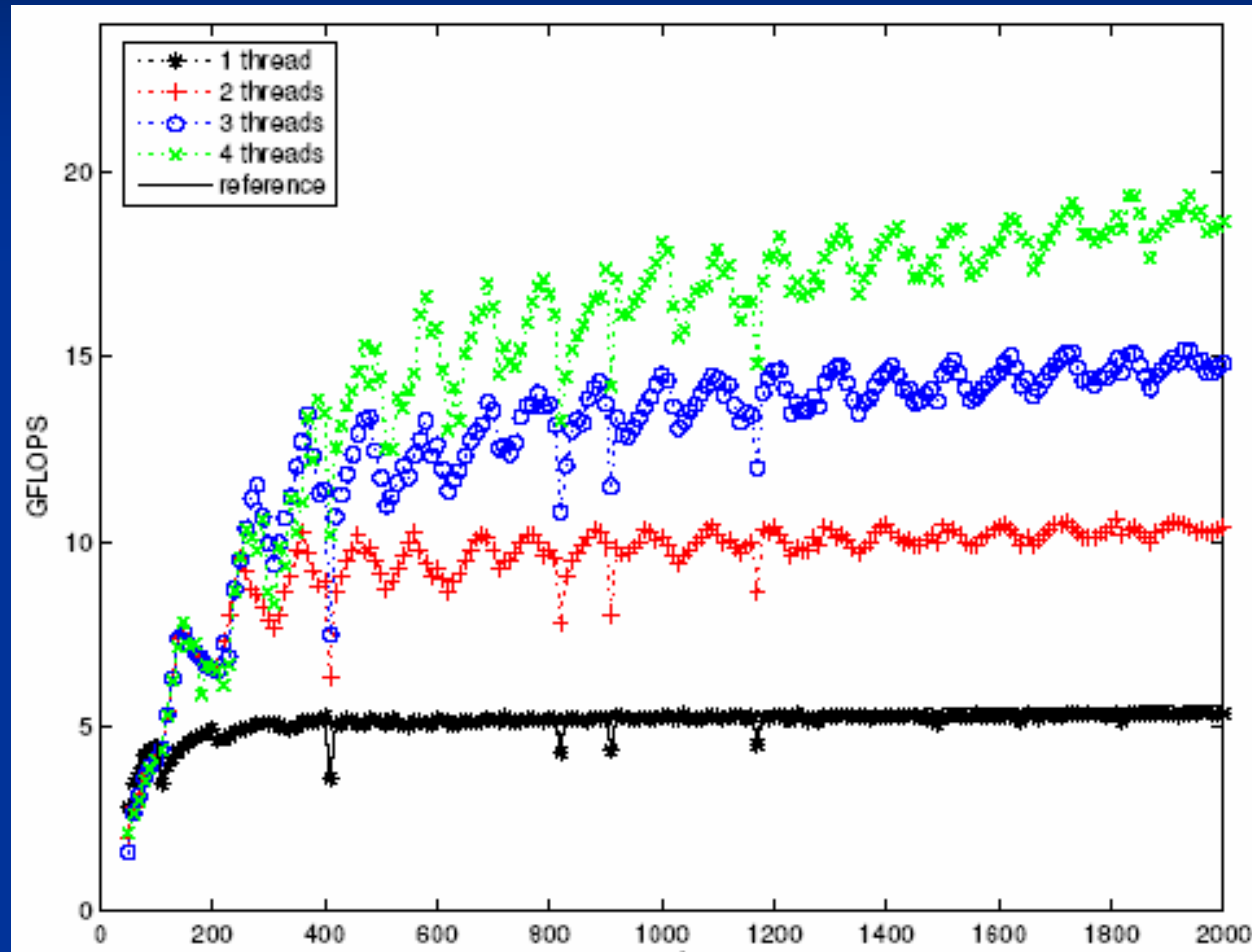
```
FLA_Cont_with_3x1_to_2x1( &AT,                A0,  
                          A1,  
                          /* ** */          /* ** */  
                          &AB,                A2,                FLA_TOP );  
FLA_Cont_with_3x3_to_2x2( &CTL, /**/ &CTR,     C00, C01, /**/ C02,  
                          C10, C11, /**/ C12,  
                          /* ***** */ /* ***** */  
                          &CBL, /**/ &CBR,     C20, C21, /**/ C22,      FLA_TL );
```

```
}
```

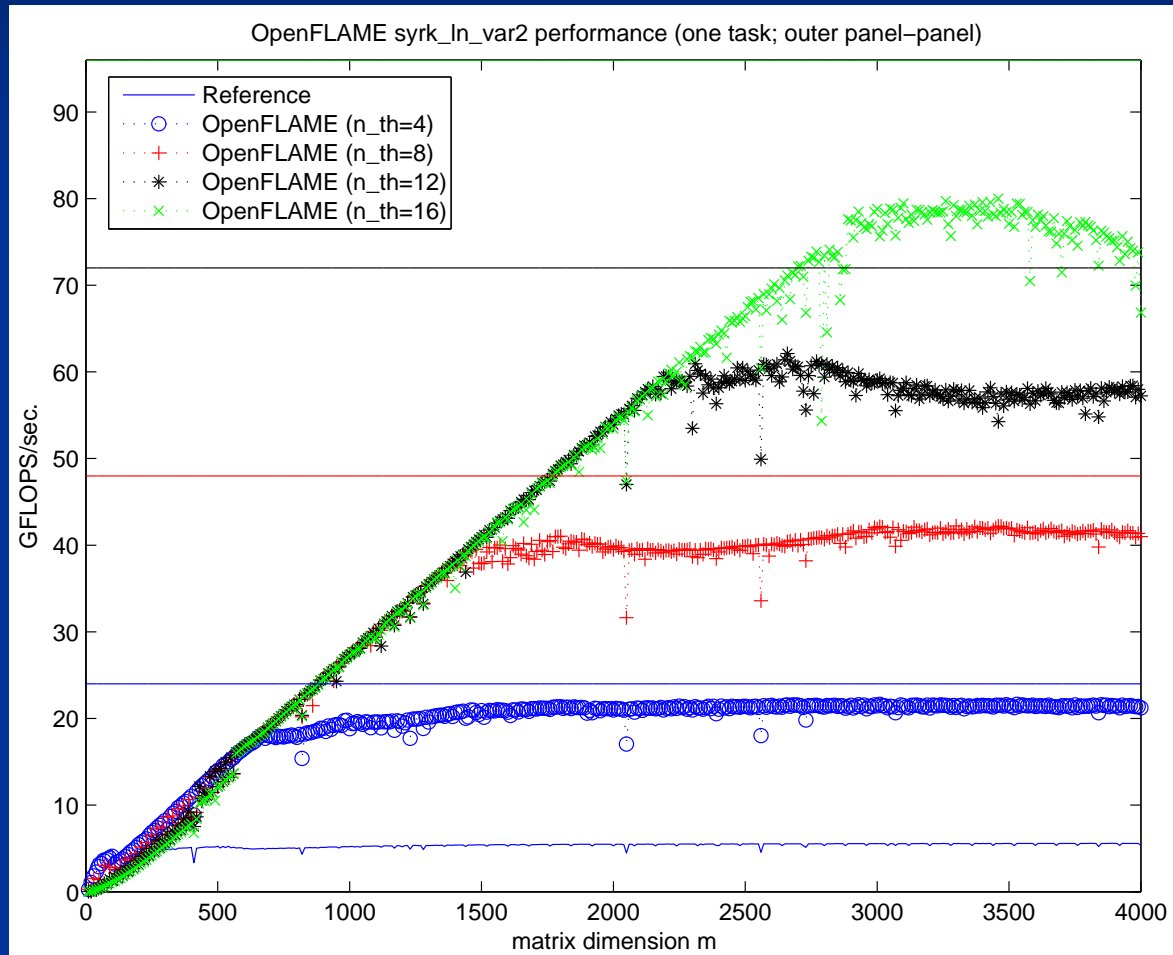
Performance HP 4CPU Itanium2 dsyrk variant 2



Performance HP 4CPU Itanium2 dsyrk variant 3



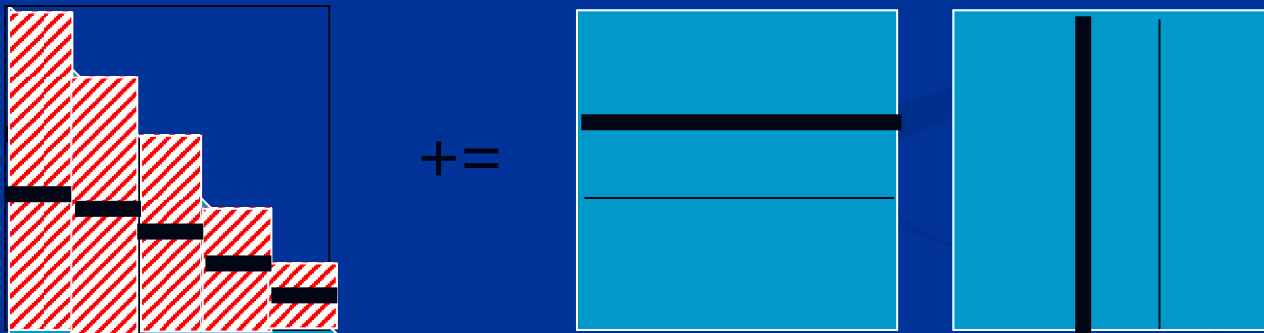
Performance on NEC 16 CPU Itanium2 system (1.5 GHz)



Work queuing

n Example:

n Symmetric rank-k update




```

#pragma intel omp parallel taskq
{
while ( FLA_Obj_length( AT ) < FLA_Obj_length( A ) ){

    b = min( FLA_Obj_length( AB ), nb_alg );

    FLA_Repart_2x1_to_3x1( AT,          &A0,
                          /* ** */    /* ** */
                          AB,          &A1,
                                          &A2,          b, FLA_BOTTOM );

    FLA_Repart_2x2_to_3x3( CTL, /**/ CTR,    &C00, /**/ &C01, &C02,
                          /* ***** */ /* ***** */
                          &C10, /**/ &C11, &C12,
                          CBL, /**/ CBR,    &C20, /**/ &C21, &C22,
                          b, b, FLA_BR );

    /*-----*/

    b2 = FLA_Obj_length( A2 )/2;

    FLA_Part_2x1( A2,          &A2_T,
                  &A2_B,          b2, FLA_TOP );

    FLA_Part_2x1( C21,        &C21_T,
                  &C21_B,          b2, FLA_TOP );

    /*-----*/
    #pragma intel omp task captureprivate(A2_T, A1, C21_T)
    {

        FLA_Gemm( FLA_NO_TRANSPOSE, FLA_TRANSPOSE,
                  ONE, A2_T, A1, ONE, C21_T );

    }
    #pragma intel omp task captureprivate(A2_B, A1, C21_B)
    {

        FLA_Gemm( FLA_NO_TRANSPOSE, FLA_TRANSPOSE,
                  ONE, A2_B, A1, ONE, C21_B );

    }

    /*-----*/
    /*-----*/

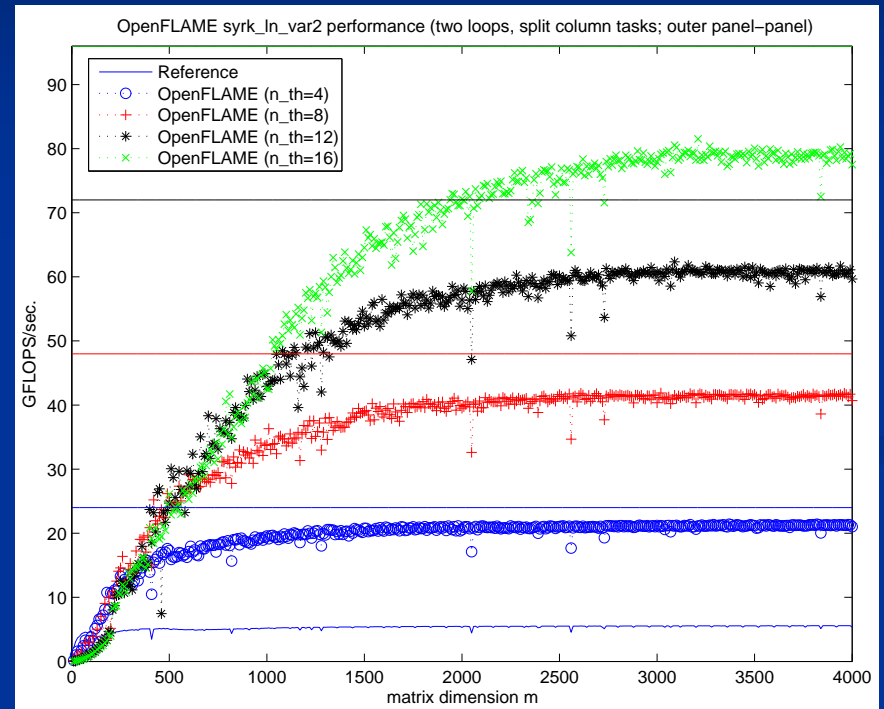
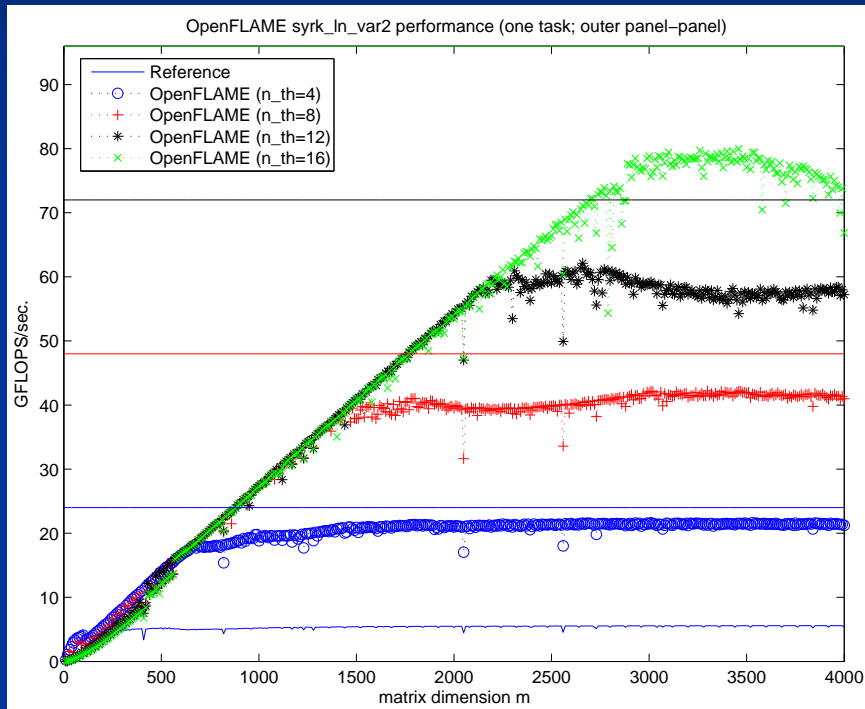
    FLA_Cont_with_3x1_to_2x1( &AT,          A0,
                              A1,
                              /* ** */    /* ** */
                              &AB,          A2,          FLA_TOP );

    FLA_Cont_with_3x3_to_2x2( &CTL, /**/ &CTR,    C00, C01, /**/ C02,
                              C10, C11, /**/ C12,
                              /* ***** */ /* ***** */
                              &CBL, /**/ &CBR,    C20, C21, /**/ C22,
                              FLA_TL );

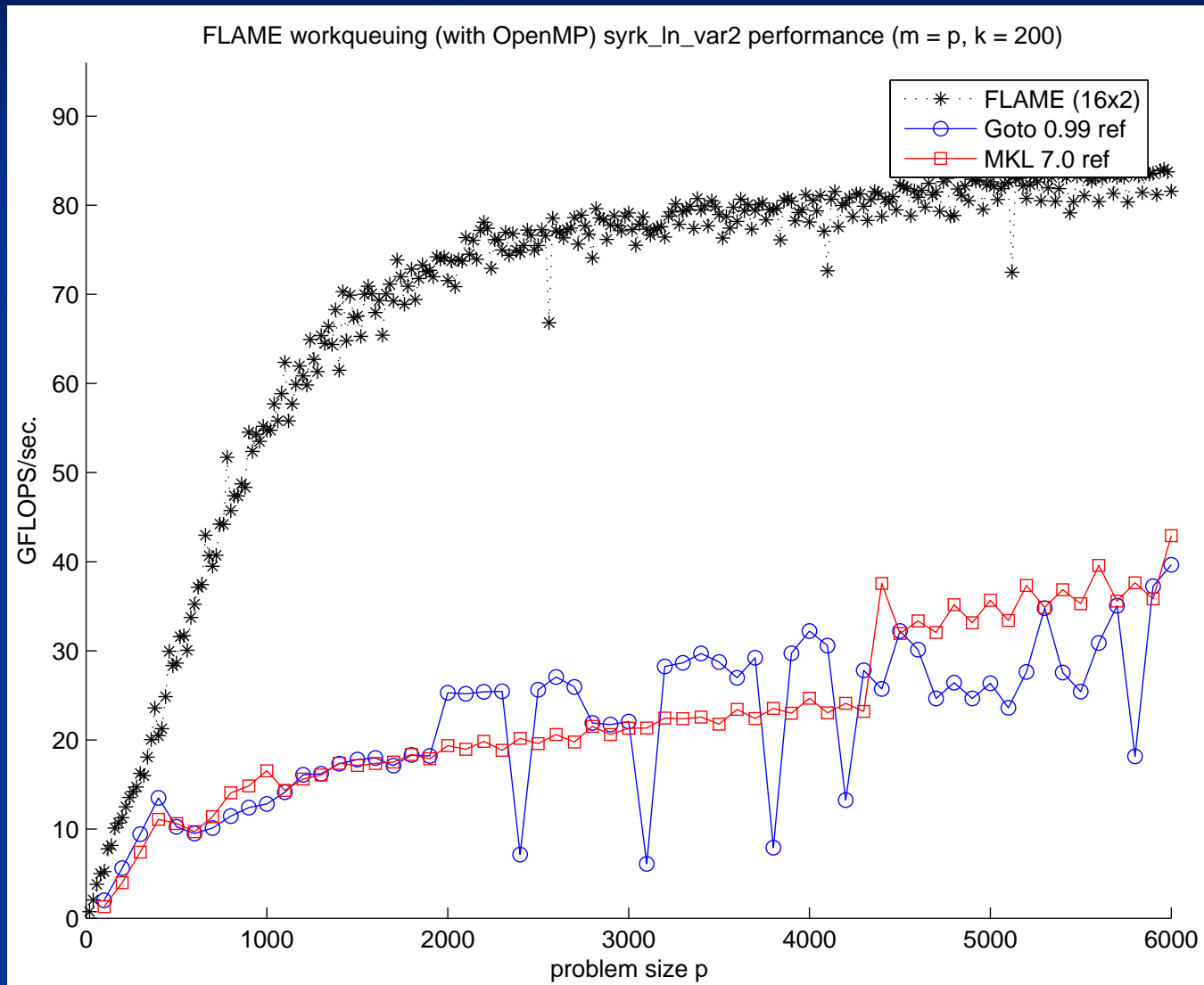
}

```

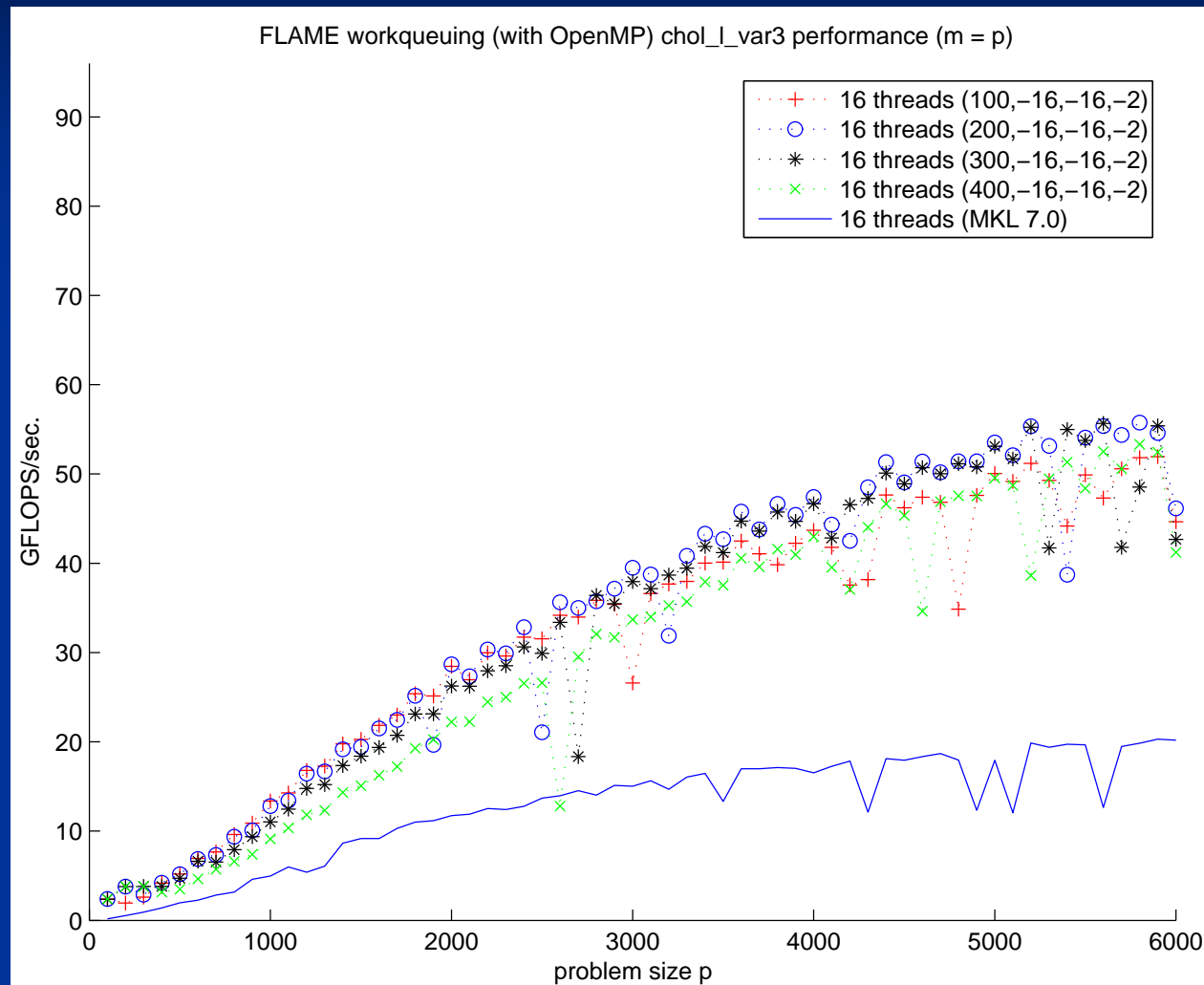
Performance on NEC 16 CPU Itanium2 system (1.5 GHz)



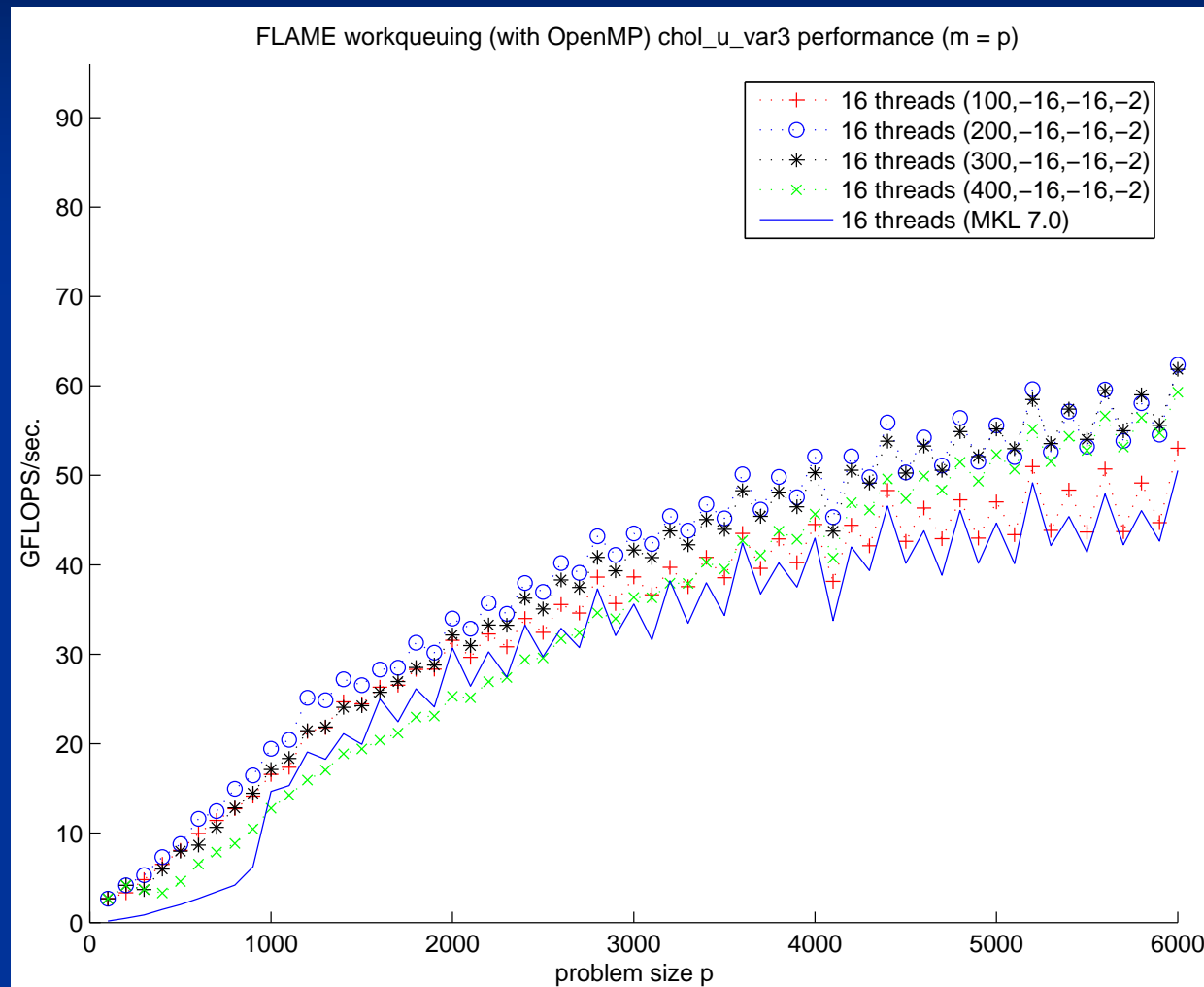
Syrk performance



Cholesky Factorization Performance

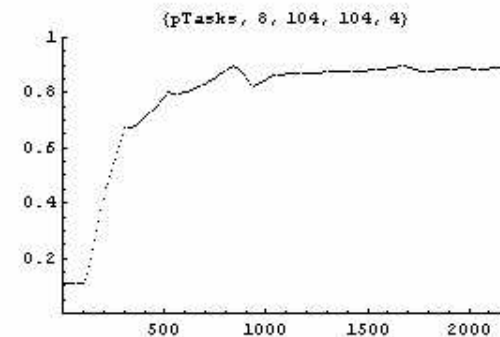
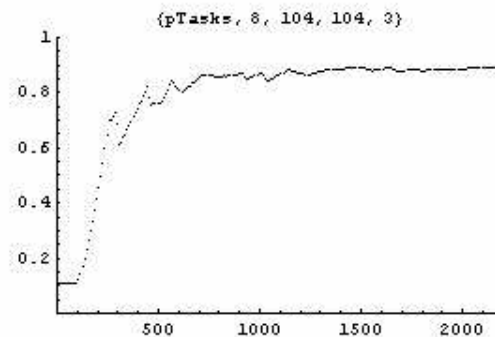
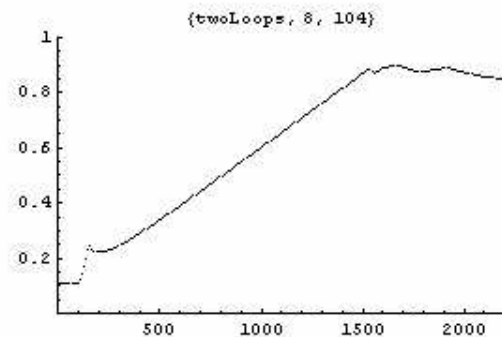
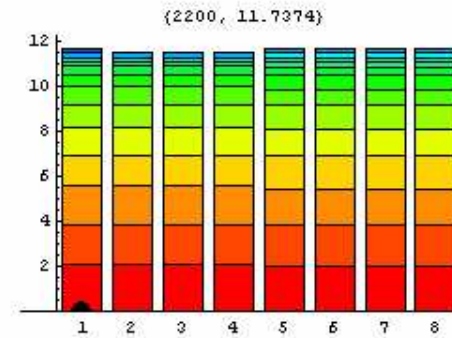
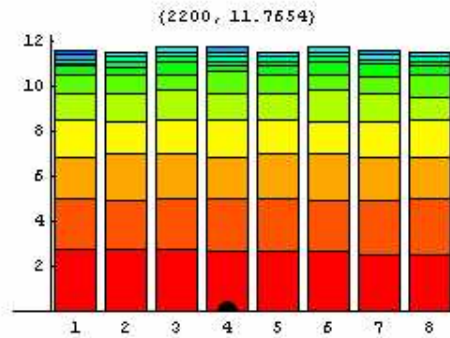
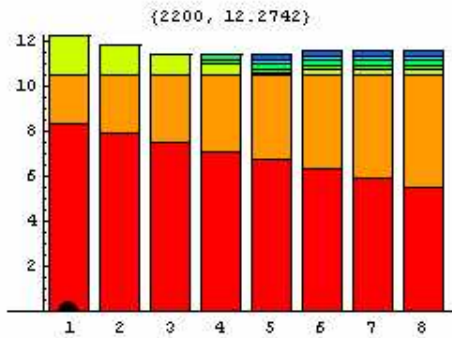
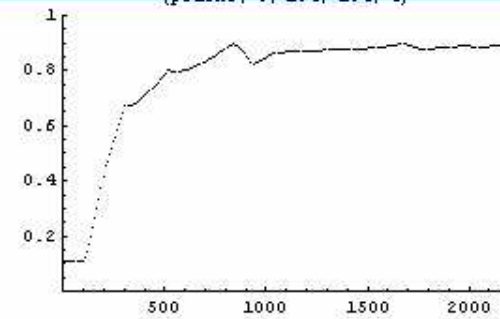
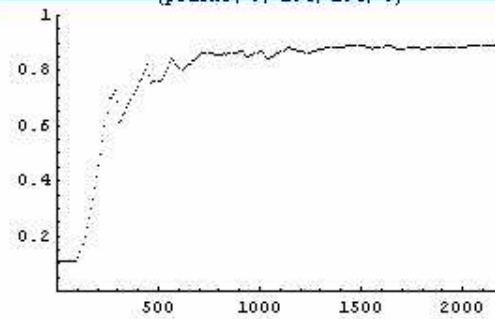
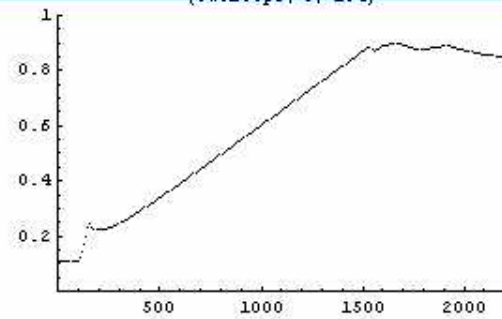


Cholesky Factorization Performance



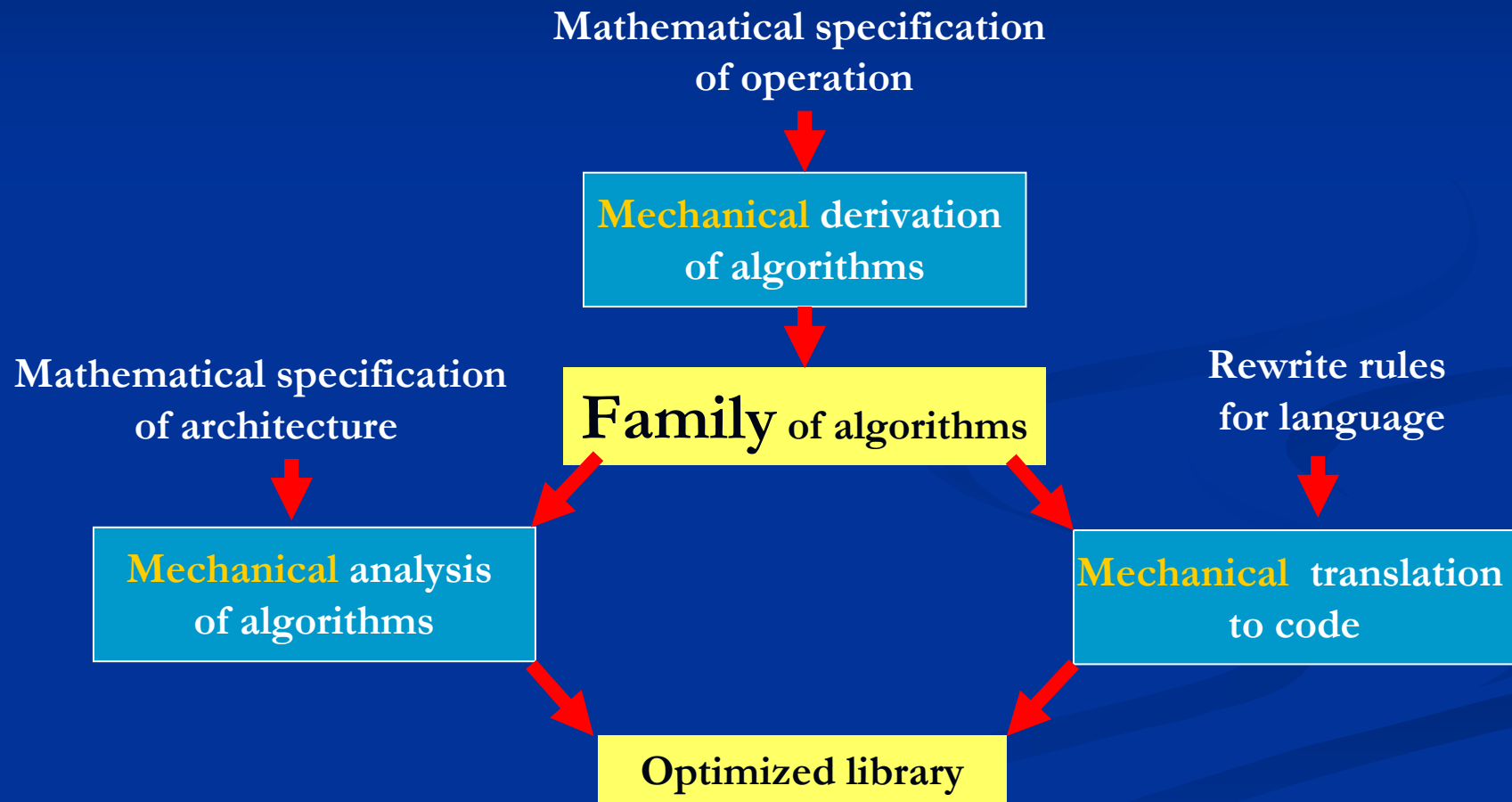
Switch to Demo

[Skip demo slides](#)



```
Show[
GraphicsArray[{
  (ListPlot[Table[{n, maxIndex[schedule[oneTask[n], 8]]}, {n, 300, 2000, 1}],
    PlotRange -> {0, 8.5}, PlotStyle -> PointSize[0.004], PlotLabel -> "1 Task", DisplayFunction -> Identity],
```

Conclusion



Conclusion

- n Mechanical generation of libraries supports
 - n New architectures
 - n Architectures currently supported: Sequential, SMP parallel, distributed memory parallel
 - n New languages
 - n Languages supported: LaTeX, C, Matlab, Fortran, C+MPI, C+OpenMP, Mathematica, Haskell, LabView's G
 - n New datastructures
 - n Datastructures supported: column-major storage, banded, dense stored by blocks, sparse hierarchical, out-of-core
 - n New operations
 - n Mechanical derivation of algorithms for all BLAS3, LAPACK, many operations in control theory

More Information

<http://www.cs.utexas.edu/users/flame/>

<http://www.cs.utexas.edu/users/flame/pubs.html>

rvdg@cs.utexas.edu

What is needed?

- n Project so far concentrates on the science that supports the approach and prototyping of tools
- n Full-blown integration of all tools requires
 - n Full-time postdoc
 - n Full-time professional programmer
 - n Part-time web developer
 - n Several graduate students
 - n Collaboration with the Texas Advanced Computing Center