

**NIST-Recommended**  
**Random Number Generator**  
**Based on ANSI X9.31 Appendix A.2.4**  
**Using the 3-Key Triple DES and AES**  
**Algorithms**

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## 1 Introduction

This document, the “NIST-Recommended Random Number Generator Based on ANSI X9.31 Appendix A.2.4 Using the 3-Key Triple DES and AES Algorithms” specifies the NIST-recommended addition to the underlying document *Digital Signatures Using Reversible Public Key Cryptography for the Financial Services Industry (rDSA)*, ANSI X9.31-1988, random number generator. It specifies how to use 3-Key Triple DES and AES as the core of the X9.31 RNG.

## 2 ANSI X9.31 Appendix A.2.4 Using 3-Key Triple DES

Let  $\text{ede}^*X(Y)$  represent the DEA multiple encryption of  $Y$  under the key  $*X$ .  
Let  $*K$  be 3-key Triple DES, 3 64 bit keys.

This  $*K$  is reserved only for the generation of pseudo random numbers.

Let  $V$  be a 64-bit seed value which is also kept secret, and XOR be the exclusive-or operator. Let  $DT$  be a date/time vector which is updated on each iteration.  $I$  is an intermediate value. A vector  $R$  is generated as follows (Note for Triple DES implementations:  $DT$ ,  $I$  and  $R$  are 64-bits each.):

$I = \text{ede}^*K(DT)$

$R = \text{ede}^*K(I \text{ XOR } V)$  and a new  $V$  is generated by  $V = \text{ede}^*K(R \text{ XOR } I)$ .

## 3 ANSI X9.31 Appendix A.2.4 Using AES

Let  $\text{ede}^*X(Y)$  represent the AES encryption of  $Y$  under the key  $*X$ .

For AES 128-bit key, let  $*K$  be a 128 bit key.

For AES 192-bit key, let  $*K$  be a 192 bit key.

For AES 256-bit key, let  $*K$  be a 256 bit key.

This  $*K$  is reserved only for the generation of pseudo random numbers.

Let  $V$  be a 128-bit seed value which is also kept secret, and XOR be the exclusive-or operator. Let  $DT$  be a date/time vector which is updated on each iteration.  $I$  is an intermediate value. A vector  $R$  is generated as follows (Note for AES implementations  $DT$ ,  $I$ , and  $R$  are 128-bits each.):

$I = \text{ede}^*K(DT)$

$R = \text{ede}^*K(I \text{ XOR } V)$  and a new  $V$  is generated by  $V = \text{ede}^*K(R \text{ XOR } I)$ .

## **Appendix A      References**

- [1]    *Digital Signatures Using Reversible Public Key Cryptography for the Financial Services Industry (rDSA)*, ANSI X9.31-1988, September 1998.